

# Bounding the set of finite dimensional correlations

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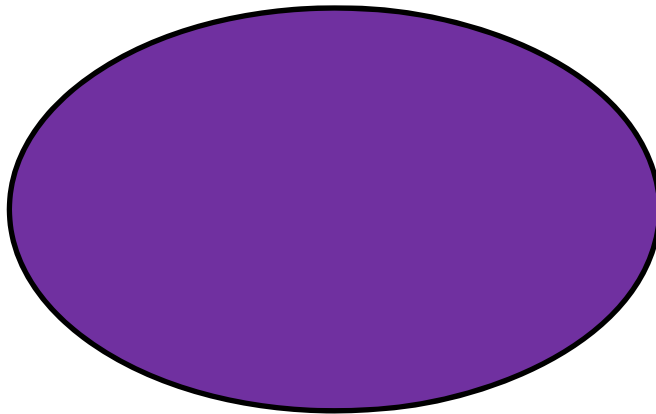
MN and T. Vértesi, Phys. Rev. Lett. 115, 020501 (2015) .

MN, A. Feix, M. Araújo and T. Vértesi, Phys. Rev. A 92, 042117 (2015).

# Optimization theory

$$\max f(x)$$

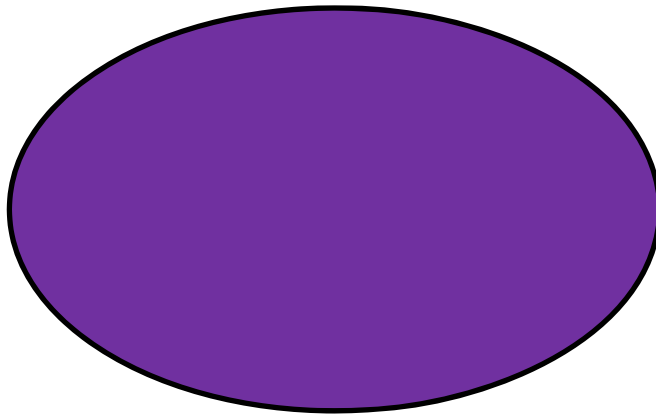
$x \in$



Variational methods

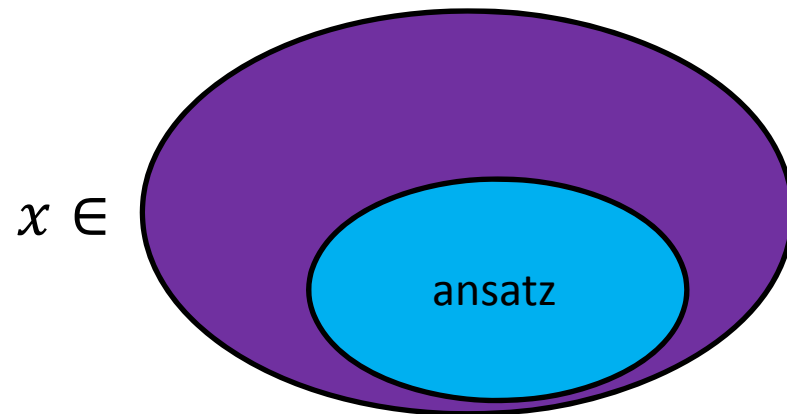
$$\max f(x)$$

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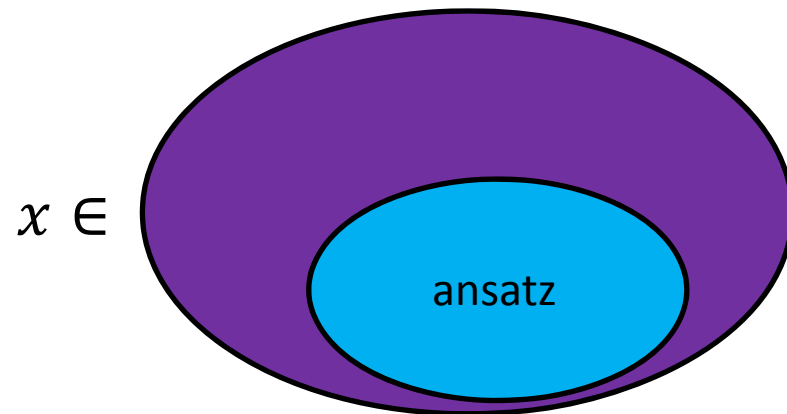
# Variational methods

$$\max f(x)$$



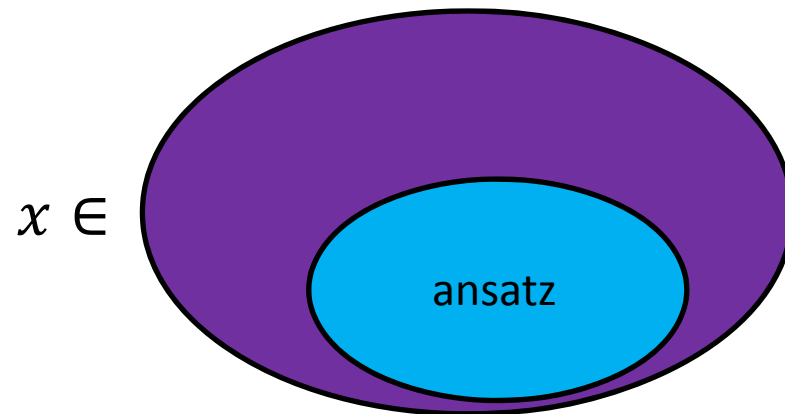
## Variational methods

$$\max f(x) \geq \mu$$



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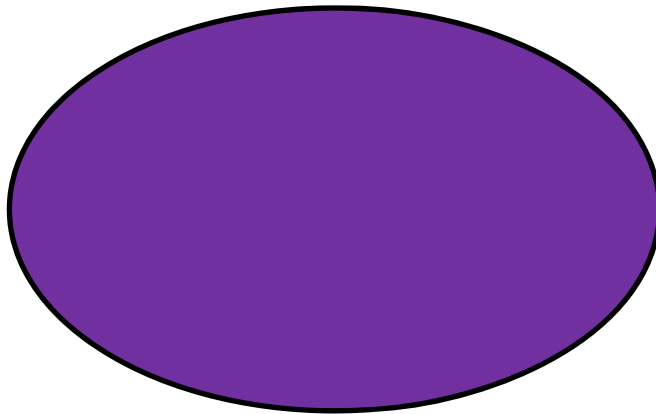


- usually, pretty straightforward
  - usually, universal
  - sometimes arise as intuitions gathered via numerical experiments. E.g.: MPS.
- } E.g.: Newton's method

Relaxations

$$\max f(x)$$

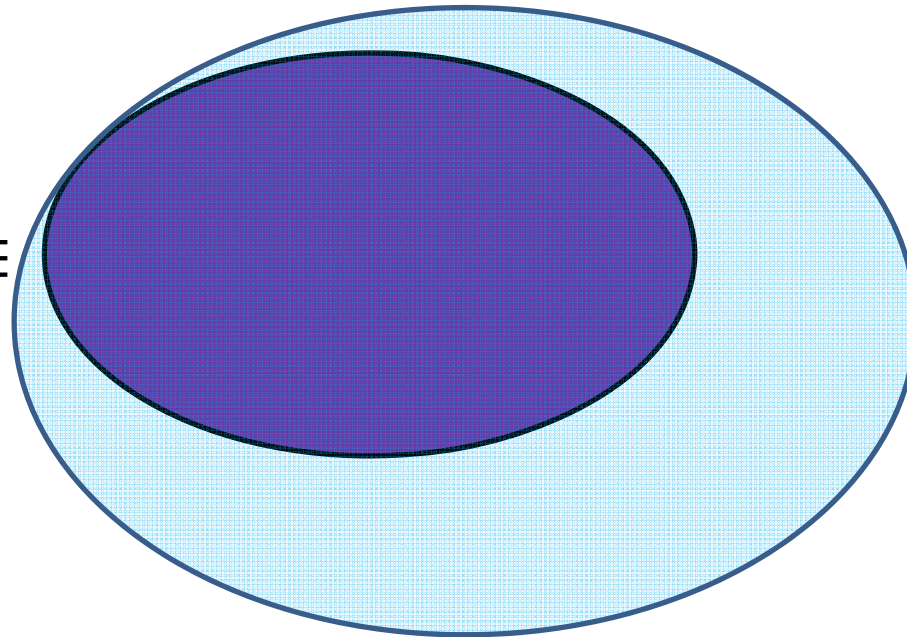
$x \in$



# Relaxations

$$\max f(x)$$

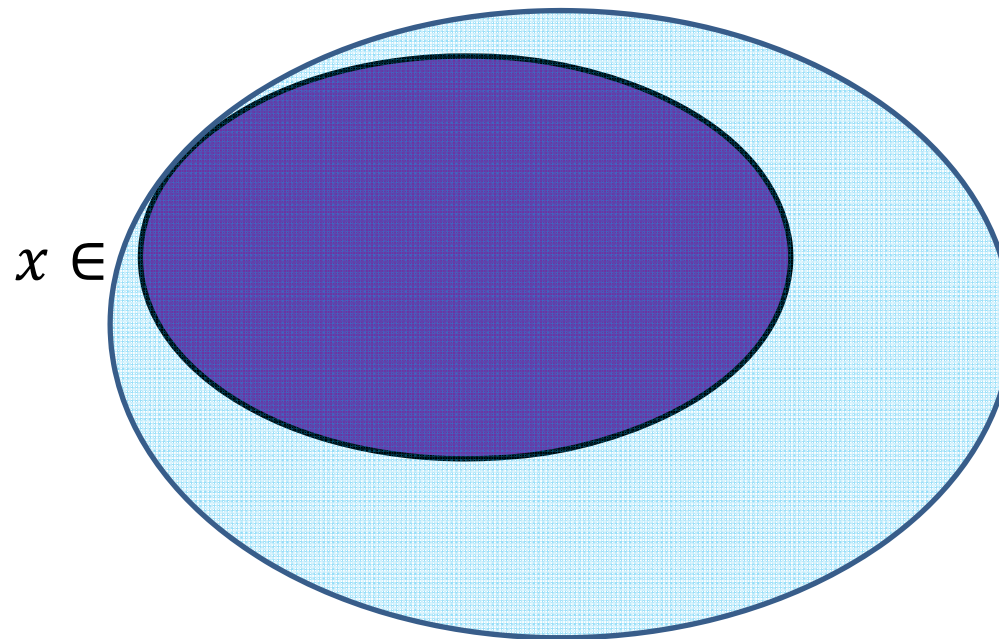
$x \in$





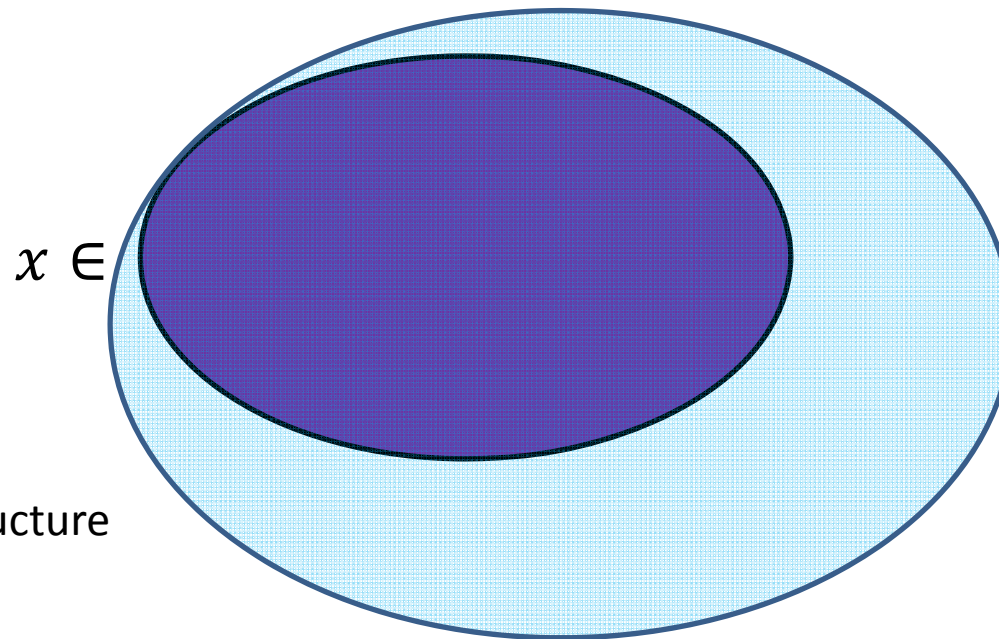
## Relaxations

$$\lambda \geq \max f(x)$$



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$$\lambda \geq \max f(x)$$



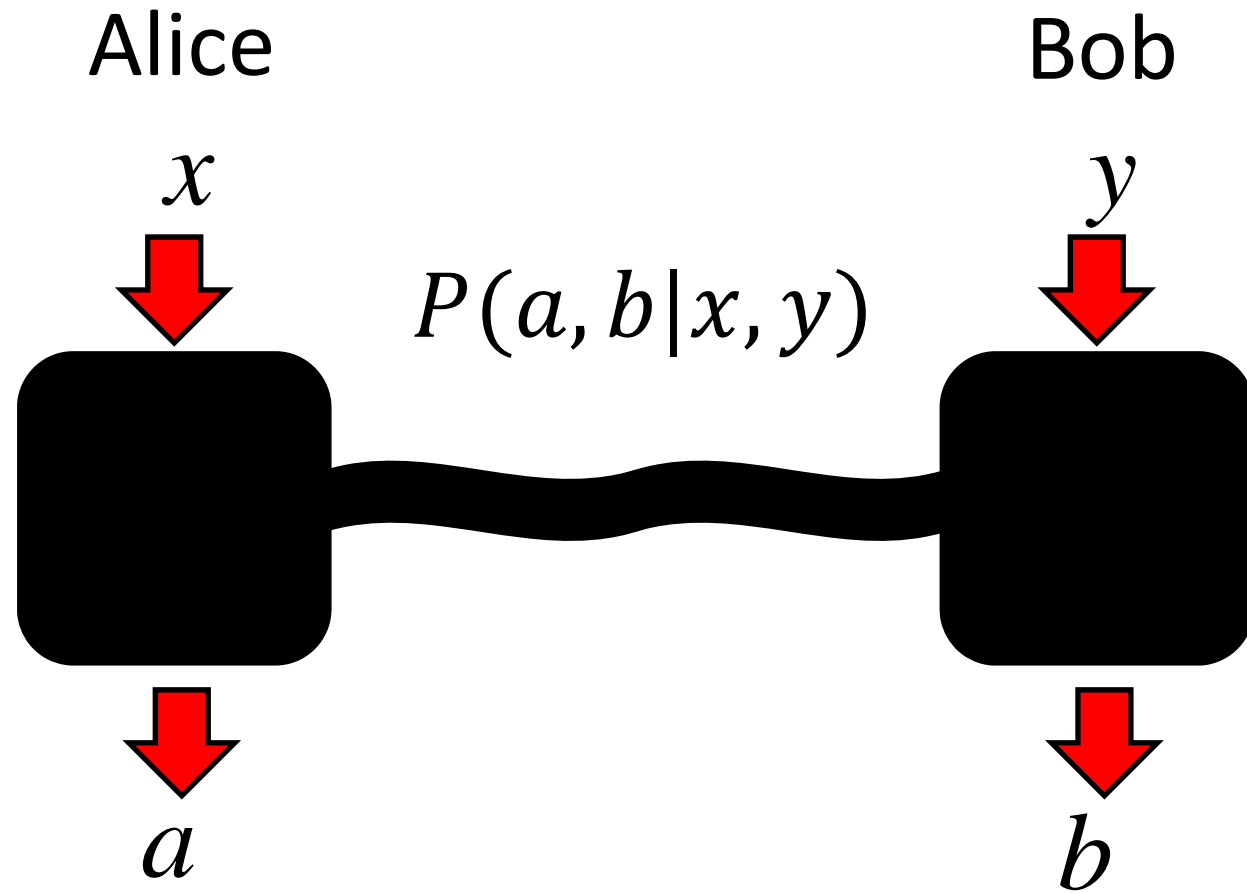
-Depend on the structure  
of the problem

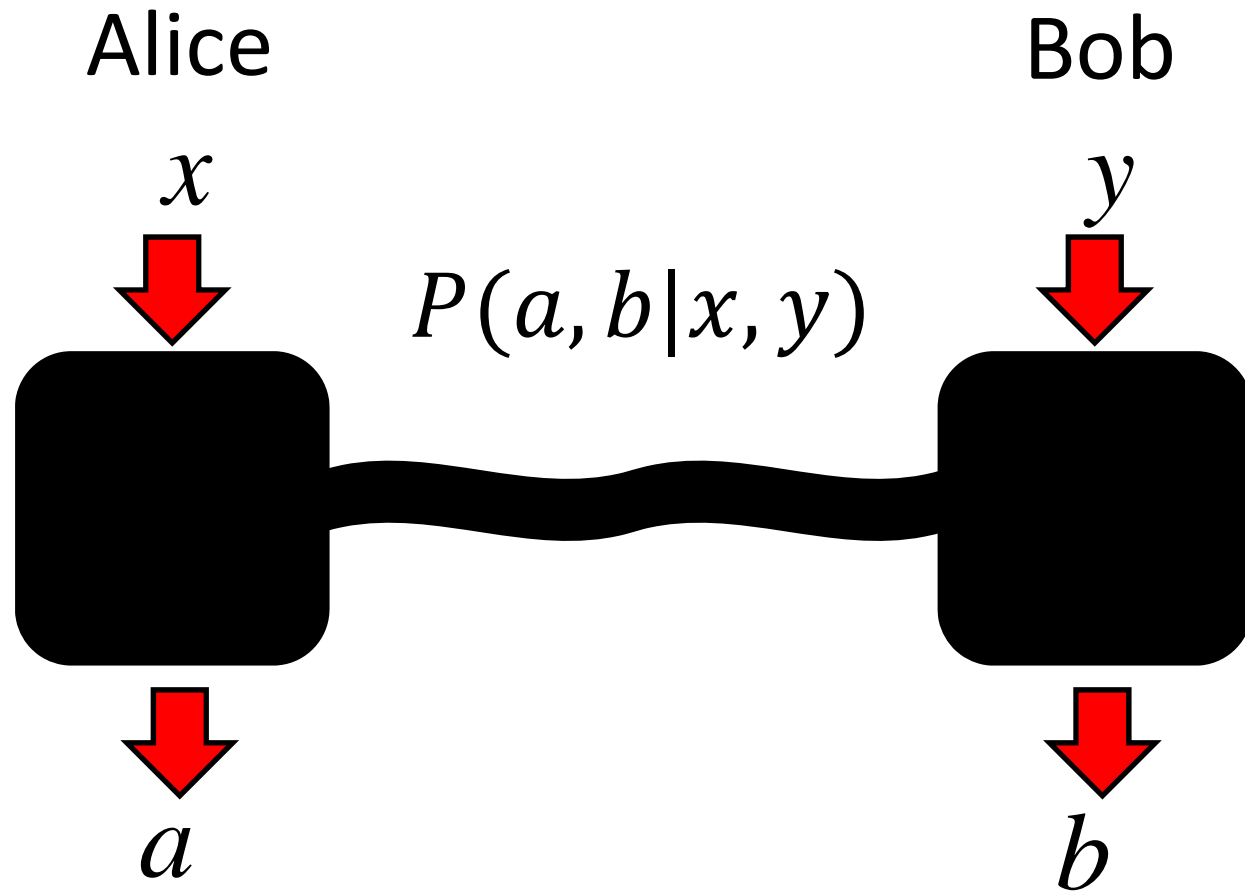
-"Aha!" idea

E.g.: the NPA hierarchy

MN, S. Pironio and A. Acín, Phys. Rev. Lett. 98, 010401 (2007).

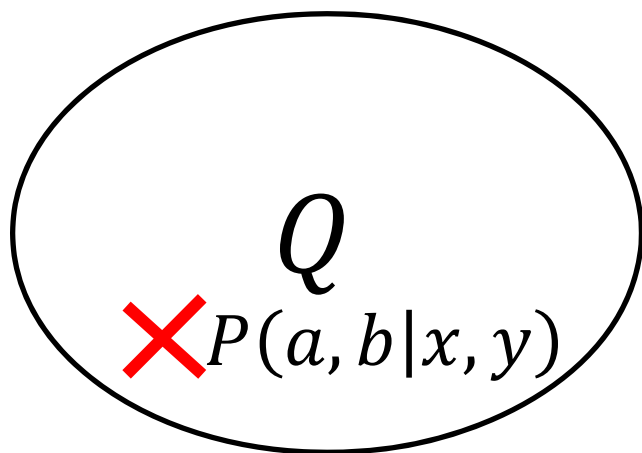
MN, S. Pironio and A. Acín, New J. Phys. 10, 073013 (2008).





Are the outputs of this experiment compatible with quantum mechanics?

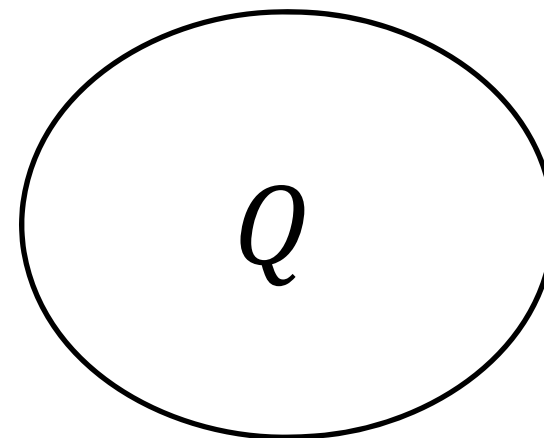
"Easy"



Variational methods

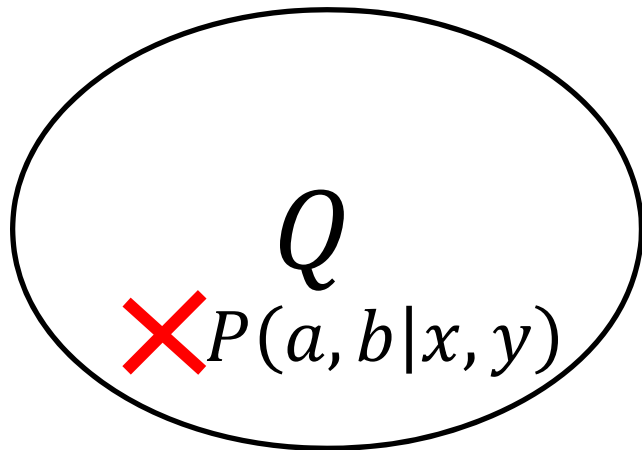
$P(a, b|x, y)$

"Difficult"



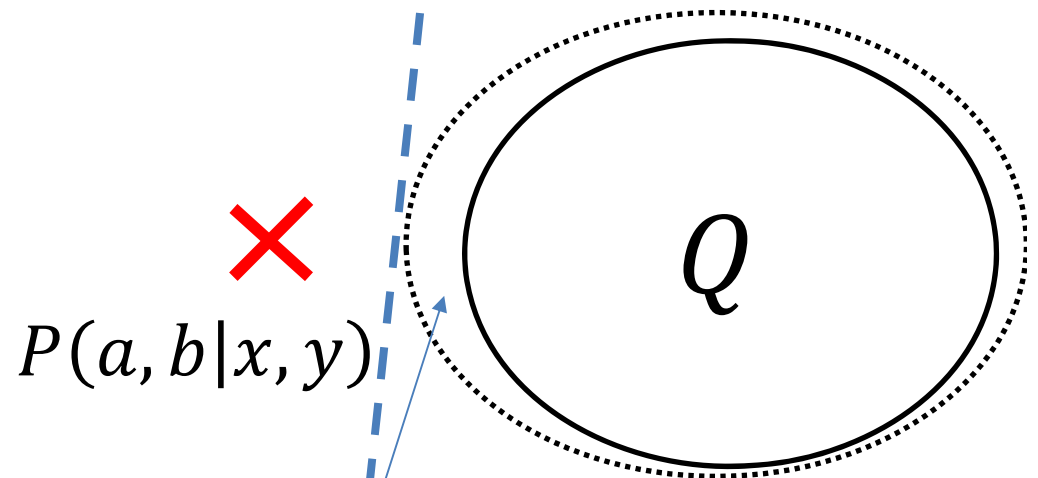
Mathematical  
coincidences

"Easy"



Variational methods

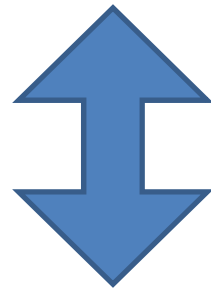
"Difficult"



Mathematical  
coincidences

relaxation?

$P(a, b|x, y)$  is not quantum

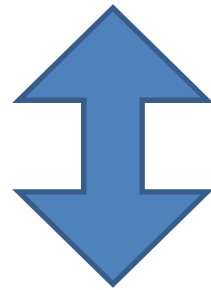


There do not exist a Hilbert space  $\mathcal{H}$ , a quantum state  $|\psi\rangle \in \mathcal{H}$  and projector operators  $\{E_a^x, F_b^y\} \subset B(\mathcal{H})$ , with  $\sum_a E_a^x = \sum_b F_b^y = \mathbb{I}$ ,  $[E_a^x, F_b^y] = 0$ , such that  $P(a, b|x, y) = \langle\psi|E_a^x F_b^y|\psi\rangle$ .

The problem resists brute force approach!!



$P(a, b|x, y)$  is not quantum



There do not exist a Hilbert space  $\mathcal{H}$ , a quantum state  $|\psi\rangle \in \mathcal{H}$  and projector operators  $\{E_a^x, F_b^y\} \subset B(\mathcal{H})$ , with  $\sum_a E_a^x = \sum_b F_b^y = \mathbb{I}$ ,  $[E_a^x, F_b^y] = 0$ , such that  $P(a, b|x, y) = \langle\psi|E_a^x F_b^y|\psi\rangle$ .

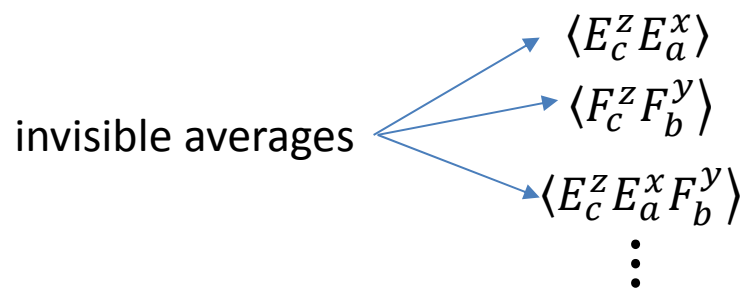
The problem resists brute force approach!!

What is a Hilbert space?

$$\psi = (\psi_1, \psi_2, \psi_3 \dots)$$

$$\sum_i |\psi_i|^2 < \infty$$

$$P(a, b|x, y) = \langle \phi, E_a^x F_b^y \phi \rangle = \langle E_a^x F_b^y \rangle$$



visible averages

Moment matrix

$$\Gamma = \begin{array}{c|cccc} & \mathbb{I} & E_a^x & E_{a'}^{x'} & F_b^y \dots \\ \hline \mathbb{I}^\dagger & 1 & & & \\ E_a^{x\dagger} & & \langle E_a^x E_{a'}^{x'} \rangle & & P(a, b|x, y) \\ E_{a'}^{x'\dagger} & & & & P(a', b|x', y) \\ F_b^{y\dagger} & & & & \\ \vdots & & & & \end{array}$$

$$\Gamma_{u,v} = \langle u^\dagger v \rangle$$

Moment matrix

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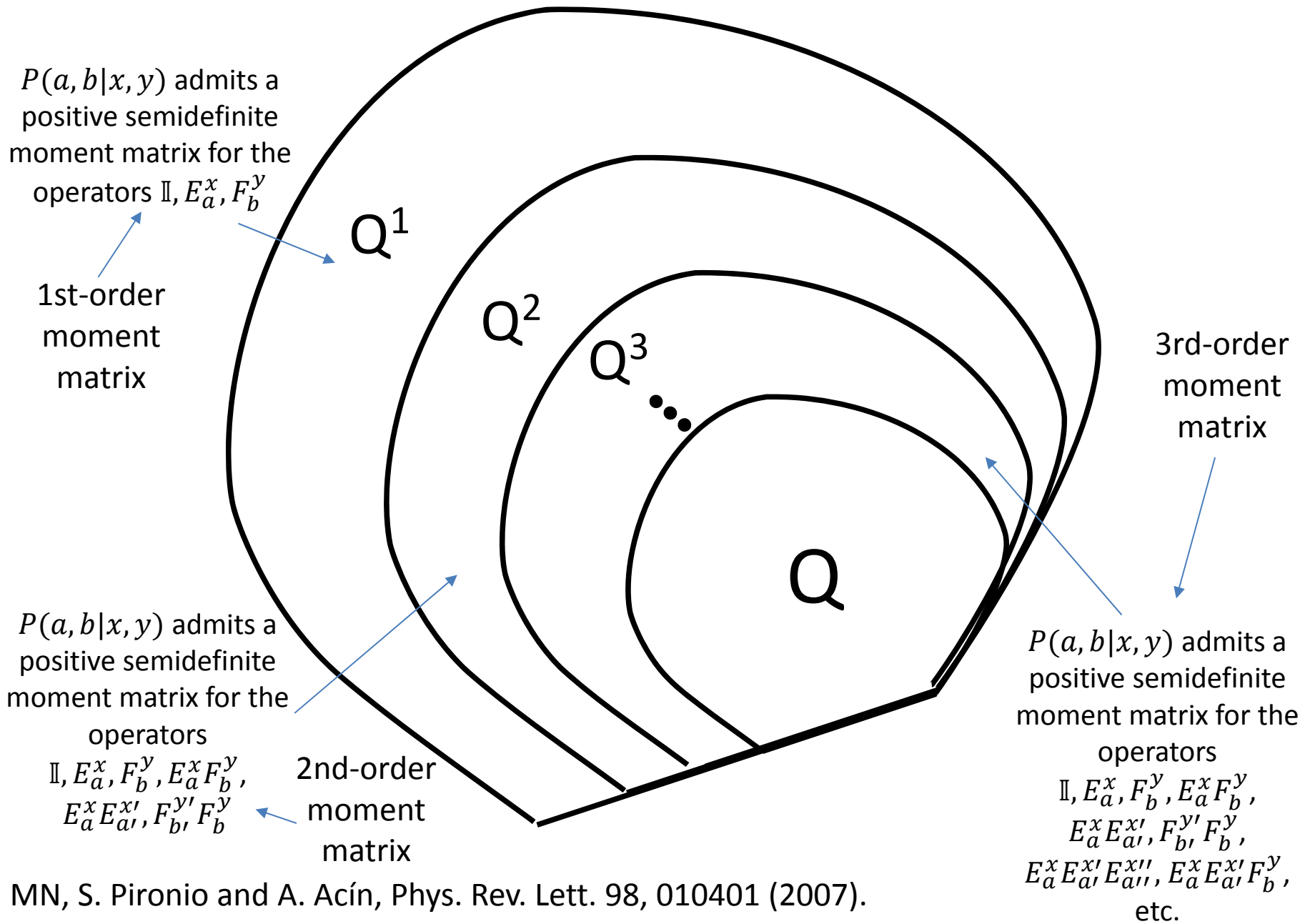
Hilbert space



Moment matrix

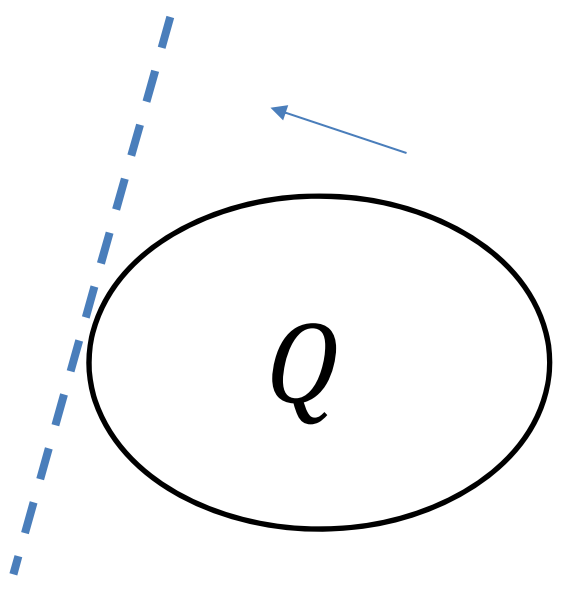
	$\mathbb{I}$	$E_a^x$	$E_{a'}^{x'}$	$F_b^y$	...
$\mathbb{I}^\dagger$	1				
$E_a^{x\dagger}$			$\langle E_a^x E_{a'}^{x'} \rangle$	$P(a, b x, y)$	
$E_{a'}^{x'\dagger}$				$P(a', b x', y)$	
$F_b^{y\dagger}$					
$\vdots$					

$\geq 0$



MN, S. Pironio and A. Acín, Phys. Rev. Lett. 98, 010401 (2007).  
 MN, S. Pironio and A. Acín, New J. Phys. 10, 073013 (2008).





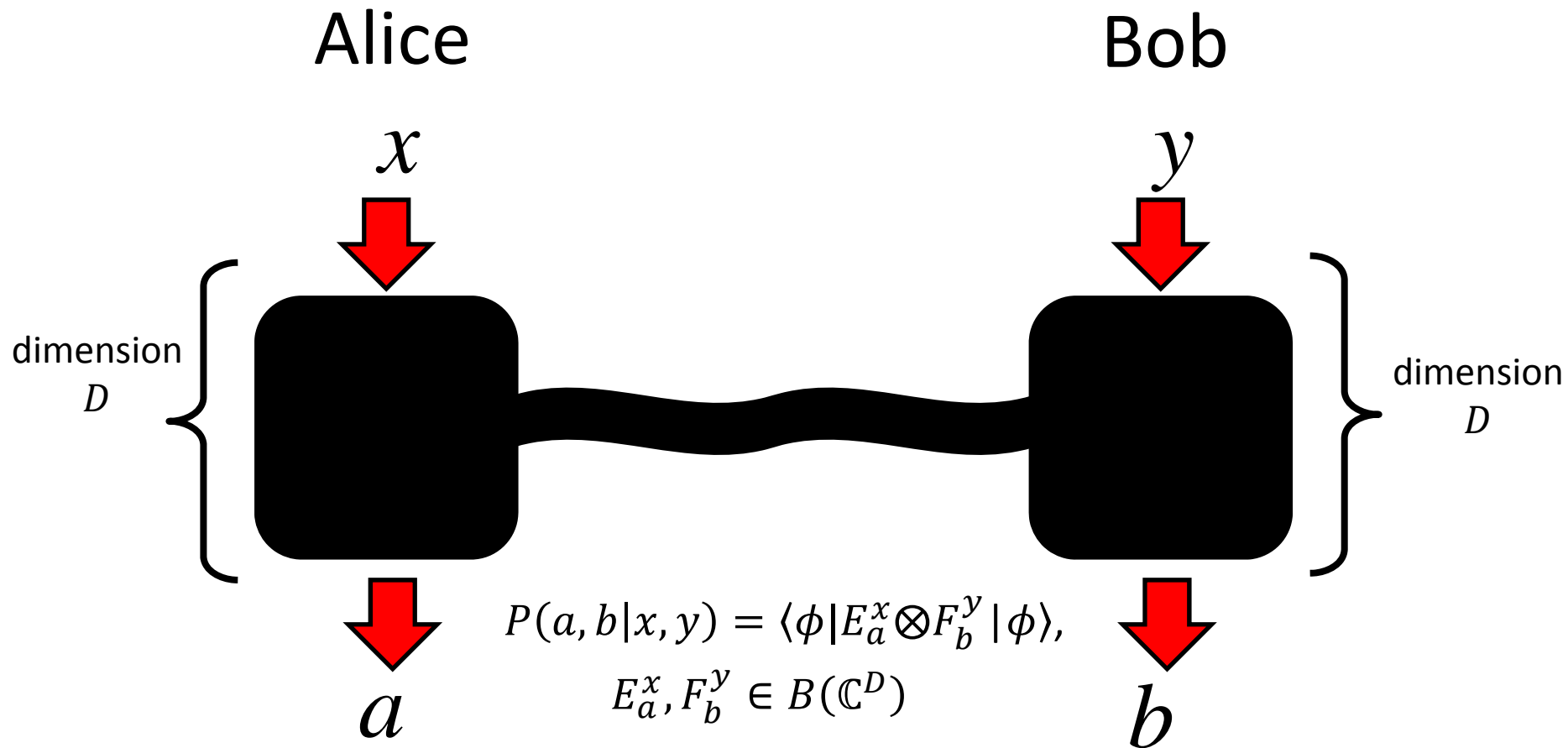
Variational methods

NPA

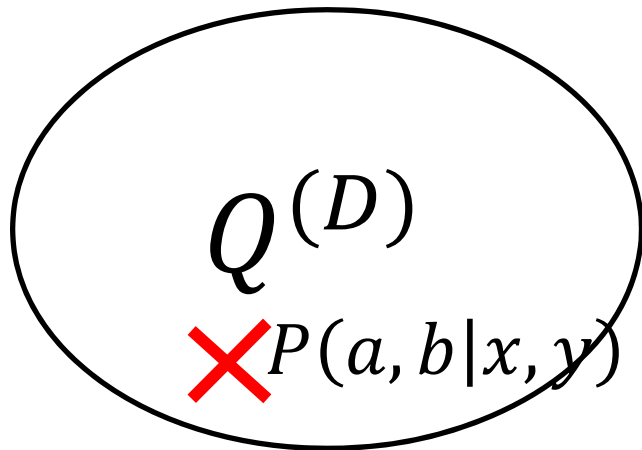
$$L \leq \max \sum c_{abxy} P(a, b|x, y) \leq U$$

# Bounding the set of finite dimensional quantum correlations

MN and T. Vértesi, Phys. Rev. Lett. 115, 020501 (2015) .

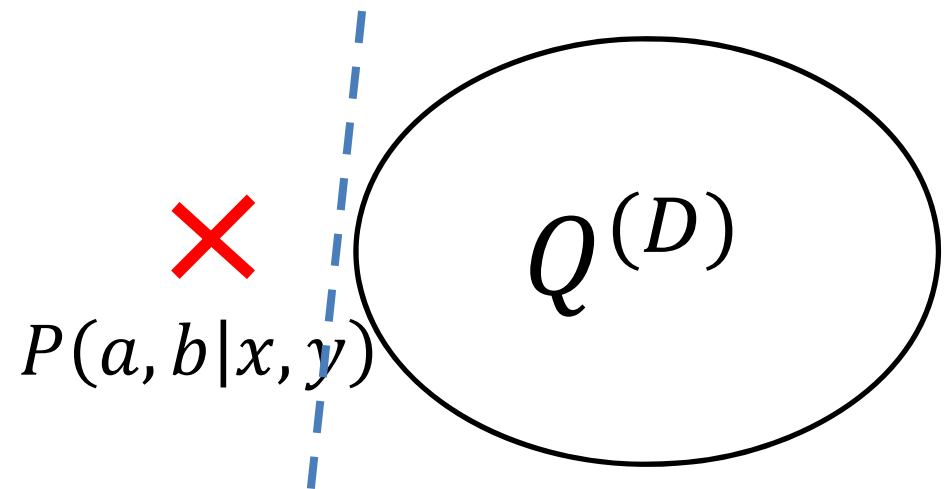


"Easy"

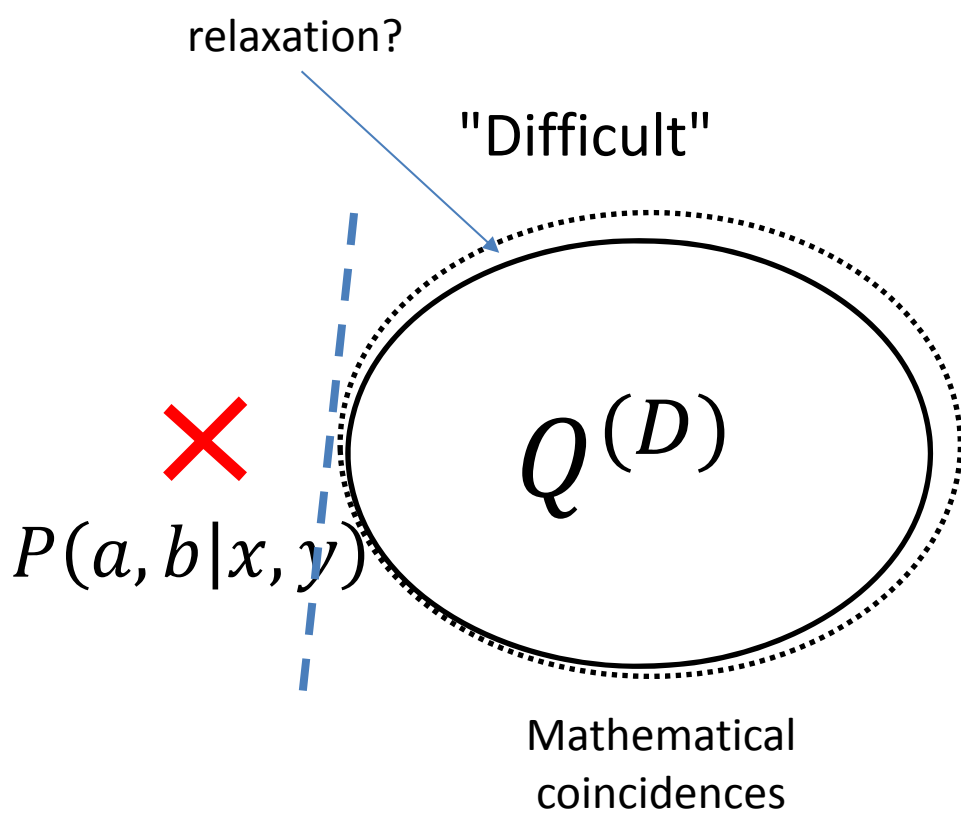
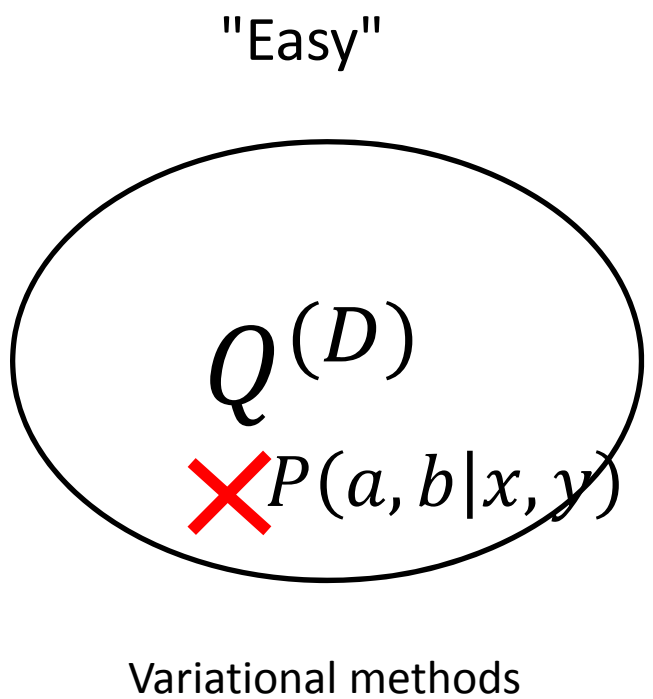


Variational methods

"Difficult"



Mathematical  
coincidences



# Quantum communication complexity

Alice

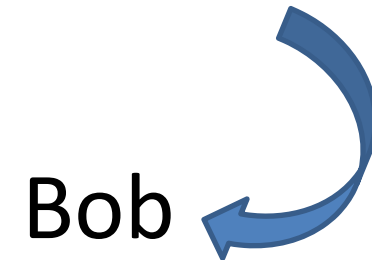
Bob

## Quantum communication complexity

$$\bar{x} \in \{0,1\}^n$$



$$\bar{y} \in \{0,1\}^n$$



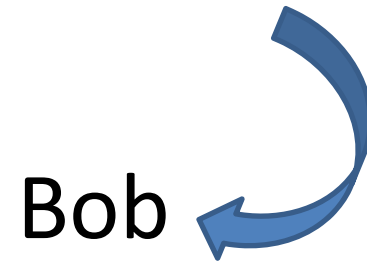
## Quantum communication complexity

$$\bar{x} \in \{0,1\}^n$$



Alice

$$\bar{y} \in \{0,1\}^n$$

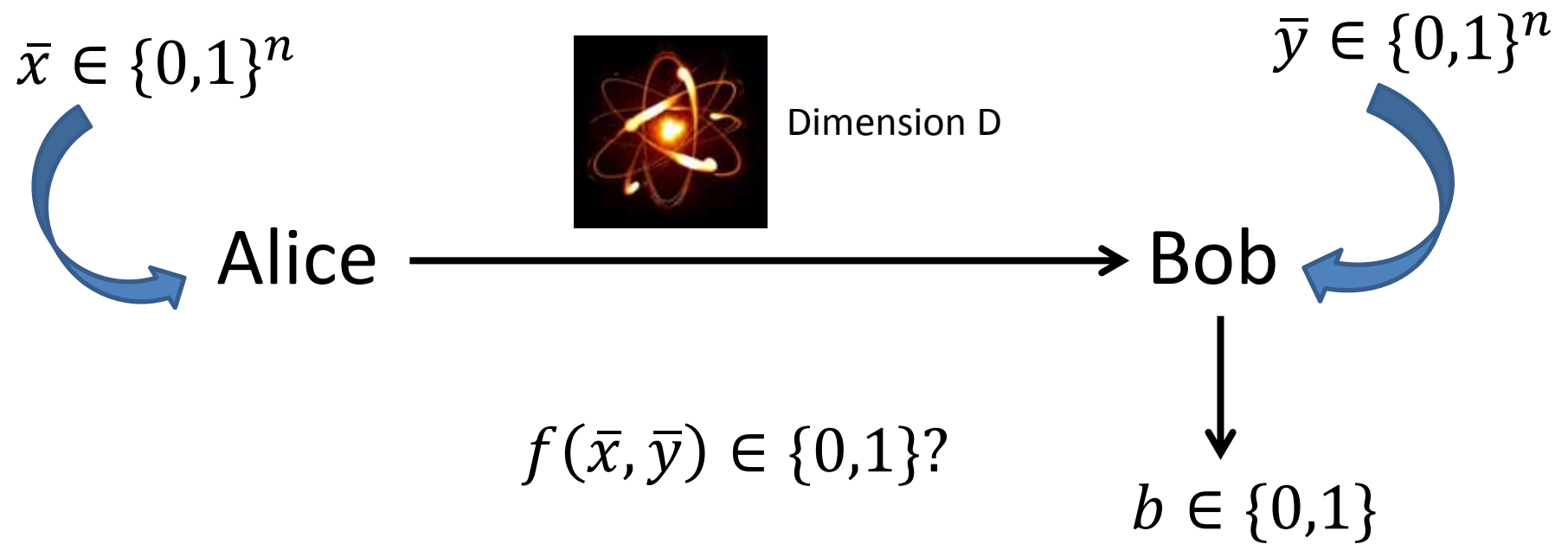


Bob

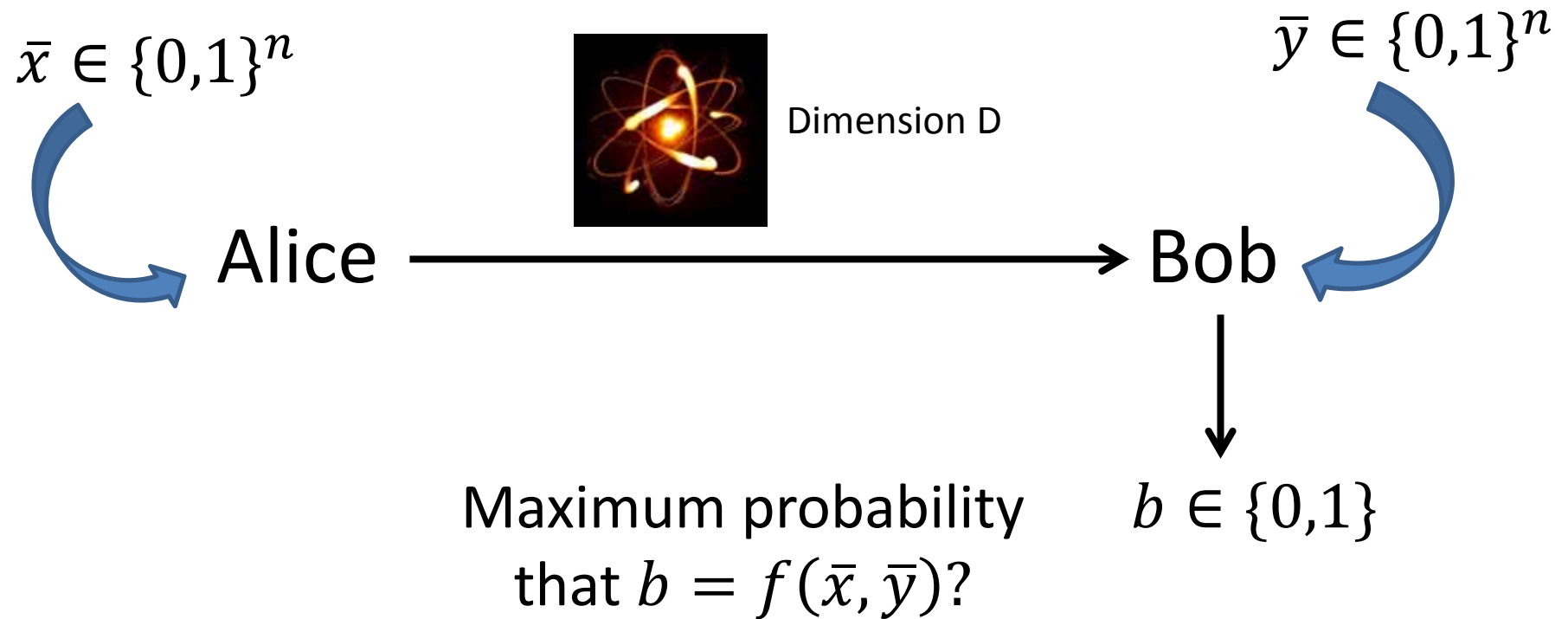
$$f(\bar{x}, \bar{y}) \in \{0,1\}?$$



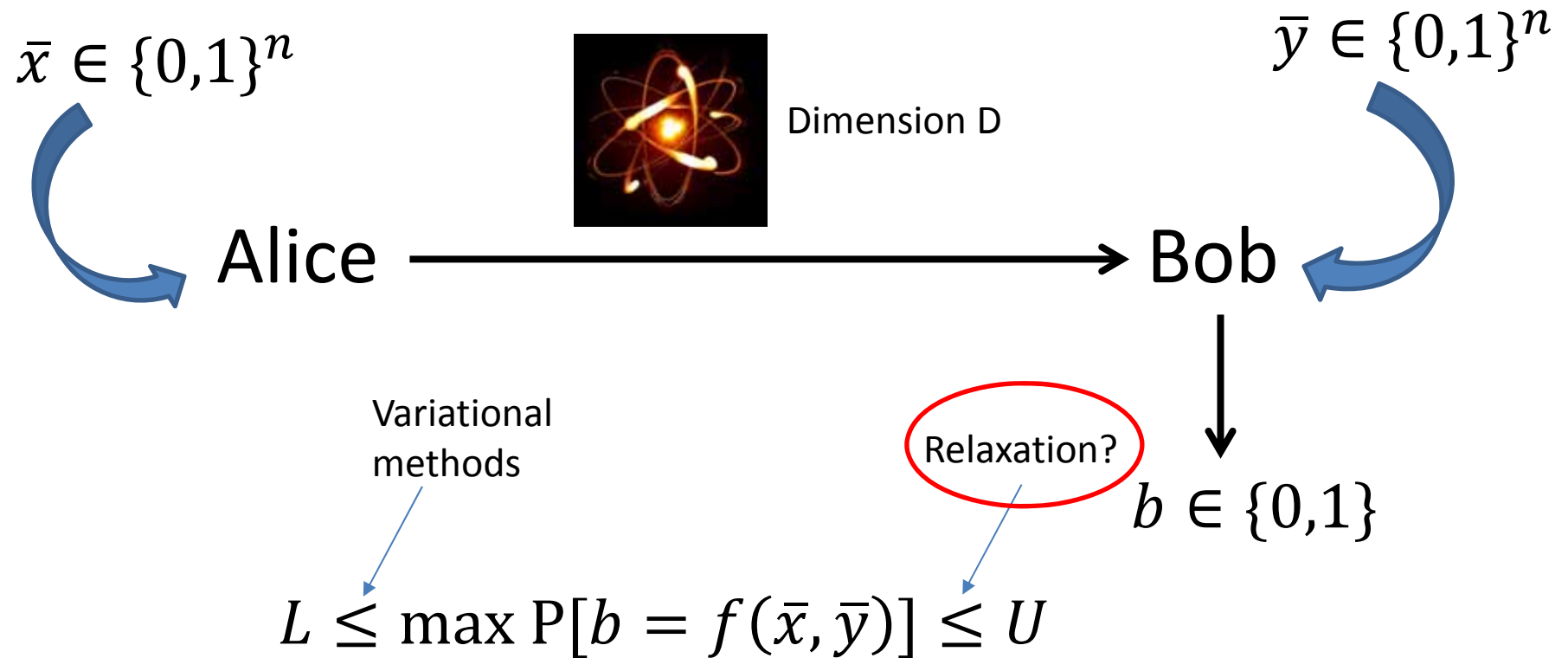
## Quantum communication complexity



## Quantum communication complexity



## Quantum communication complexity



$$p^* = \frac{1}{2^{2n}} \max \sum_{x,y} \text{tr}(F_{f(x,y)}^y \rho_x)$$

$$\text{s.t. } F_b^y, \rho_x \in B(\mathbb{C}^D),$$

$$(F_b^y)^2 = F_b^y, \\ \rho_x \geq 0, \text{tr}(\rho_x) = 1$$

$$\max \left\langle \phi, \sum_{a,b,x,y} C_{a,b}^{x,y} E_a^x \otimes F_b^y \phi \right\rangle$$

$$\text{s.t. } \phi \in \mathcal{H}^{\otimes 2}, \langle \phi, \phi \rangle = 1$$

$$E_a^x, F_b^y \in B(\mathcal{H}), \text{ projectors,}$$

$$\sum_a E_a^x = \sum_b F_b^y = 1,$$

$$\dim(\mathcal{H}) \leq D$$

Brute force theoretically possible, but impractical

What is a  $D$ -dimensional Hilbert space?

$$\psi = (\underbrace{\psi_1, \dots, \psi_D}_{D=\# \text{ of entries}})$$

$$X_i = \begin{pmatrix} x_{11}^i & \dots & x_{1D}^i \\ \vdots & \ddots & \vdots \\ x_{D1}^i & \dots & x_{DD}^i \end{pmatrix} \quad D=\# \text{ of columns}$$

Moment matrix

$$\Gamma = \begin{array}{c|cccc}
 & \mathbb{I} & E_a^x & E_{a'}^{x'} & F_b^y \dots \\
 \hline
 \mathbb{I}^\dagger & 1 & & & \\
 E_a^{x\dagger} & & \langle E_a^x E_{a'}^{x'} \rangle & P(a, b|x, y) & \\
 E_{a'}^{x'\dagger} & & P(a'|x') & P(a', b|x', y) & \\
 F_b^{y\dagger} & & & & \\
 \vdots & & & & 
 \end{array} \succeq 0$$

$\Gamma_{u,v} = \langle u^\dagger v \rangle$

How to incorporate dimension constraints?

Matrix polynomial identities

$$F(X) = 0, \forall X_1, \dots, X_n \in B(\mathbb{C}^d), d \leq D$$



## Matrix polynomial identities

$$F(X) = 0, \forall X_1, \dots, X_n \in B(\mathbb{C}^d), d \leq D$$

$$\text{E.g.: } \left\{ \begin{array}{l} D = 1 \longrightarrow [X_1, X_2] \end{array} \right.$$

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E.g.:

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Central polynomial (commutes with everything)

Standard Identity

$X_1, \dots, X_{2D}, D \times D$  matrices



$$F_D(X) = \sum_{\pi \in S_{2D}} \operatorname{sgn}(\pi) X_{\pi(1)} \cdots X_{\pi(2D)} = 0$$

Moment matrix

$$\Gamma = \begin{array}{c|cccc} & \mathbb{I} & E_a^x & E_{a'}^{x'} & F_b^y \dots \\ \hline \mathbb{I}^\dagger & 1 & & & \\ E_a^{x\dagger} & & \langle E_a^x E_{a'}^{x'} \rangle & P(a, b|x, y) & \\ E_{a'}^{x'\dagger} & P(a'|x') & & P(a', b|x', y) & \\ F_b^{y\dagger} & & & & \\ \vdots & & & & \end{array} \succeq 0$$

Moment matrix

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$$F(X) = 0, \forall X_1, \dots, X_n \in B(\mathbb{C}^d), d \leq D \quad \Rightarrow \quad \langle F(E_a^x, F_b^y) \rangle = 0 \quad \Rightarrow \quad \sum_{i,j} C_{i,j}^F \Gamma_{i,j} = 0$$

Moment matrix

$$\Gamma = \begin{array}{c|cccc} & \mathbb{I} & E_a^x & E_{a'}^{x'} & F_b^y \dots \\ \hline \mathbb{I}^\dagger & 1 & & & \\ E_a^{x\dagger} & & \langle E_a^x E_{a'}^{x'} \rangle & P(a, b|x, y) & \\ E_{a'}^{x'\dagger} & & P(a'|x') & P(a', b|x', y) & \\ F_b^{y\dagger} & & & & \\ \vdots & & & & \end{array} \succeq 0$$



Moment matrices of D-level quantum systems are subject to non-trivial linear constraints!!!

D-dimensional  
Hilbert space



$$\Gamma \geq 0$$

$$\Gamma \in \mathcal{S}^D$$



D-dimensional  
Hilbert space



$$\Gamma \geq 0$$

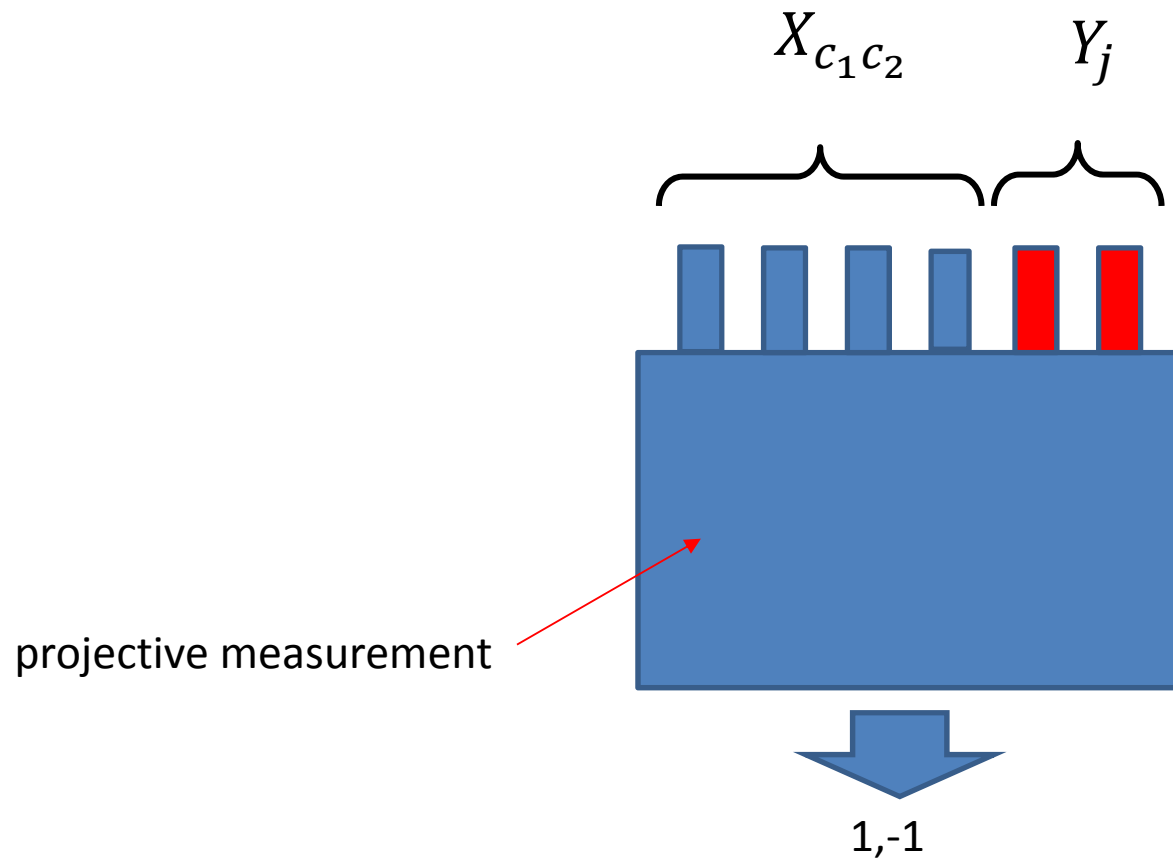
$$\Gamma \in \mathcal{S}^D ?$$

Toy problem

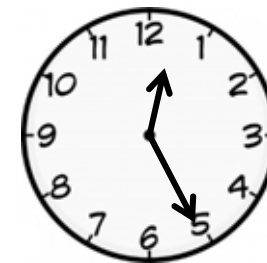
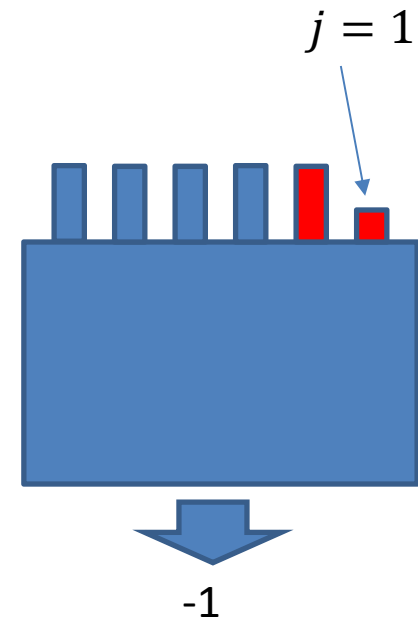
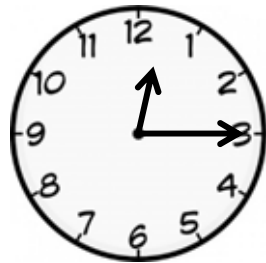
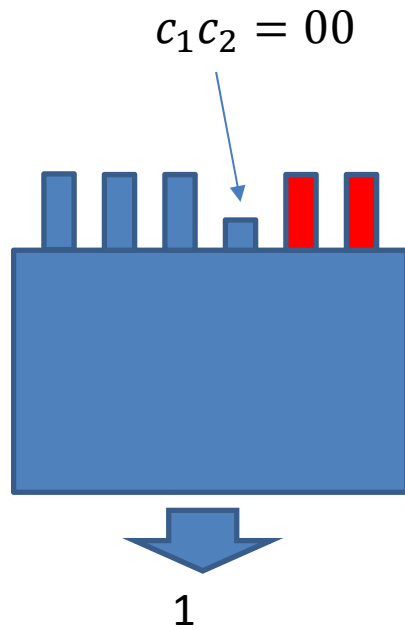
$$\max \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} \langle X_{c_1 c_2} Y_j X_{c_1 c_2} \rangle$$

$$\text{s.t. } X_{c_1 c_2}^2 = Y_j^2 = 1$$

## Temporal quantum correlations



## Temporal quantum correlations



Toy problem

$$p^* = \max \langle \phi | \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} X_{c_1 c_2} Y_j X_{c_1 c_2} | \phi \rangle$$

$$\text{s.t. } X_{c_1 c_2}^2 = Y_j^2 = 1,$$

$$\langle \phi | \phi \rangle = 1$$

Toy problem

$$p^* = \max \langle \phi | \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} X_{c_1 c_2} Y_j X_{c_1 c_2} | \phi \rangle$$

$$\text{s.t. } X_{c_1 c_2}^2 = Y_j^2 = 1,$$

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In the dimension-free case, this kind of problems can be solved with a single SDP

S. Pironio, M. Navascués and A. Acín, SIAM J. Optim. 20, 5, 2157-2180 (2010).

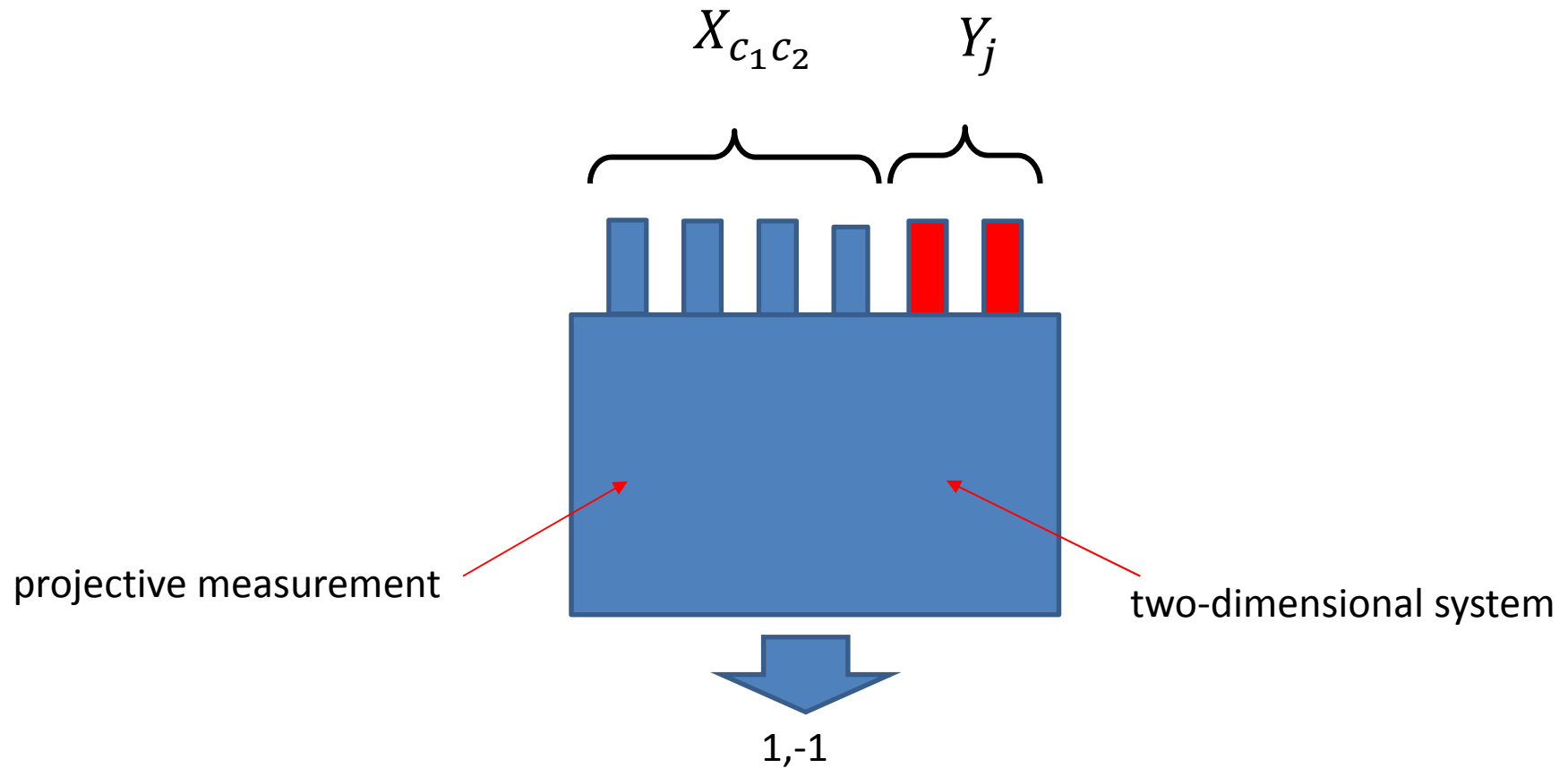
C. Budroni, T. Moroder, M. Kleinmann, O. Gühne, Phys. Rev. Lett. 111, 020403 (2013)



$$p^* = 8$$



# Temporal quantum correlations



Toy problem

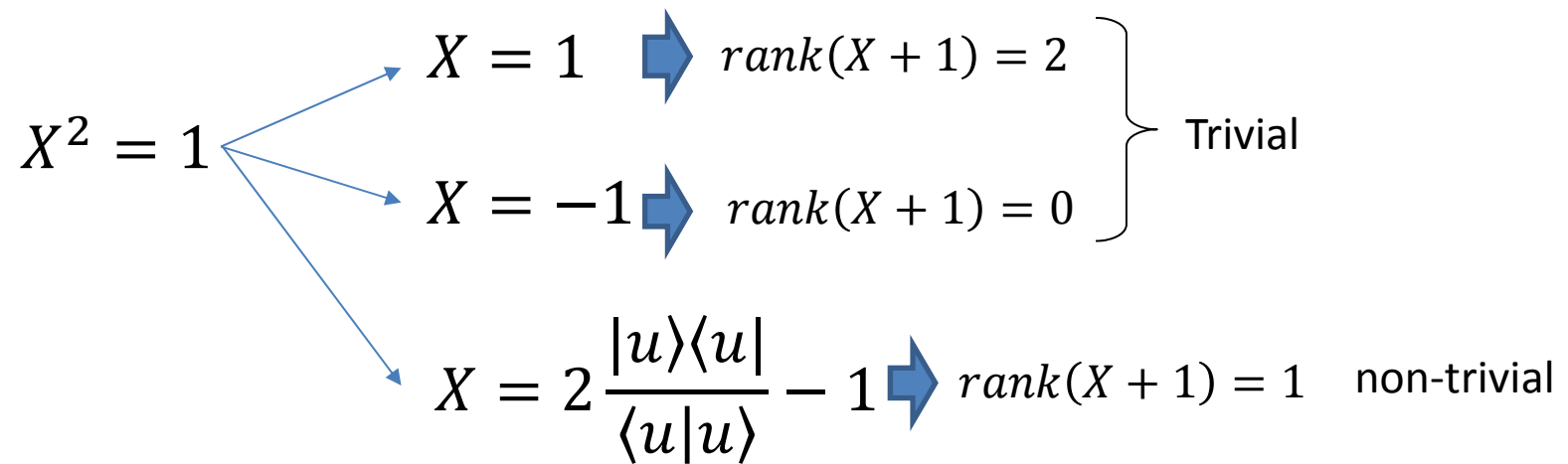
$$p^* = \max \langle \phi | \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} X_{c_1 c_2} Y_j X_{c_1 c_2} | \phi \rangle$$

$$\text{s.t. } X_{c_1 c_2}^2 = Y_j^2 = 1,$$

$$\langle \phi | \phi \rangle = 1,$$

$$|\phi\rangle \in \mathbb{C}^2, X_{c_1 c_2}, Y_j \in B(\mathbb{C}^2)$$

Solving the toy problem (I): divide the problem into classes



A random instance can be generated easily

Solving the toy problem (I): divide the problem into classes

729 classes

Each class labeled by a  
vector  $\vec{r} \in \{0,1,2\}^6$



$$\mathit{rank}(X_{c_1 c_2} + 1) = r_{c_1 c_2}$$

$$\mathit{rank}(Y_j + 1) = r_j$$

Classed toy problem

$$p_{\vec{r}}^* = \max \langle \phi | \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} X_{c_1 c_2} Y_j X_{c_1 c_2} | \phi \rangle$$

$$\text{s.t. } X_{c_1 c_2}^2 = Y_j^2 = 1,$$

$$\langle \phi | \phi \rangle = 1,$$

$$|\phi\rangle \in \mathbb{C}^2, X_{c_1 c_2}, Y_j \in B(\mathbb{C}^2),$$

$$\text{rank}(X_{c_1 c_2} + 1) = r_{c_1 c_2},$$

$$\text{rank}(Y_j + 1) = r_j$$

$$\max \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} \Gamma_{X_{c_1 c_2} Y_j X_{c_1 c_2}}$$

$$\text{s.t. } \Gamma_{1,1} = 1$$

$$\Gamma \geq 0,$$

$$\Gamma \in S_{D, \vec{r}}$$

Solving the toy problem (II): Identify  $S_{D,\vec{r}}$

$$\max \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} \Gamma_{X_{c_1 c_2} Y_j X_{c_1 c_2}}$$

$$\text{s.t. } \Gamma_{1,1} = 1$$

$$\Gamma \geq 0,$$

$$\Gamma \in S_{D,\vec{r}}$$

High level description

$$i = 1$$



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Generate random dichotomic operators  $X_{c_1 c_2}^i, Y_j^i \in B(\mathbb{C}^2)$  (with app. ranks) and vector  $\phi^i \in \mathbb{C}^2$

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Build moment matrix  $\Gamma^j$

$$\begin{aligned}\Gamma_{\mathbb{I}, \mathbb{I}}^i &= \langle \phi^i | \mathbb{I}_d | \phi^i \rangle = 1, \\ \Gamma_{X_{c_1 c_2}, Y_j}^i &= \langle \phi^i | X_{c_1 c_2}^i, Y_j^i | \phi^i \rangle, \\ &\vdots\end{aligned}$$

# High level description

$$i = 1$$

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Build moment matrix  $\Gamma^j$

$$\Gamma_{\mathbb{I}, \mathbb{I}}^i = \langle \phi^i | \mathbb{I}_d | \phi^i \rangle = 1,$$
$$\Gamma_{X_{c_1 c_2}, Y_j}^i = \langle \phi^i | X_{c_1 c_2}^i, Y_j^i | \phi^i \rangle,$$

⋮



$$i = i + 1$$

Repeat until the moment matrix  $\Gamma^{(N+1)}$  is a linear combination of  $\Gamma^{(1)}, \dots, \Gamma^{(N)}$

# High level description

$$i = 1$$

Generate random dichotomic operators  $X_{c_1 c_2}^i, Y_j^i \in B(\mathbb{C}^2)$  (with app. ranks) and vector  $\phi^i \in \mathbb{C}^2$



Build moment matrix  $\Gamma^j$

$$\Gamma_{\mathbb{I}, \mathbb{I}}^i = \langle \phi^i | \mathbb{I}_d | \phi^i \rangle = 1,$$
$$\Gamma_{X_{c_1 c_2}, Y_j}^i = \langle \phi^i | X_{c_1 c_2}^i, Y_j^i | \phi^i \rangle,$$

⋮



$$i = i + 1$$

At that point, *any* feasible moment matrix  $\Gamma$  must be a linear combination of  $\Gamma^{(1)}, \dots, \Gamma^{(N)}$



$$\mathcal{S}_r^D = \text{span}(\Gamma^{(1)}, \dots, \Gamma^{(N)})$$



SDP relaxation

$$\max \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} \Gamma_{X_{c_1 c_2} Y_j X_{c_1 c_2}}$$

$$\text{s.t. } \Gamma_{1,1} = 1$$

$$\Gamma \geq 0,$$

$$\Gamma = \sum_{s=1}^N c_s \Gamma^s$$

SDP relaxation

$$\max \sum_{j=1,2} \sum_{c_1, c_2=0,1} (-1)^{c_j} \Gamma_{X_{c_1 c_2} Y_j X_{c_1 c_2}}$$

s.t.  $\Gamma_{1,1} = 1$

$\Gamma \geq 0,$  Lack of strict feasibility,  
i.e., there are no  $\Gamma > 0$

Problems with SDP solvers



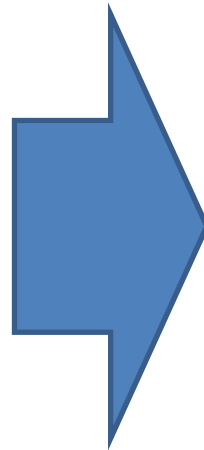
$$\Gamma = \sum_{s=1}^N c_s \Gamma^s \rightarrow |c_s| \gg 1$$

## Implementation tips



$$\frac{1}{N} \sum_{s=1}^N \Gamma^s = \Gamma_{mix} \geq 0$$

Isometry  $\bar{V}$  to the  
support of  $\Gamma_{mix}$



$$\Gamma \geq 0 \quad \Rightarrow \quad \bar{V}^\dagger \Gamma \bar{V} \geq 0$$

Why?

$$\bar{V}^\dagger \Gamma_{mix} \bar{V} > 0$$

Strict positivity

## Implementation tips



Use (modified) Gram-Schmidt as you generate the matrices  $\Gamma^1, \Gamma^2 \dots$

Sequence of orthogonal matrices  $\tilde{\Gamma}^1, \tilde{\Gamma}^2 \dots$

at  $s=N+1, \tilde{\Gamma}^s = 0$ , up to computer precision

$$\Gamma = \sum_{s=1}^N c_s \Gamma^s \quad \Rightarrow \quad \Gamma = \sum_{s=1}^N \tilde{c}_s \frac{\tilde{\Gamma}_s}{\sqrt{\text{tr}(\tilde{\Gamma}_s^2)}}$$



SDP relaxation

$$\begin{aligned} \min \quad & \sum_{a,b,x,y} C_{a,b}^{x,y} \Gamma_{E_a^x, F_b^y} \\ \text{s.t.} \quad & \Gamma_{1,1} = 1, \\ & \Gamma \geq 0, \\ & \Gamma = \sum_{s=1}^N c_s \Gamma^s \end{aligned}$$



$$\begin{aligned} \min \quad & \sum_{a,b,x,y} C_{a,b}^{x,y} \Gamma_{E_a^x, F_b^y} \\ \text{s.t.} \quad & \Gamma_{1,1} = 1, \\ & \bar{V}^\dagger \Gamma \bar{V} \geq 0, \\ & \Gamma = \sum_{s=1}^N \tilde{c}_s \tilde{\Gamma}^s \end{aligned}$$



Free dimensionality

$$p^* = 8$$

Second relaxation, D=2

$$p^2 = 5.656854 \dots$$

Generalization

NPO problem with dimension constraints

$$p^* = \min_{\mathcal{H}, X, \phi} \langle \phi, p(X)\phi \rangle$$

$$\text{s.t. } X_1, \dots, X_n \in B(\mathcal{H}),$$
$$q_i(X) \geq 0, i = 1, \dots, s$$

$$\phi \in \mathcal{H}$$

$$\langle \phi, \phi \rangle = 1$$

$$\dim(\mathcal{H}) \leq D$$

$$\begin{aligned} p_k^* &= \min \sum_w p_w y_w \\ \text{s.t. } & y_1 = 1, \\ & M_k(y) \geq 0, \\ & M_{k-\lfloor d_i/2 \rfloor}(q_i y) \geq 0, \\ & y = \sum_{s=1}^{j-1} c_s y^{(s)}, c_s \in \mathbb{C}. \end{aligned}$$



$$p_k^* \leq p^*$$

$$\begin{aligned}
p_k^* &= \min \sum_w p_w y_w \\
\text{s.t. } & y_1 = 1, \\
& M_k(y) \geq 0, \\
& M_{k-\lfloor d_i/2 \rfloor}(q_i y) \geq 0, \\
& y = \sum_{s=1}^{j-1} c_s y^{(s)}, c_s \in \mathbb{C}.
\end{aligned}$$

Archimedean  
condition

$$C - \sum_r X_r^\dagger X_r + X_r X_r^\dagger = \sum_s f_s f_s^\dagger + \sum_{s,i} g_{s,i} q_i g_{s,i}^\dagger$$

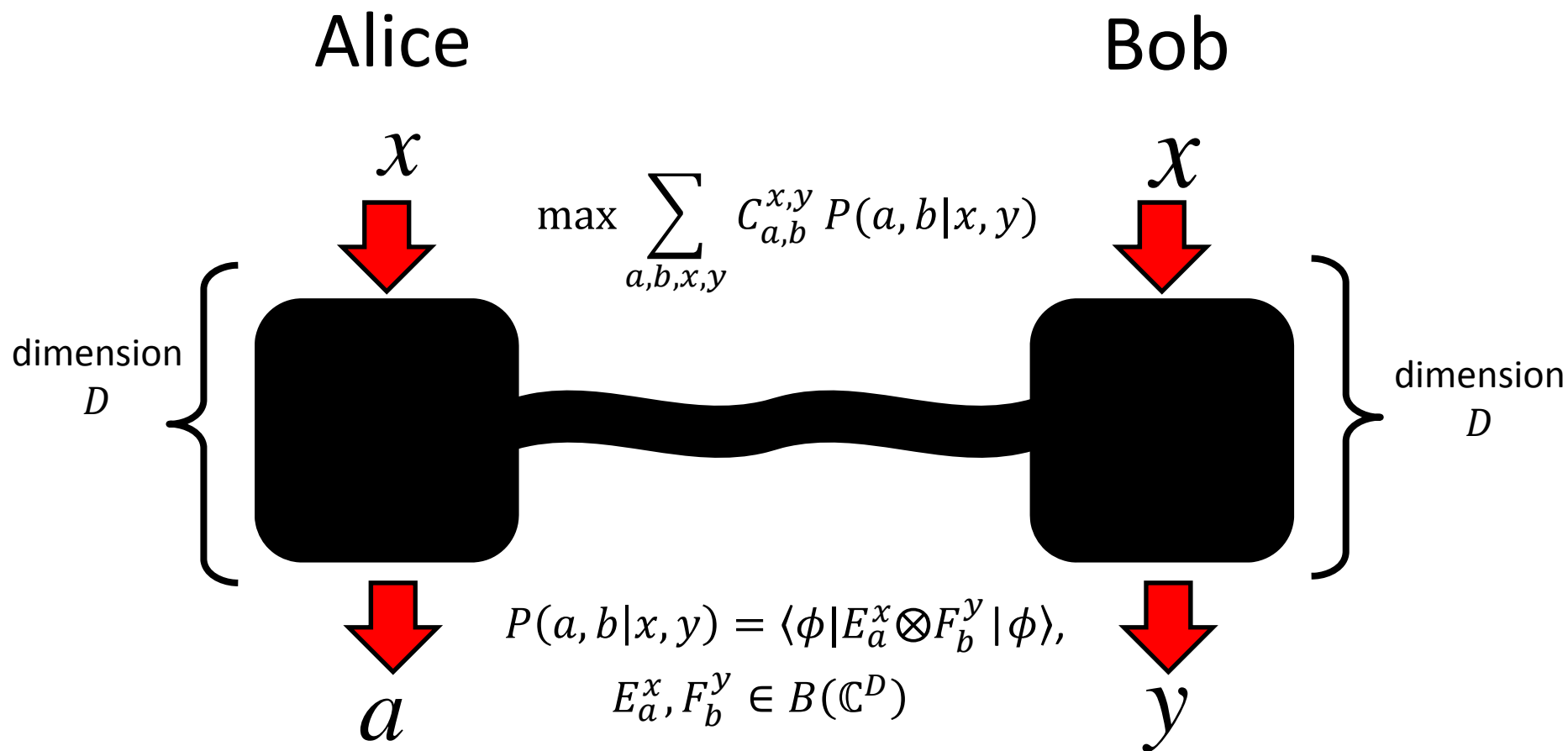


$$\lim_{k \rightarrow \infty} p_k^* = p^*$$

Related hierarchies



SDP hierarchy for quantum non-locality under dimension constraints



$$\max \left\langle \phi, \sum_{a,b,x,y} C_{a,b}^{x,y} E_a^x \otimes F_b^y \phi \right\rangle$$

$$\text{s.t. } \phi \in \mathcal{H}^{\otimes 2}, \langle \phi, \phi \rangle = 1$$

$$E_a^x, F_b^y \in \mathcal{B}(\mathcal{H}), \text{ projectors,}$$

$$\sum_a E_a^x = \sum_b F_b^y = 1,$$

$$\dim(\mathcal{H}) \leq D$$

$$\max \left\langle \phi, \sum_{a,b,x,y} C_{a,b}^{x,y} E_a^x \otimes F_b^y \phi \right\rangle$$

$$\text{s.t. } \phi \in \mathcal{H}^{\otimes 2}, \langle \phi, \phi \rangle = 1$$

$$E_a^x, F_b^y \in \mathcal{B}(\mathcal{H}), \text{ projectors,}$$

$$\sum_a E_a^x = \sum_b F_b^y = 1,$$

$$\dim(\mathcal{H}) \leq D,$$

$$\text{rank}(E_a^x) = r_a^x, \text{rank}(F_b^y) = t_b^y$$

# High level description

$$i = 1$$

Generate random projectors

$E_a^{i,x}, F_b^{i,y} \in B(\mathbb{C}^D)$  (with app. ranks)  
and vector  $\phi^i \in \mathbb{C}^{D^2}$



Build moment matrix  $\Gamma^i$

$$\Gamma_{\mathbb{I},\mathbb{I}}^i = \langle \phi^i | \mathbb{I}_d \otimes \mathbb{I}_d | \phi^i \rangle = 1,$$

$$\Gamma_{E_a^x, F_b^y}^i = \langle \phi^i | E_a^{i,x} \otimes F_b^{i,y} | \phi^i \rangle,$$

⋮



$$i = i + 1$$

Repeat until the  
moment matrix  $\Gamma^{(N+1)}$  is  
a linear combination of  
 $\Gamma^{(1)}, \dots, \Gamma^{(N)}$

SDP hierarchy

$$\max \sum_{a,b,x,y} C_{a,b}^{x,y} \Gamma_{E_a^x, F_b^y}$$

$$\text{s.t. } \Gamma_{1,1} = 1$$

$$\Gamma \geq 0, \Gamma = \sum_{s=1}^j c_s \Gamma^s$$

Convergence is guaranteed



```
' Sends text and returns with result including the  
' '>' Time out is 100 ms  
Public Function SendText(ByVal text)  
Dim timeNow  
With Form1  
SendText = ""  
.textIn = "" ' clear buffer  
.Sout(text) ' send it  
timeNow = .t1  
While .t1 < timeNow + 10  
If InStr(.textIn, ">") > 0 Then  
SendText = .textIn  
Exit While  
End If  
Application.DoEvents()  
End While  
End With  
End Function
```

```
MS-DOS Prompt - BBASIC  
FILE EDIT VIEW EXEC Run Debug Options Help  
ROBOT.BAS  
...  
Public Interface Drive routines  
Dim Parameters Computer Technology 1995  
GetStart 2.4  
...  
DIM SHARED Fischer*(0 To 511)  
DEF SEG = VARSEG(Fischer*(0))  
DEF SEG = Fischer*(512), UARPTR(Fischer*(0))  
FischerInstall  
HandOpen 38  
Ready 38  
...  
HandOpen 38  
Ready 38  
...  
F1=F1=Help2 <F8=Win=Out> <F2=Sub> <F3=Run> <F8=Step> 0064400
```

```
q.Vertices(descendant_index).state_in = q.Vertices(our_state_index).state_in  
q.Vertices(descendant_index).state_out = q.Vertices(our_state_index).state_out  
...  
q.Vertices(descendant_index).state = "i" if sev_gs < q.Vertices(descendant_index).state  
q.Vertices(descendant_index).state = "o" if sev_gs > q.Vertices(descendant_index).state  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vertices(descendant_index).path  
q.Vertices(descendant_index).gs = sev_gs // Close  
...  
q.Vertices(descendant_index).state = "i" if sev_gs < q.Vertices(descendant_index).state  
q.Vertices(descendant_index).state = "o" if sev_gs > q.Vertices(descendant_index).state  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vertices(descendant_index).path  
q.Vertices(descendant_index).gs = sev_gs // Close  
...  
q.Vertices(descendant_index).fn = Calculate(q.Vertices(descendant_index).qsall)
```

NPA

Pironio-Bell  $\leq 0.2532$

Pironio-Bell  $\leq 0.2071$

D=2

NPA

$I_{3322} \leq 0.250875$

$I_{3322} \leq 0.25$

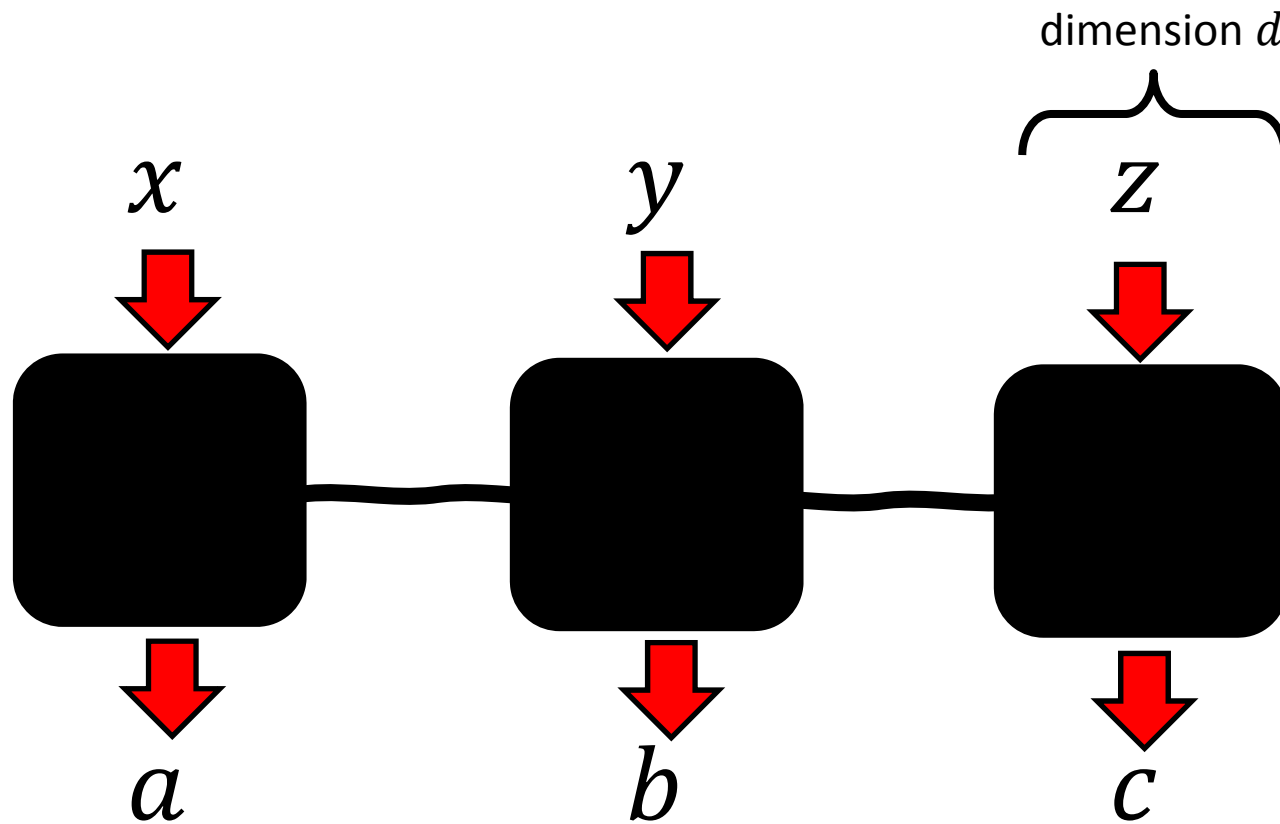
D=2, 3

$$I_{4,4} = E_1^A + E_{1,1} + E_{1,2} + E_{2,1} - E_{2,2} + E_{3,3} + E_{3,4} + E_{4,3} - E_{4,4} \leq 5.9907$$

quantum mechanics

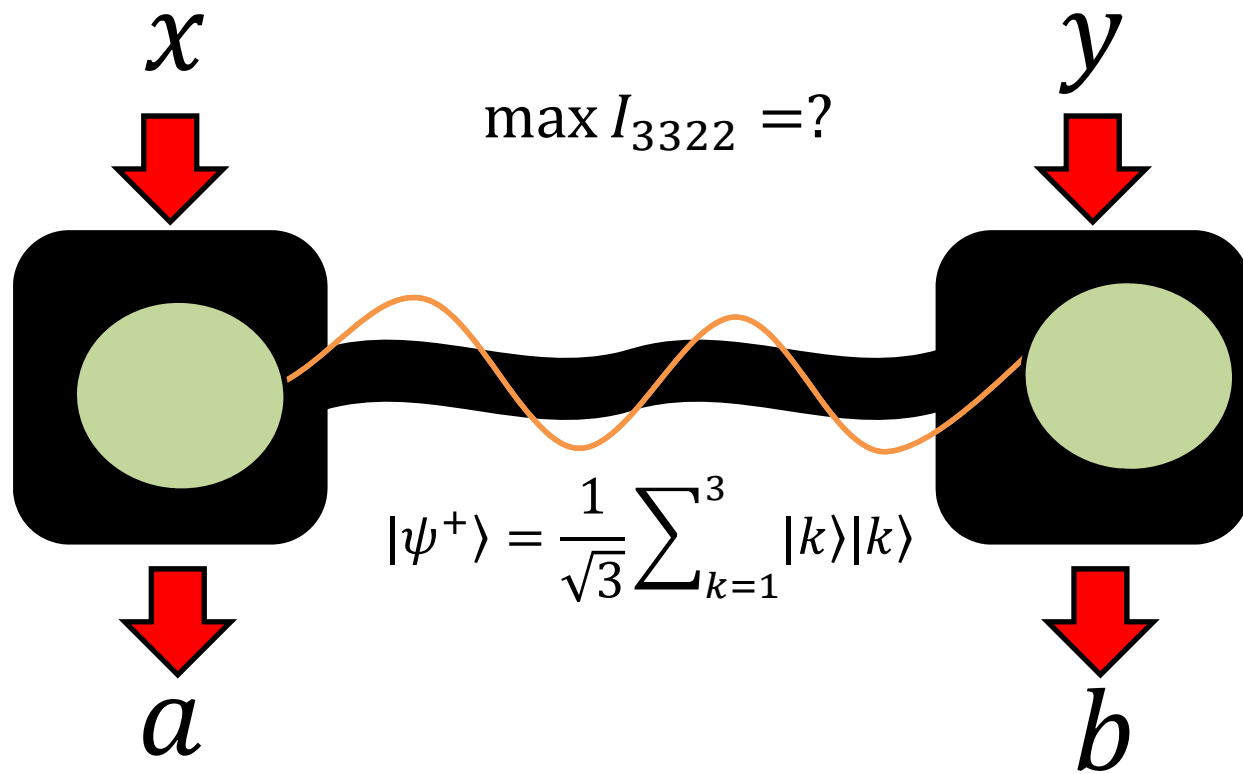
$$I_{4,4} = E_1^A + E_{1,1} + E_{1,2} + E_{2,1} - E_{2,2} + E_{3,3} + E_{3,4} + E_{4,3} - E_{4,4} \leq 5.8310$$

D=2, 3



$$P(a, b, c | x, y, z)$$





# High level description

$$i = 1$$

Generate random projectors

$$E_a^{i,x}, F_b^{i,y} \in B(\mathbb{C}^D) \text{ (with app. ranks)}$$



Build moment matrix  $\Gamma^i$

$$\Gamma_{\mathbb{I},\mathbb{I}}^i = \langle \psi^+ | \mathbb{I}_d \otimes \mathbb{I}_d | \psi^+ \rangle = 1,$$

$$\Gamma_{E_a^x, F_b^y}^i = \langle \psi^+ | E_a^{i,x} \otimes F_b^{i,y} | \psi^+ \rangle,$$

⋮



$$i = i + 1$$

Repeat until the moment matrix  $\Gamma^{(N+1)}$  is a linear combination of  $\Gamma^{(1)}, \dots, \Gamma^{(N)}$

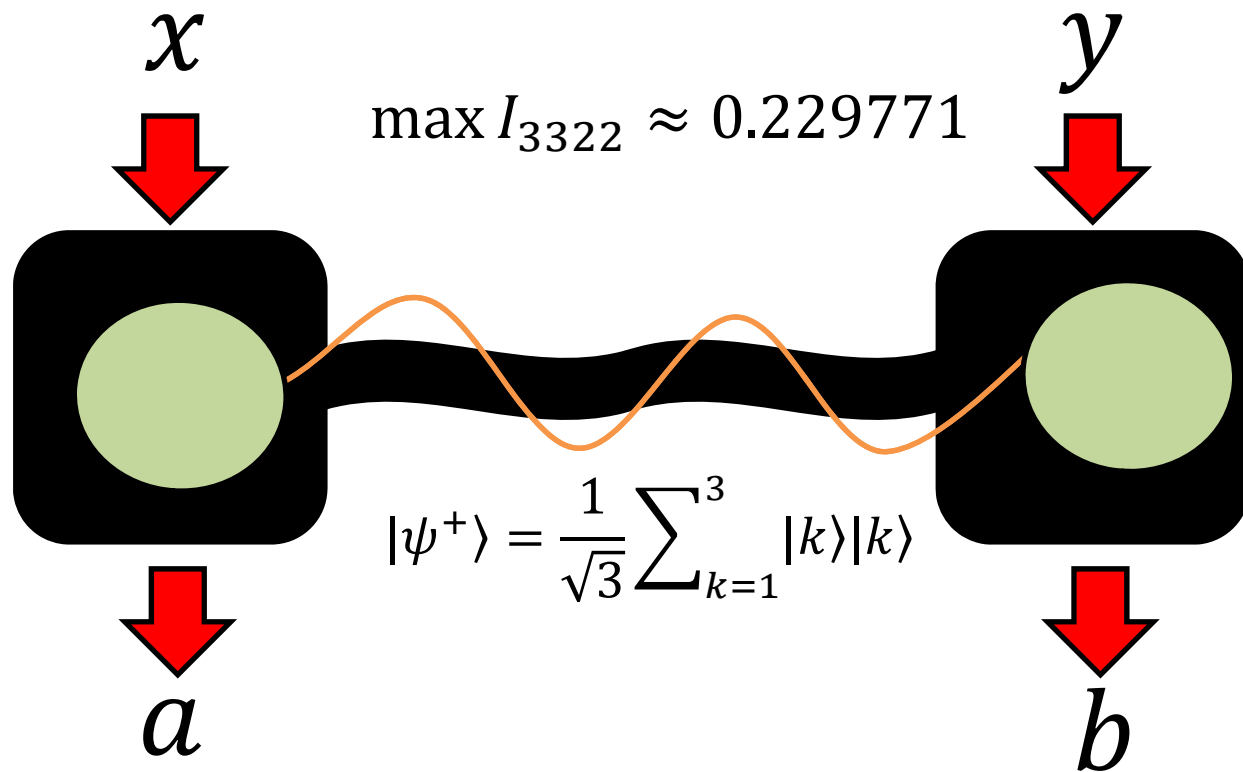




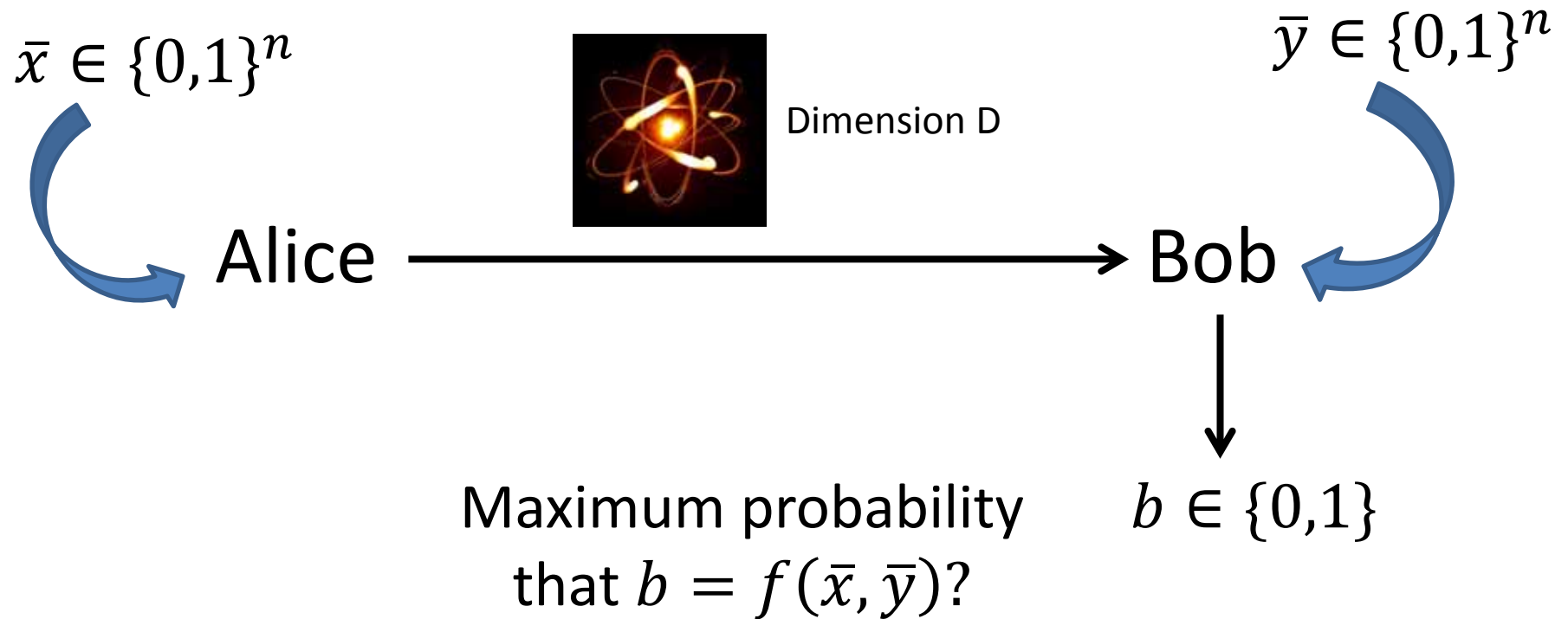
```
' Sends text and returns with result including the
' '>' Time out is 100 ms
Public Function SendText(ByVal text)
Dim timeNow
With Form1
SendText = ""
.textIn = "" ' clear buffer
.Sout(text) ' send it
timeNow = .t1
While .t1 < timeNow + 10
If InStr(.textIn, ">") > 0 Then
SendText = .textIn
Exit While
End If
Application.DoEvents()
End While
End With
End Function
```

```
MS-DOS Prompt - BBASIC
FILE EDIT VIEW EXEC RUN DEBUG OPTIONS
ROBOT.BAS
...
PUBLIC Interface Drive routines
DIM Parameters Computer Technology 1995
GETSTAR 2.4
...
DIM SHARED Fischer*(R TO 511)
DEF SEG = VARSEG(Fischer*(0))
DEF SEG = VARSEG(Fischer*(14), UARPTR(Fischer*(0)))
FischerInstall
HandOpen 38
Ready 38
...
HandOpen 38
Ready 38
...
Immediate
0004:00
```

```
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
...
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
...
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
...
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
q.Vertices(descendant_index).state_index = q.Stateq.Vertices(descendant_index).state_index
```



# SDP hierarchy for quantum communication complexity



$$p^* = \frac{1}{2^{2n}} \max \sum_{x,y} \text{tr}(F_{f(x,y)}^y \rho_x)$$

such that  $F_b^y, \rho_x \in B(\mathbb{C}^D)$ ,

$$(F_b^y)^2 = F_b^y,$$
$$\rho_x \geq 0, \text{tr}(\rho_x) = 1$$

# High level description

$$i = 1$$

Generate random projectors  $\rho_x^i, F_b^{i,y} \in B(\mathbb{C}^D)$  (with app. ranks)



Build moment matrix  $\Gamma^i$

$$\Gamma_{\mathbb{I}, \mathbb{I}}^i = \text{tr}(\mathbb{I}) = D,$$

$$\Gamma_{\rho_x, F_b^y}^i = \text{tr}(\rho_x^i F_b^{i,y}),$$

⋮



$$i = i + 1$$

Repeat until the moment matrix  $\Gamma^{(N+1)}$  is a linear combination of  $\Gamma^{(1)}, \dots, \Gamma^{(N)}$



SDP hierarchy

$$\max \frac{1}{2^{2n}} \sum_{b,x,y} \Gamma_{\rho_x, F_f^y(x,y)}$$

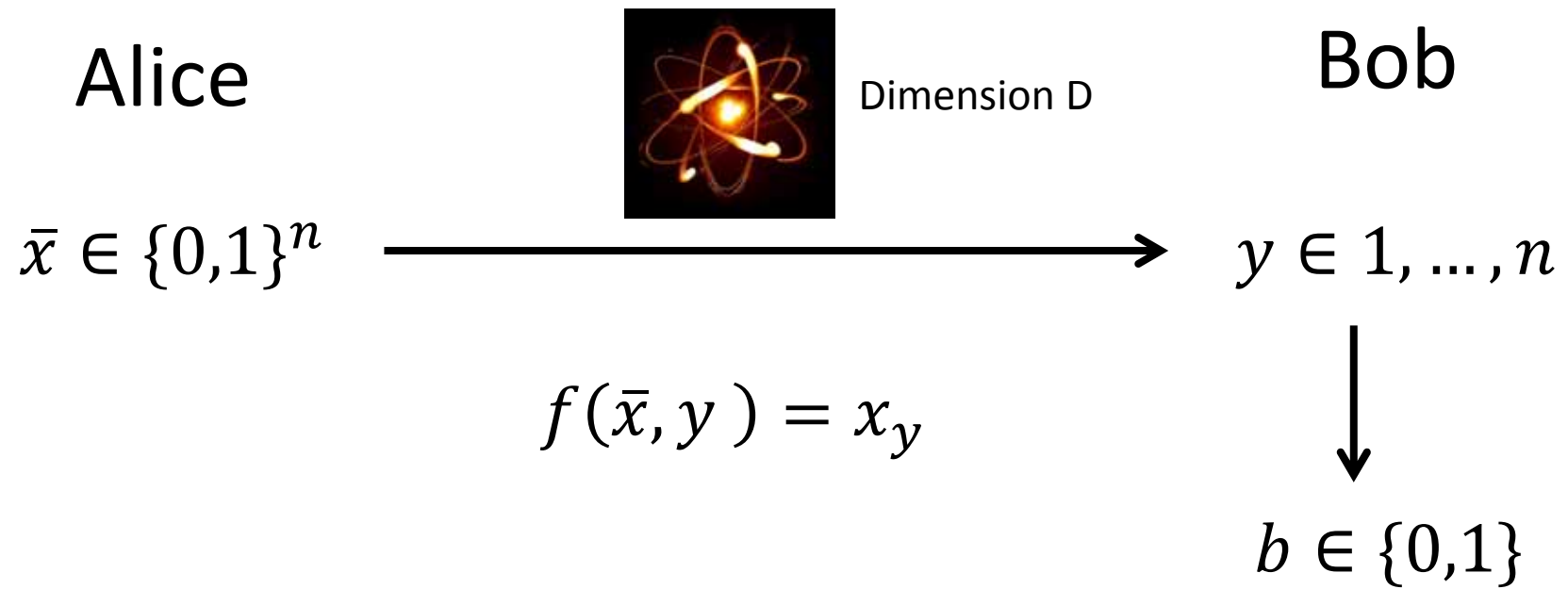
$$\text{s.t. } \Gamma_{1,1} = 1$$

$$\Gamma \geq 0, \Gamma = \sum_{s=1}^j c_s \Gamma^s$$

No proof of convergence!!!



# Random Access Codes (RAC)





```
' Sends text and returns with result including the  
' '>' Time out is 100 ms  
Public Function SendText(ByVal text)  
Dim timeNow  
With Form1  
SendText = ""  
.textIn = "" ' clear buffer  
.Sout(text) ' send it  
timeNow = .t1  
While .t1 < timeNow + 10  
If InStr(.textIn, ">") > 0 Then  
SendText = .textIn  
Exit While  
End If  
Application.DoEvents()  
End While  
End With  
End Function
```

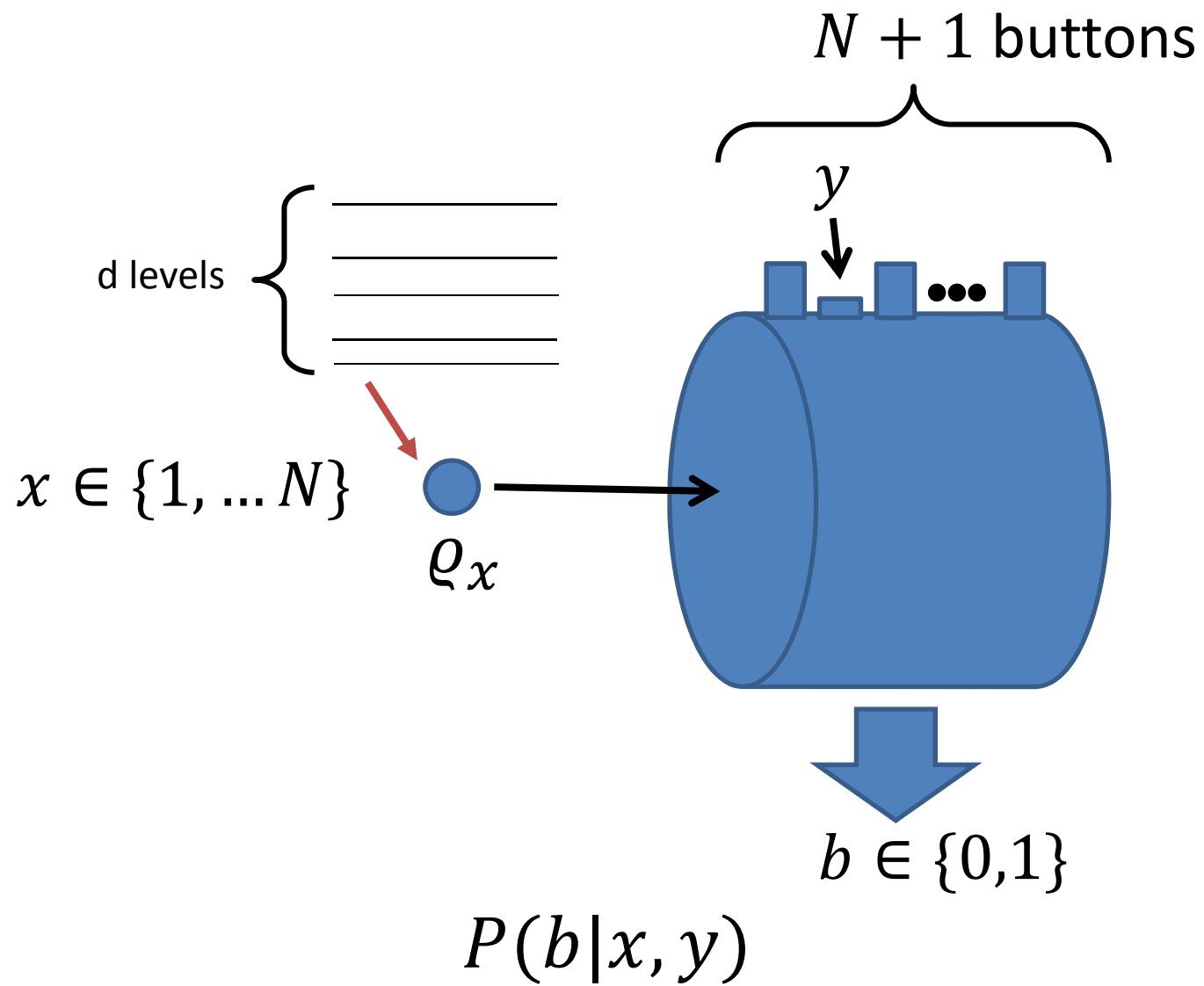
```
MS-DOS Prompt - BBASIC  
FILE EDIT VIEW EXEC RUN DEBUG OPTIONS HELP  
ROBOT.BAS  
...  
Public Interface Drive routines  
Dim Parameters Computer Technology 1995  
GetStart 2.4  
...  
DIM SHARED Fischer*(0 To 511)  
DEF SEG = VARSEG(Fischer*(0))  
DEF SEG = Fischer*(512), UARPTR(Fischer*(0))  
FischerInstall  
HandOpen 38  
Ready 38  
...  
HandOpen 38  
Ready 38  
...  
[F1=Help] [F2=Window] [F3=Sub] [F4=Run] [F5=Step] 00:44:00
```

```
q.Vertices(descendant_index).state_in = q.Vertices(our_state_index).state_in  
q.Vertices(descendant_index).state_out = q.Vertices(our_state_index).state_out  
...  
q.Vertices(descendant_index).state = "1" if sev_gs < q.Vertices(descendant_index).state_in  
q.Vertices(descendant_index).state = "0" if sev_gs < q.Vertices(descendant_index).state_out  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vertices(descendant_index).state  
...  
q.Vertices(descendant_index).state = "1" if sev_gs < q.Vertices(descendant_index).state_in  
q.Vertices(descendant_index).state = "0" if sev_gs < q.Vertices(descendant_index).state_out  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vertices(descendant_index).state  
...  
q.Vertices(descendant_index).fn = Calculate(q.Vertices(descendant_index).state)
```

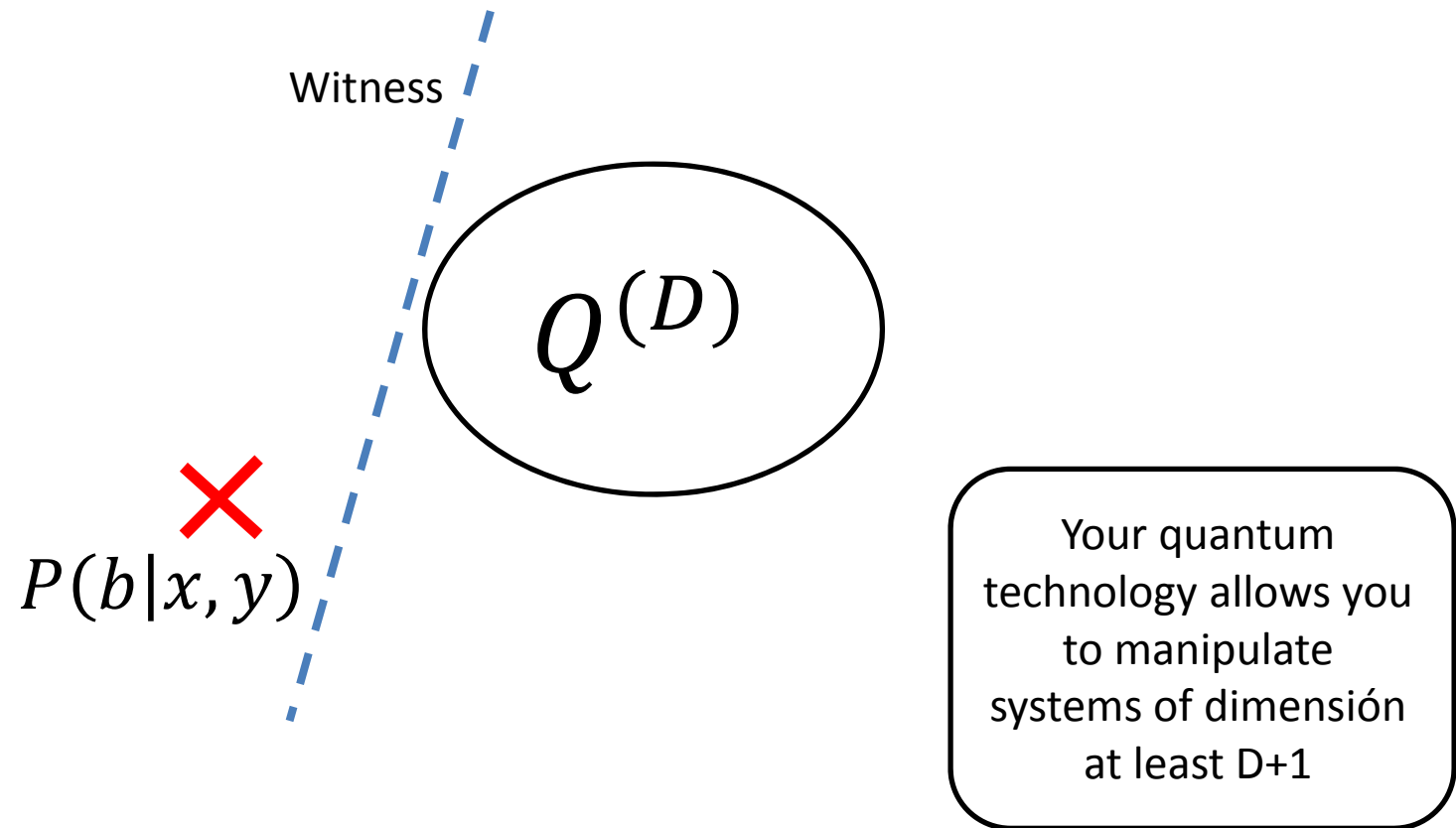
Random Access Codes (RAC)

$D$	2	3	4	5	6	7
LB	0.788675	0.832273	0.908248	0.924431	0.951184	0.969841
UB	0.788675	0.832273	0.908248	0.924445	0.954123	0.969841

TABLE I: Lower and upper bounds on  $P_{\max}(3 \rightarrow \log_2(D))$ .



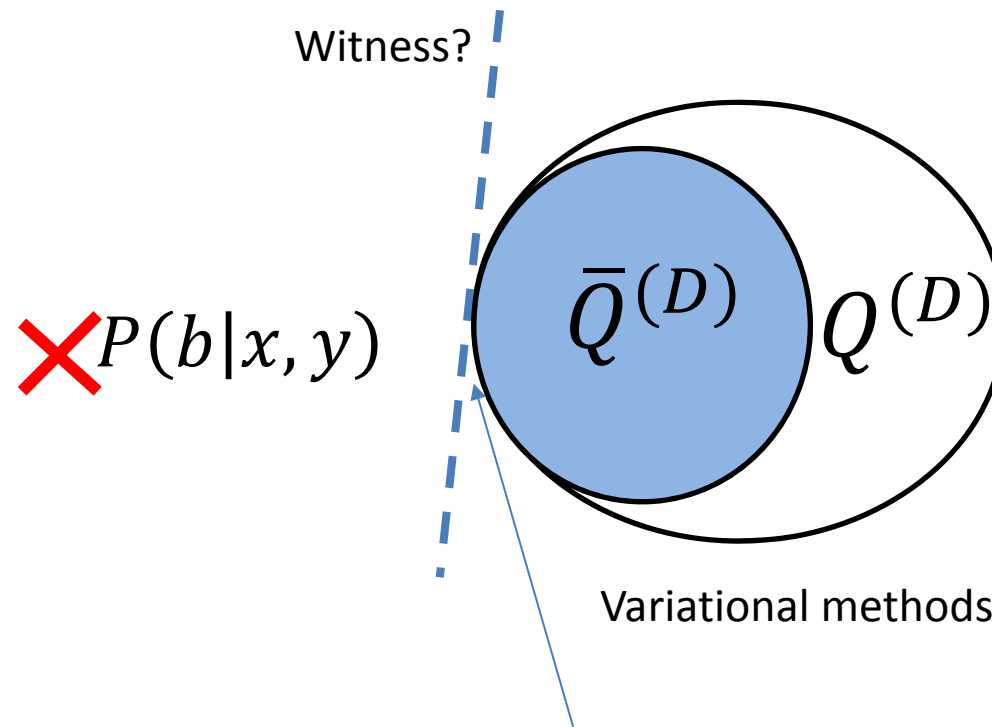
## Prepare-and-measure dimension witnesses



	$C_2$ (bit)	$Q_2$ (qubit)	$C_3$ (trit)	$Q_3$ (qutrit)	$C_4$ (quat)
$I_3$	3	$1 + 2\sqrt{2}$	5	5	5
$I_4$	5	6	7	7.9689	9

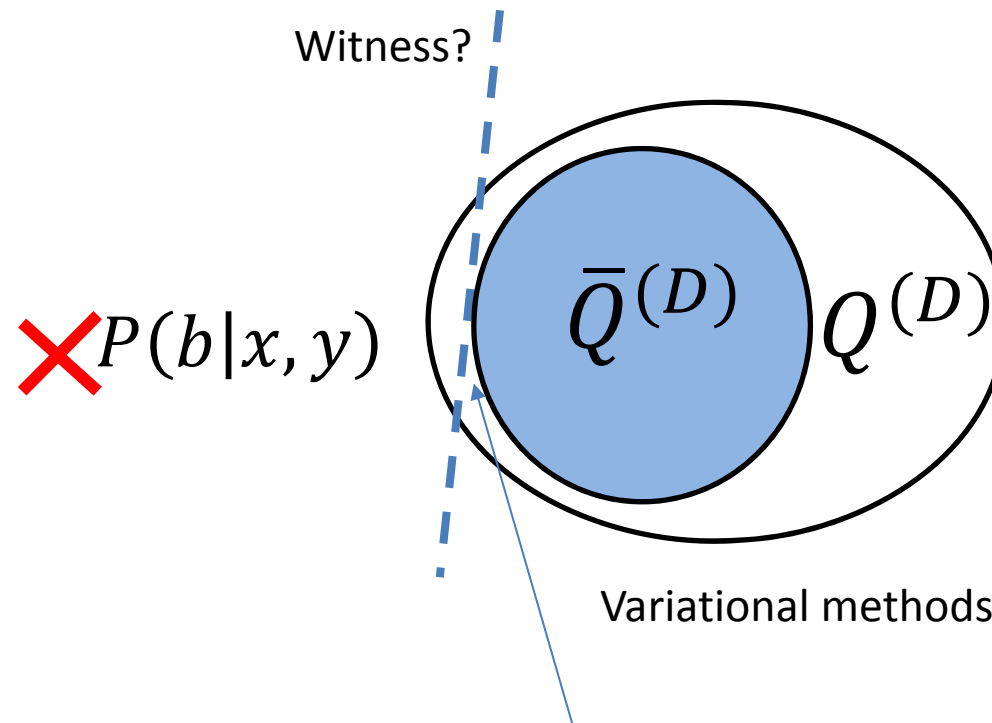
TABLE I: Classical and quantum bounds for the dimension witnesses  $I_3$  and  $I_4$ . Notably, these witnesses can distinguish classical and quantum systems of given dimensions.

In the lab



hope: the inequality is satisfied by the bigger set

In the lab



hope: the inequality is satisfied by the bigger set





```
' Sends text and returns with result including the  
' '>' Time out is 100 ms  
Public Function SendText(ByVal text)  
Dim timeNow  
With Form1  
SendText = ""  
.textIn = "" ' clear buffer  
.Sout(text) ' send it  
timeNow = .t1  
While .t1 < timeNow + 10  
If InStr(.textIn, ">") > 0 Then  
SendText = .textIn  
Exit While  
End If  
Application.DoEvents()  
End While  
End With  
End Function
```

```
MS-DOS Prompt - BBASIC  
FILE EDIT VIEW EXEC RUN DEBUG OPTIONS HELP  
ROBOT.BAS  
...  
Public Interface Drive routines  
Dim Parameters Computer Technology 1995  
GetStart 2.4  
...  
DIM SHARED Fischer*(0 To 511)  
DEF SEG = VARSEG(Fischer*(0))  
DEF SEG = Fischer*(512), UARPTR(Fischer*(0))  
FischerInstall  
HandOpen 38  
Ready 38  
...  
HandOpen 38  
Ready 38  
...  
F1=F1+Help2 <F8=Win+ou> <F2=Sub> <F3=Run> <F8=Step> 0064700
```

```
q.Vertices(descendant_index).state = q.Vertices(our_state_index).state  
...  
q.Vertices(descendant_index).state = "1"  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vert  
q.Vertices(descendant_index).qs = sev_qs // Check  
...  
q.Vertices(descendant_index).state = "1"  
q.Vertices(descendant_index).path = q.Vertices(our_state_index).path + q.Vert  
q.Vertices(descendant_index).qs = sev_qs // Check  
...  
q.Vertices(descendant_index).fn = Calculate(q.Vertices(descendant_index).qs)
```

	$C_2$ (bit)	$Q_2$ (qubit)	$C_3$ (trit)	$Q_3$ (qutrit)	$C_4$ (quat)
$I_3$	3	$1 + 2\sqrt{2}$	5	5	5
$I_4$	5	6	7	7.9689	9

TABLE I: Classical and quantum bounds for the dimension witnesses  $I_3$  and  $I_4$ . Notably, these witnesses can distinguish classical and quantum systems of given dimensions.

We proved that all these witnesses were sound,  
hence validating of the conclusions of the  
experimental papers

M. Hendrych, R. Gallego, M. Micuda, N. Brunner, A. Acín and J. P. Torres, Nat. Phys. 8, 588 (2012).  
J. Ahrens, P. Badziag, A. Cabello, M. Bourennane, Nat. Phys. 8, 592 (2012).

# Hierarchy for real quantum systems

$$i = 1$$

Generate *real* random projectors  $\rho_x^i, F_b^{i,y} \in B(\mathbb{C}^D)$  (with app. ranks)



Build moment matrix  $\Gamma^i$

$$\begin{aligned}\Gamma_{\mathbb{I},\mathbb{I}}^i &= \text{tr}(\mathbb{I}) = 1, \\ \Gamma_{\rho_x, F_b^y}^i &= \text{tr}(\rho_x^i F_b^{i,y}), \\ &\vdots\end{aligned}$$



$$i = i + 1$$

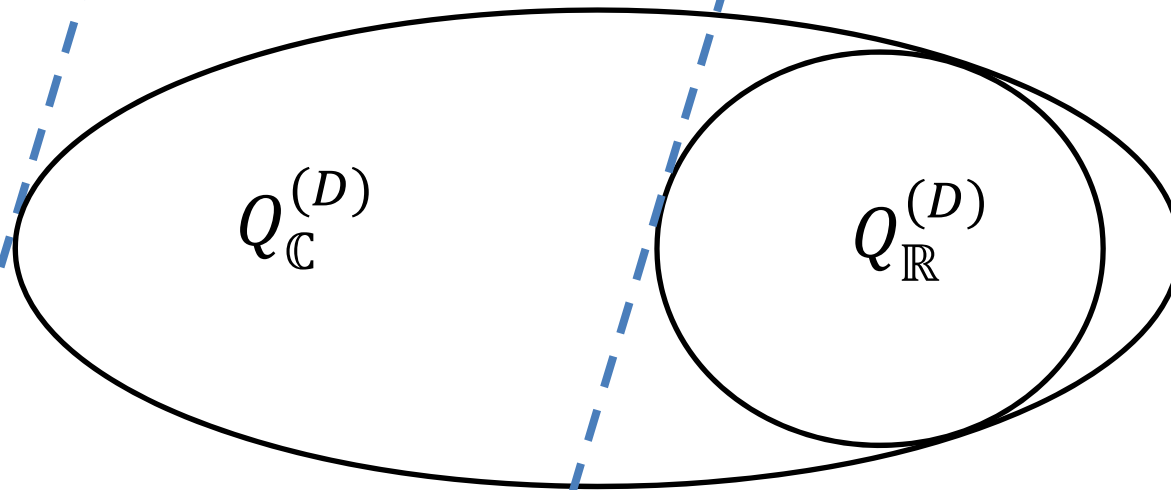
Repeat until the moment matrix  $\Gamma^{(N+1)}$  is a linear combination of  $\Gamma^{(1)}, \dots, \Gamma^{(N)}$




# Real vs. Complex quantum mechanics

Your technology allows you to go beyond complex quantum mechanics in dimension  $d$

$P(b|x, y)$



Your technology allows you to go beyond real quantum mechanics in dimension  $d$

$$\max \sum_{x,y,b} w_b^{x,y} P(b|x, y)$$


# Random Access Codes (RAC)

Alice



Dimension  $D$

Bob



$y \in 1, \dots, n$

$$f(\bar{x}, y) = x_y$$



$b \in \{0,1\}$

complex qubits $P_{max}(3 \rightarrow 1) \approx 0.788675$	real qubits $P_{max}(3 \rightarrow 1) \approx 0.7696723$
---	---

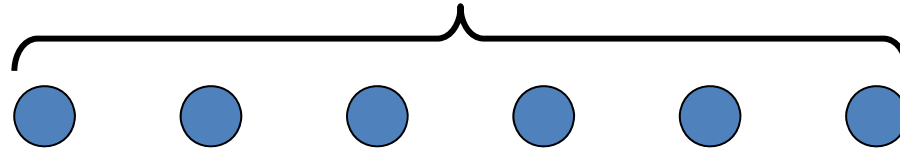
# The structure of Matrix Product States

MN and T. Vértesi, arXiv:1509.04507.

s-local term  
(acts on particles  
 $j, \dots, j+s-1$ )

$$H = \sum_{j=1}^{n-s} h_j$$

k particles



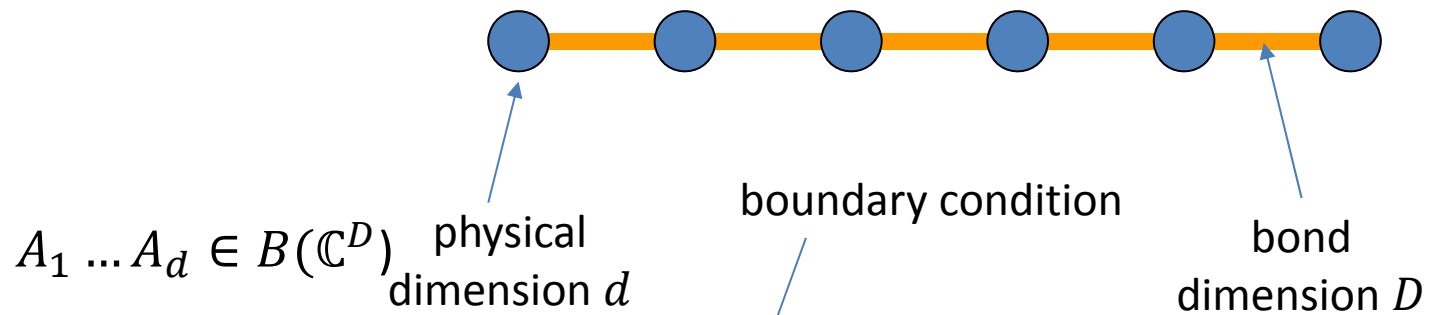
physical  
dimension  $d$

ground state  $|\varphi\rangle \in \mathbb{C}^{d^k}$



Impossible to even  
*store* for high  $k$

## Matrix product states



$$|\psi\rangle = \sum_{i_1, \dots, i_k} \text{tr}(\sigma A_{i_1} \dots A_{i_k}) |i_1\rangle \otimes \dots \otimes |i_k\rangle$$

D. Perez-Garcia, F. Verstraete, M.M. Wolf and J.I. Cirac, Quantum Inf. Comput. 7, 401 (2007).



## Matrix product states



Features:

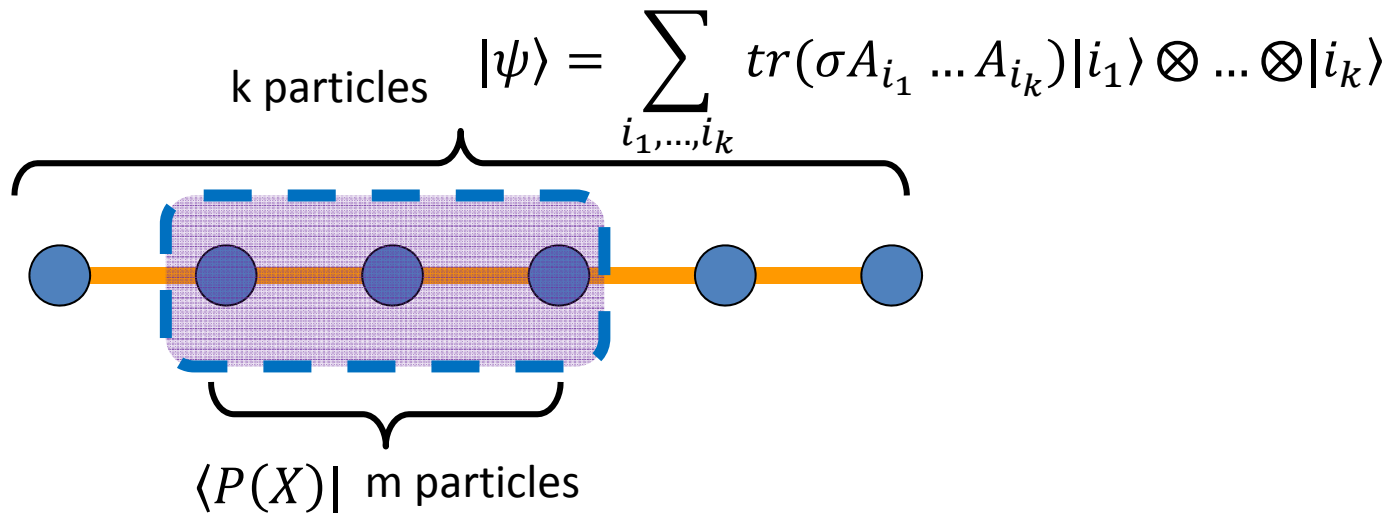
- a) Efficient computation of expectation values
- b) Calculations in the thermodynamical limit  $k \rightarrow \infty$

## Correspondence between polynomials and states

$P(X) = P(X_1, \dots, X_d)$ ,  
homogeneous polynomial  
of degree  $m$

$$P(X) = \sum_{\vec{i}} p_{\vec{i}} X_{i_1} \dots X_{i_m} \quad \Rightarrow \quad |P(X)\rangle = \sum_{\vec{i}} p_{\vec{i}}^* |i_1, \dots, i_m\rangle$$

## Overlap between a polynomial and a MPS




$$\langle P(X) | \psi \rangle = \sum_{\substack{i_1, \dots, i_s \\ i_{s+m+1}, \dots, i_k}} \text{tr}(\sigma A_{i_1} \dots A_{i_s} P(A) A_{i_{s+m+1}} \dots A_{i_k}) |i_1, \dots, i_s, i_{s+m+1}, \dots, i_k\rangle$$

↑ "defect"

Existence of annihilation operators

$$\langle P(X)|\psi\rangle = \sum_{\substack{i_1, \dots, i_s, \\ i_{s+m+1}, \dots, i_k}} \text{tr}(\sigma A_{i_1} \dots A_{i_s} P(A) A_{i_{s+m+1}} \dots A_{i_k}) |i_1, \dots, i_s, i_{s+m+1}, \dots, i_k\rangle$$

$P(X)$ , MPI for dimension  $D$    $\langle P(X)|\psi\rangle = 0$ , for *all* MPS with bond dimension  $D$ !!

There exist local operators  $h$  which annihilate all MPS of bond dimension  $D$  or smaller


Actually, for high  $m$ , almost all  $m$ -local operators are annihilation operators, because

Local dimension of  $m$  particles  $d^m$

Local dimension of MPS subspace  $poly(m)$

Existence of cut-and-glue operators

$$\langle P(X)|\psi\rangle = \sum_{\substack{i_1, \dots, i_s, \\ i_{s+m+1}, \dots, i_k}} \text{tr}(\sigma A_{i_1} \dots A_{i_s} P(A) A_{i_{s+m+1}} \dots A_{i_k}) |i_1, \dots, i_s, i_{s+m+1}, \dots, i_k\rangle$$

$P(X)$ , central polynomial for dimension D
 
 $\langle P(X)|\psi\rangle = p(A) |\psi'\rangle$

scalar

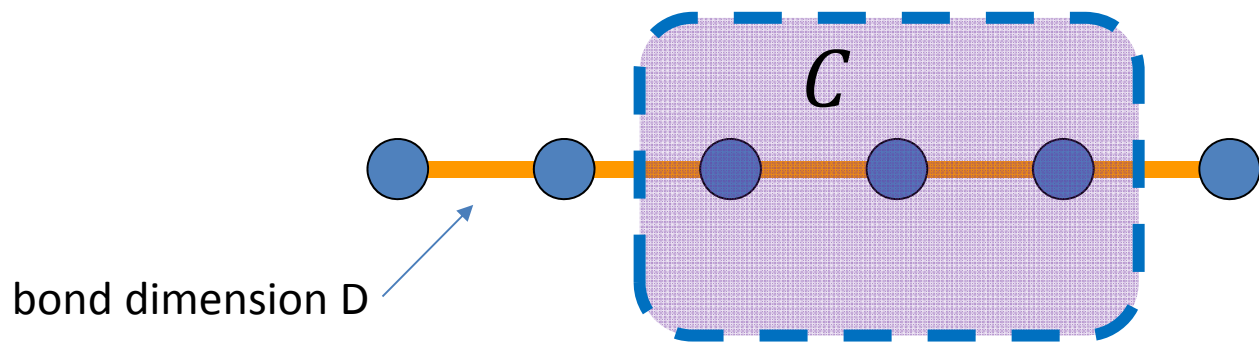
$$|\psi'\rangle = \sum_{\substack{i_1, \dots, i_s, \\ i_{s+m+1}, \dots, i_k}} \text{tr}(\sigma A_{i_1} \dots A_{i_s} A_{i_{s+m+1}} \dots A_{i_k}) |i_1, \dots, i_s, i_{s+m+1}, \dots, i_k\rangle$$



$\{P_j(X)\}$ , basis of  
homogeneous central  
polynomials for  
dimension  $D$  of degree  $m$



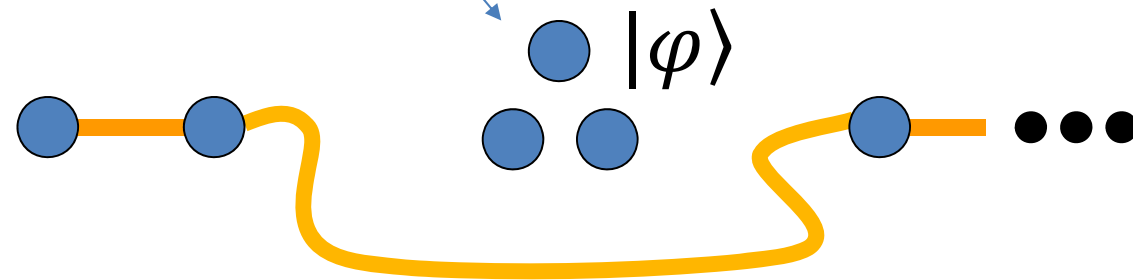
$$C = \sum_j |j\rangle\langle P_j(X)|$$



$$|\psi\rangle = \sum_{i_1, \dots, i_k} \text{tr}(\sigma A_{i_1} \dots A_{i_k}) |i_1\rangle \otimes \dots \otimes |i_k\rangle$$

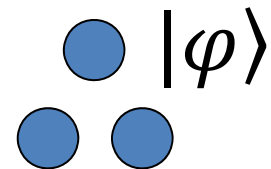
# C, cut-and-glue operator

m particles projected onto a pure state (cut from the chain)



Two ends glued back

C, cut-and-glue operator



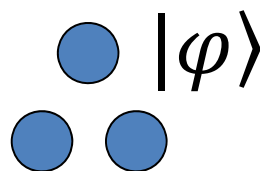
$$|\psi\rangle = \sum_{i_1, \dots, i_k} \text{tr}(\sigma A_{i_1} \dots A_{i_{k-m}}) |i_1\rangle \otimes \dots \otimes |i_{k-m}\rangle$$

C, cut-and-glue operator for dimension D

$\rho_k$ , k-site MPS of bond dimension D



$(C \otimes 1)\rho_k(C^\dagger \otimes 1)$ , separable



$$|\psi\rangle = \sum_{i_1, \dots, i_k} \text{tr}(\sigma A_{i_1} \dots A_{i_{k-m}}) |i_1\rangle \otimes \dots \otimes |i_{k-m}\rangle$$

Theoretical complexity of the SDP  
hierarchy and connection with  
MPS

## Complexity of implementing the $k$ th relaxation

SDP constraint

$$\Gamma_k = \left( \quad \right) \left. \vphantom{\left( \quad \right)} \right\} \exp(k)$$

$$|c^k| = \exp(k) \quad \text{free variables}$$

Free-dimensional problems have the same complexity as dimension-constrained ones!?

## Connection with MPSs

$$\Gamma_k = \sum_{i_1, \dots, i_k} \sum_{j_1, \dots, j_k} \text{tr}(\tilde{X}_{j_n} \dots \tilde{X}_{j_1} |\psi\rangle \langle \psi| \tilde{X}_{i_1} \dots \tilde{X}_{i_k}) |\vec{i}\rangle \langle \vec{j}|$$

Feasible moment matrices are extendible MPSs of bond dimension  $D!!!$



# Complexity of implementing the kth relaxation

SDP constraint

$$\Gamma_k = \left[ \begin{array}{c} \text{SDP constraint} \\ \text{SDP constraint} \\ \text{SDP constraint} \end{array} \right] \left. \vphantom{\left[ \begin{array}{c} \text{SDP constraint} \\ \text{SDP constraint} \\ \text{SDP constraint} \end{array} \right]} \right\} \exp(k)$$

SuppEspan{poly(k) k-site MPS of bond dimension D}

$$|c^k| = \exp(k) \quad \text{free variables}$$

$c^k \in \text{span}\{\text{poly(k) 2k-site MPS of bond dimension D}\}$

Complexity of implementing the  $k$ th relaxation

$$\tilde{\Gamma}^k = \left( \quad \right) \} \text{poly}(k)$$

$$|\tilde{c}^k| = \text{poly}(k)$$

Finite dimensions are exponentially  
easier to characterize than infinite  
dimensions!!!

Extra PSD conditions to boost speed of convergence

$$\Gamma^k = \sum p_j \rho_{MPS}^j$$

Extra PSD conditions to boost speed of convergence

$$\Gamma^k = \sum p_j \rho_{MPS}^j$$

cut-and-glue-operator



$C\Gamma^k C^\dagger$ , separable operator

Extra PSD conditions to boost speed of convergence

$$\Gamma^k = \sum p_j \rho_{MPS}^j$$

cut-and-glue-operator



$C\Gamma^k C^\dagger$ , separable operator



New PSD conditions

$C\Gamma^k C^\dagger$ , PPT

# Conclusions

-Simple, easy-to-program technique to enforce dimension constraints in noncommutative polynomial optimization.

-Plenty of numerical evidence suggests that it is effective.

-It can be combined with other hierarchies, such as MLHG or BFS.

T. Moroder, J.-D. Bancal, Y.-C. Liang, M. Hofmann and O. Gühne, Phys. Rev. Lett. 111, 030501 (2013).

M. Berta, O. Fawzi and V. B. Scholz, arXiv:1506.08810.

THE END