

Causality, Consistency, Complexity

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Based on collaboration and discussion with:

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Alberto Montina, Benno Salwey and Andreas Winter*

Tainan, Taiwan, December 2015

Overview

Causality

Overview

Causality



Correlations

Overview

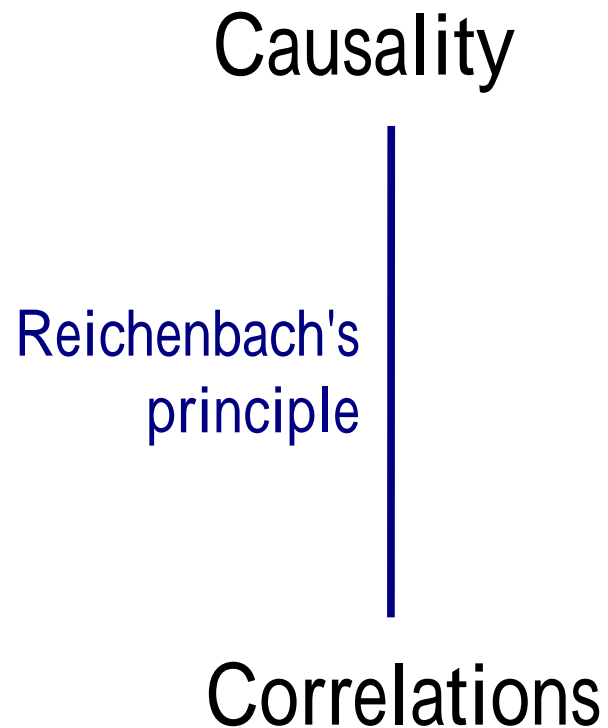
Causality

Reichenbach's
principle

Correlations



Overview



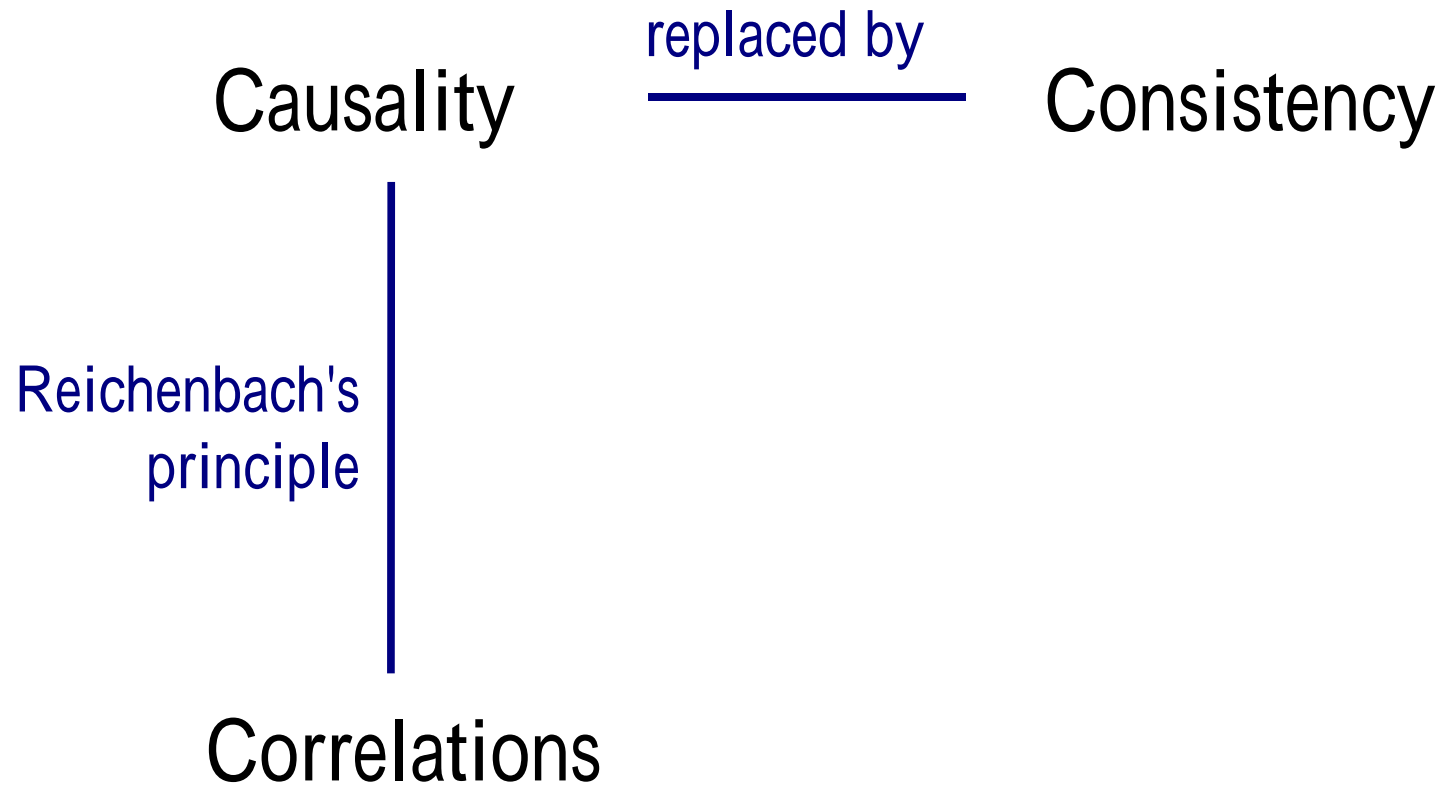
The law of causality [...] is a relic of a bygone age, surviving, like the monarchy, only because it is erroneously supposed to do no harm

Bertrand Russell

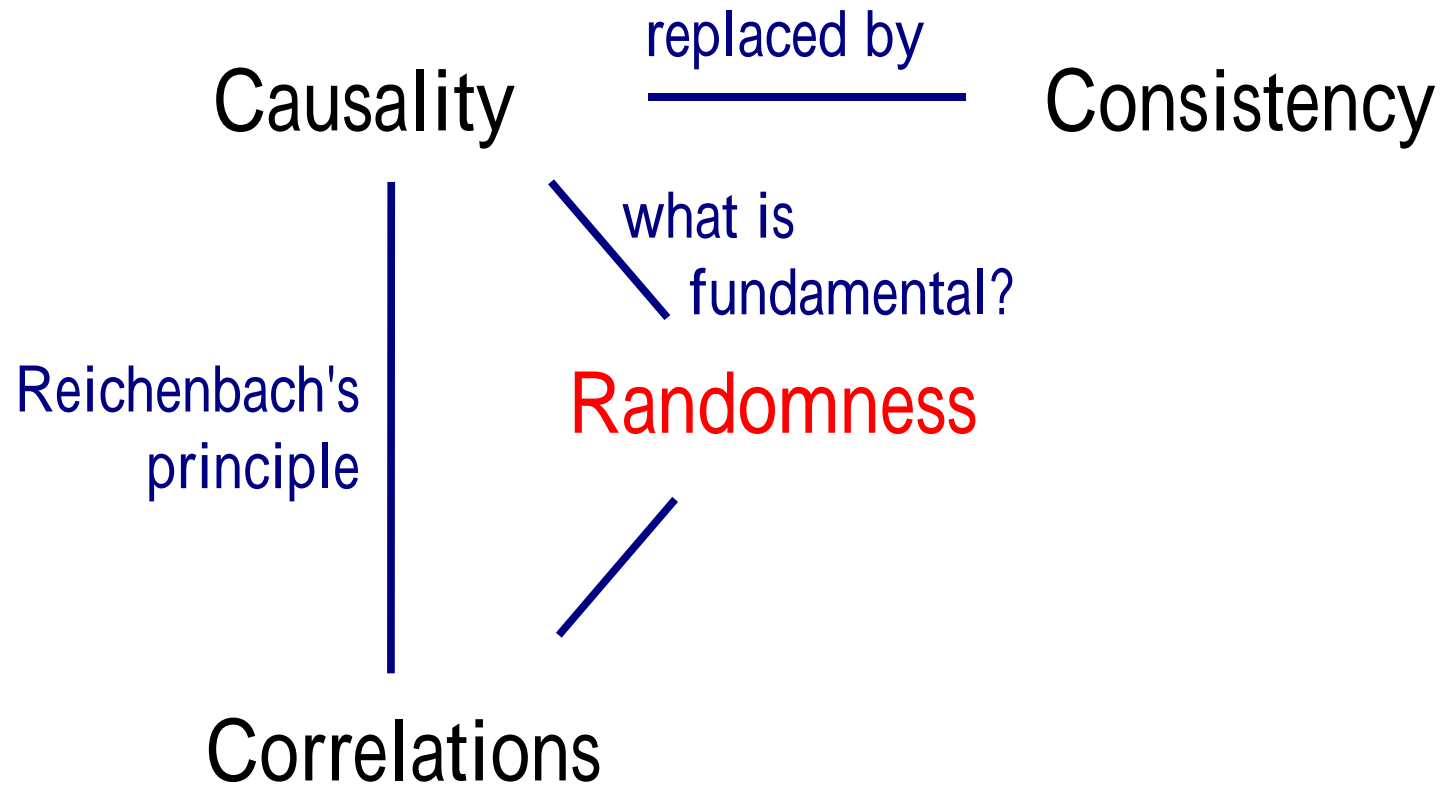


1913

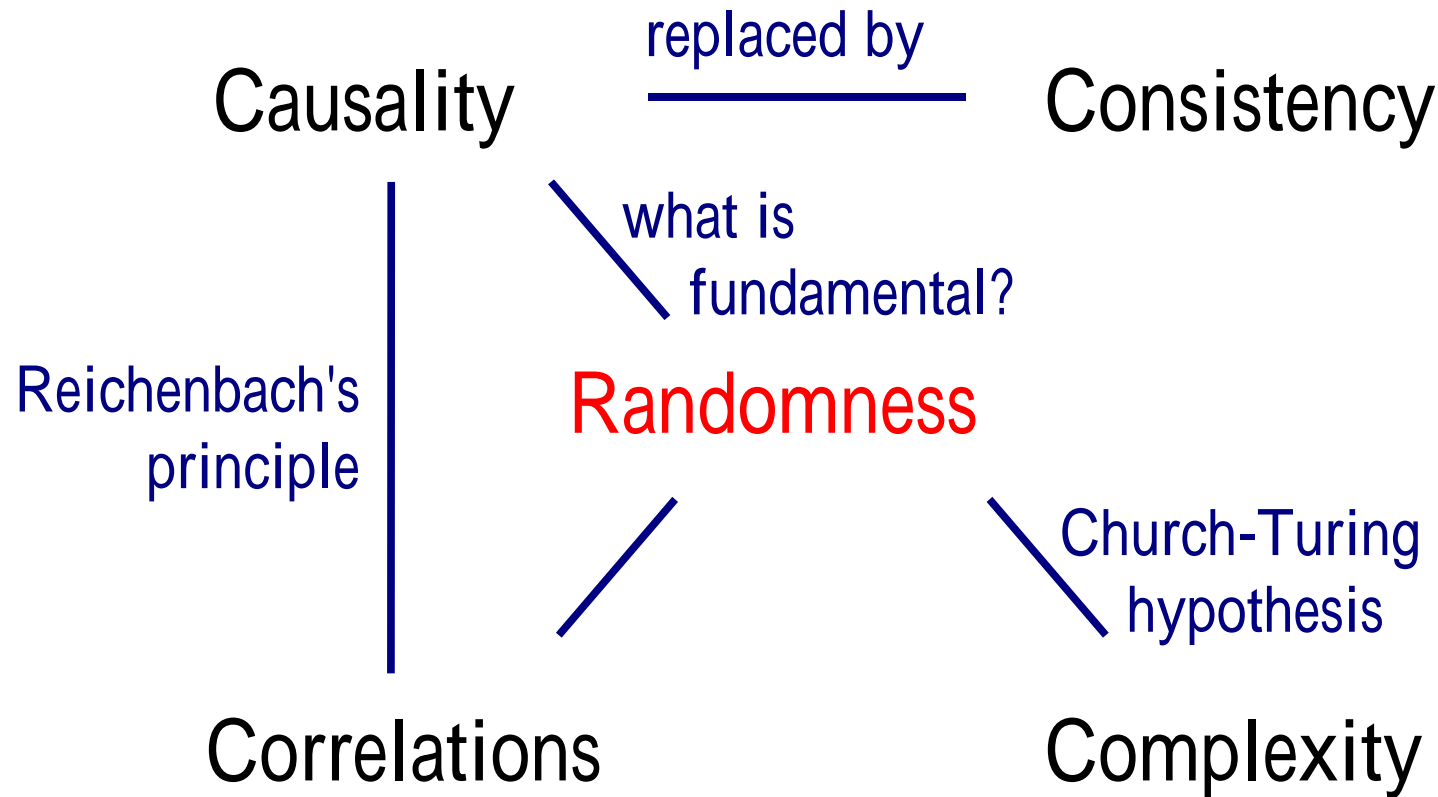
Overview



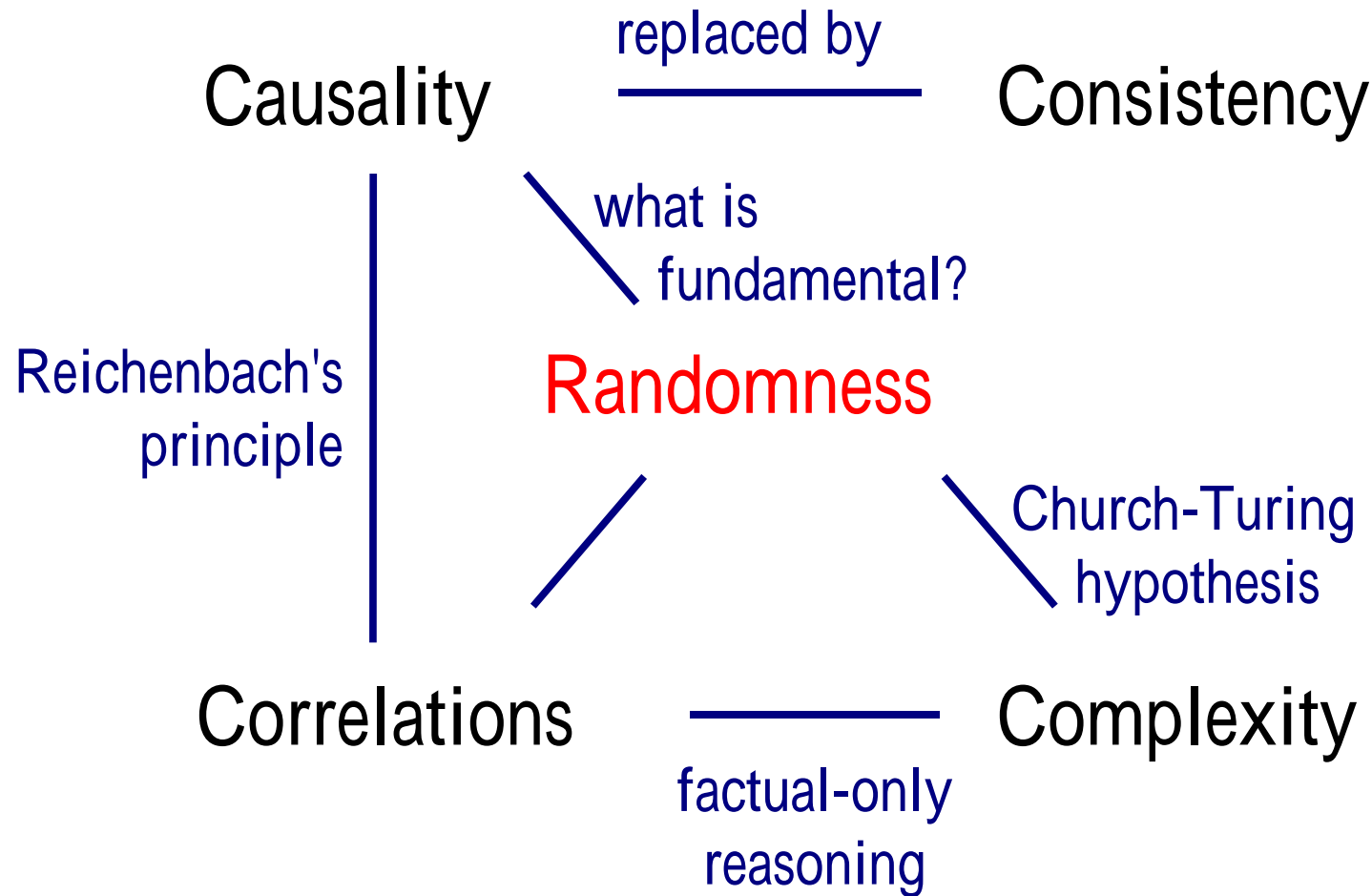
Overview



Overview



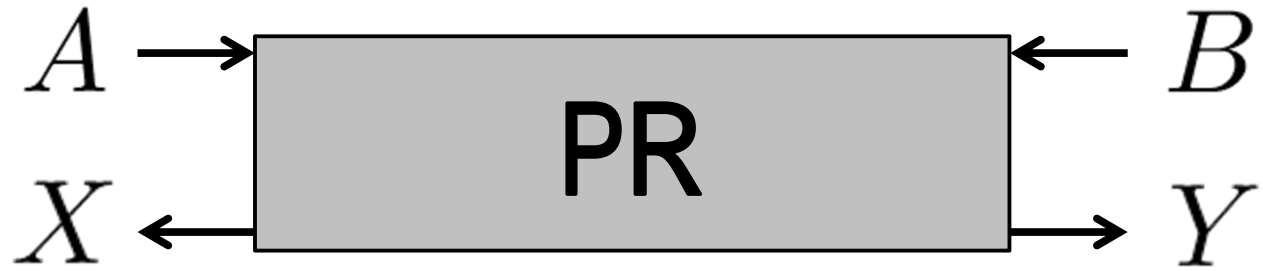
Overview



Non-Local Correlations



John Bell
1964



$$X \oplus Y = A \odot B$$



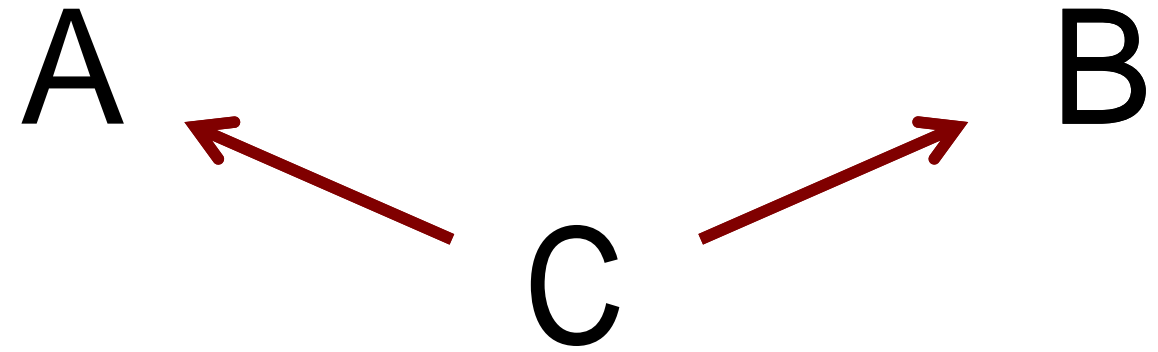
1994
Sandu Popescu



Daniel Rohrlich

Explaining Correlations in a Causal Structure

1) Common Cause



2) Influence



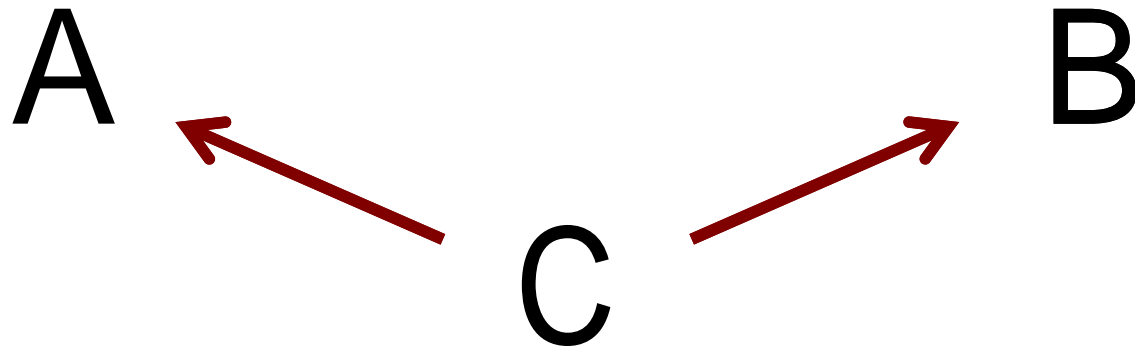
1956

Hans Reichenbach



Explaining Correlations in a Causal Structure

1) Common Cause



2) Influence



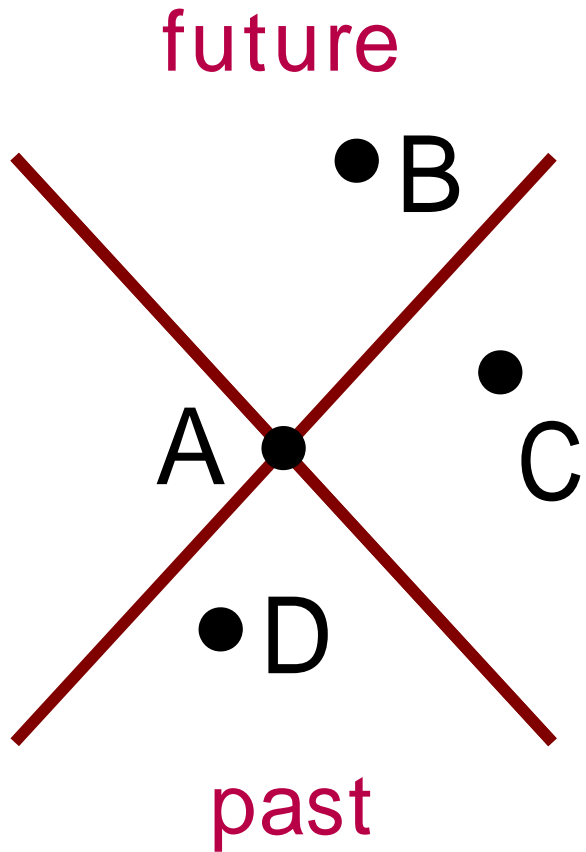
3) Drop the Causal Structure?

1956

Hans Reichenbach



Causality and Randomness



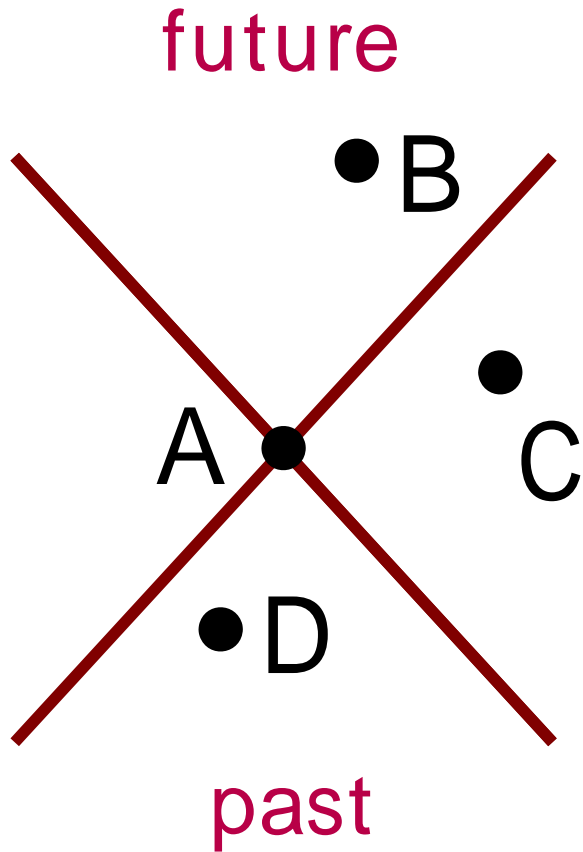
A is **free** if it is independent of C and D (of all except its **future**)



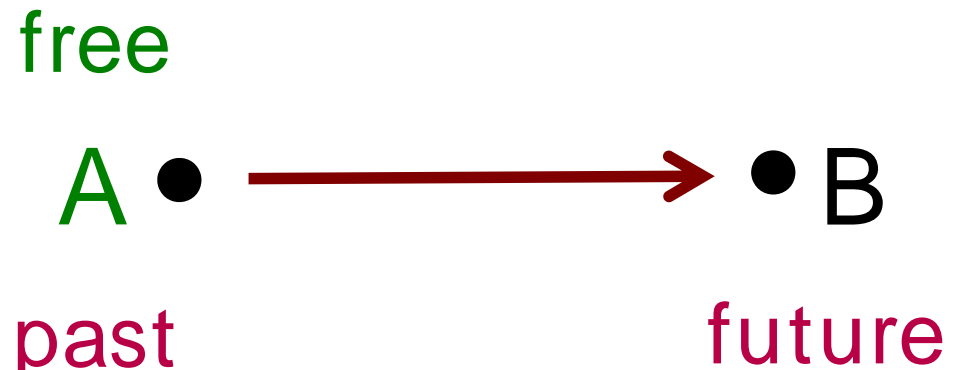
Renner

Colbeck

Causality and Randomness

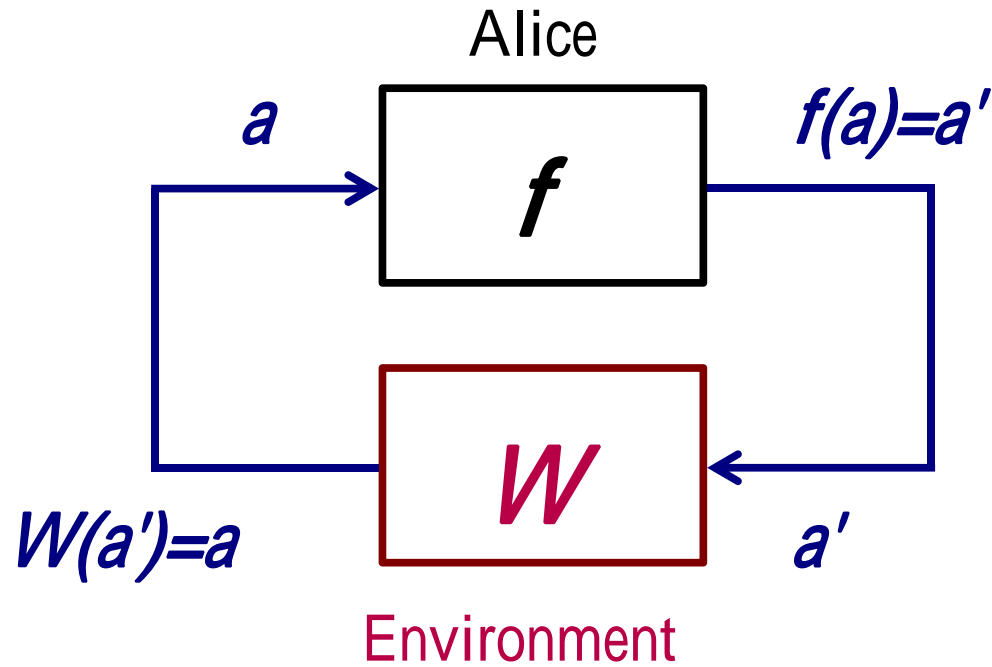


If A is **free** and B is correlated with A, then B is in A's **future**



A is **free** if it is independent of C and D (of all except its **future**)

Dropping the Causal Structure



Oreshkov



Costa



Brukner



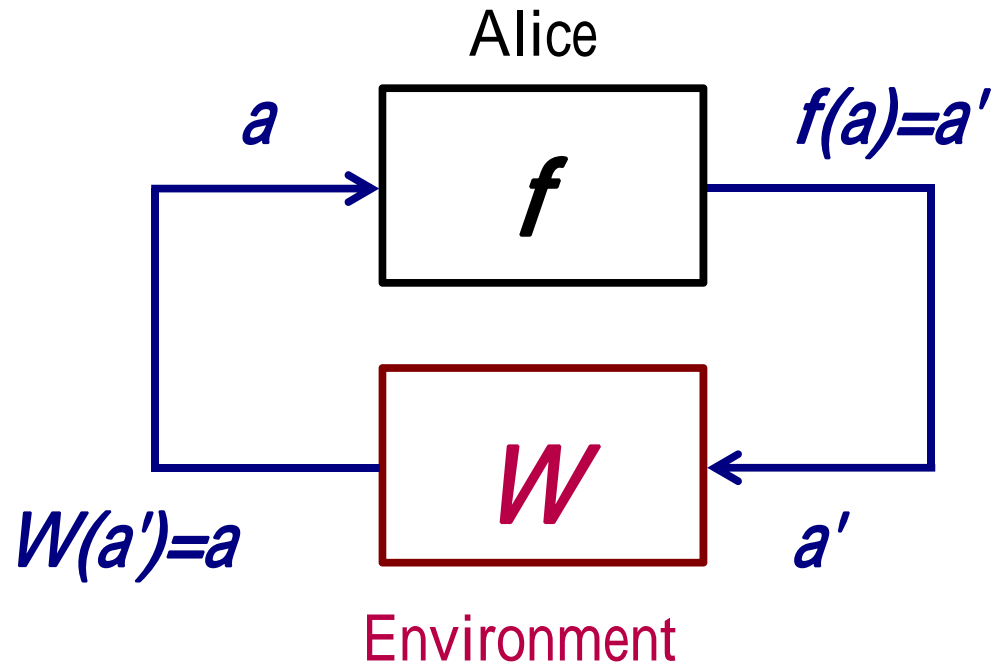
Baumeler



Feix

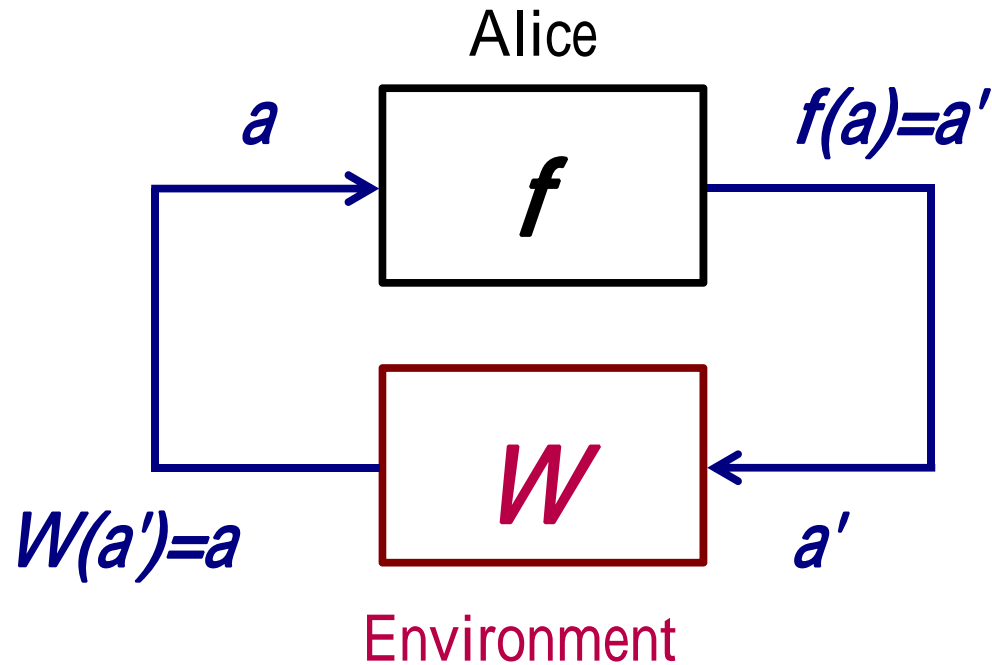


Dropping the Causal Structure



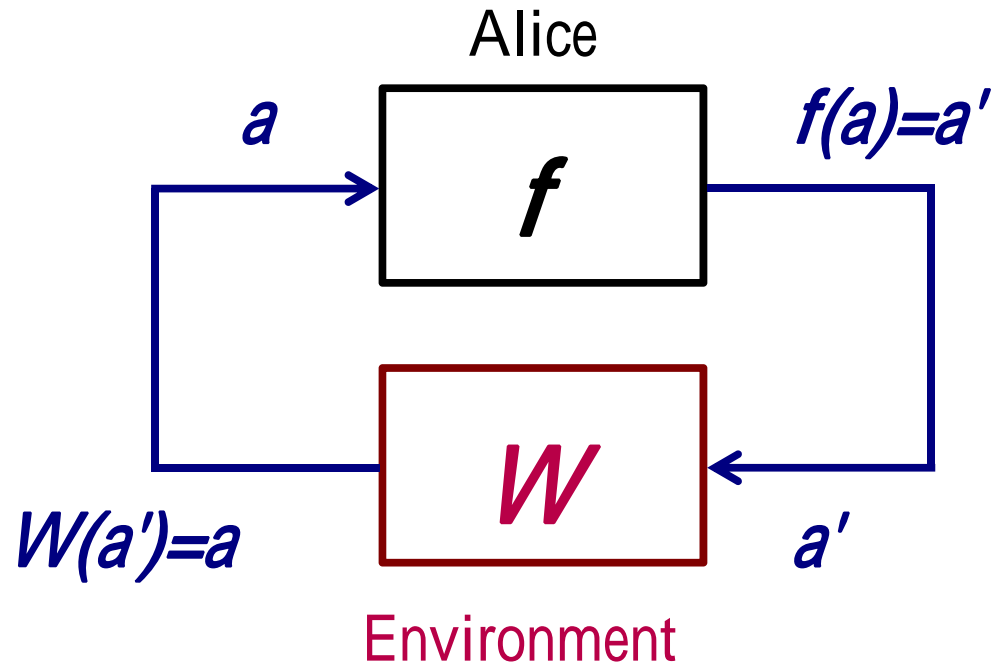
W is *consistent* if no logical contradiction arises, whatever Alice does

Dropping the Causal Structure



Example 1: $W(a') = a \oplus 1$
Contradiction for $f(a) = a$
("Grandfather paradox")

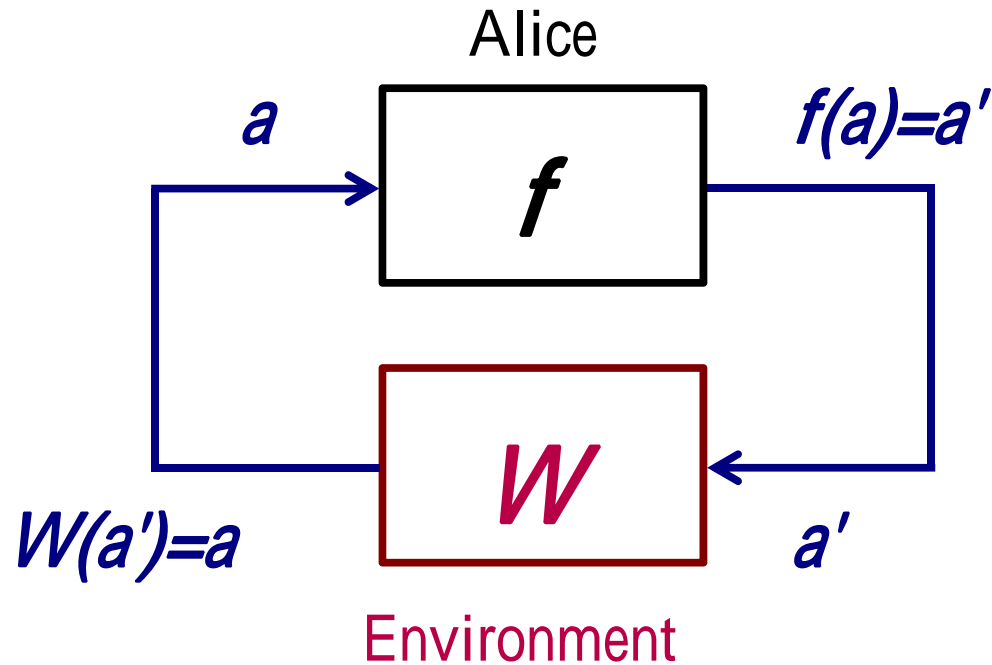
Dropping the Causal Structure



Example 2: $W(a') = a$

Contradiction for $f(a) = a \oplus 1$

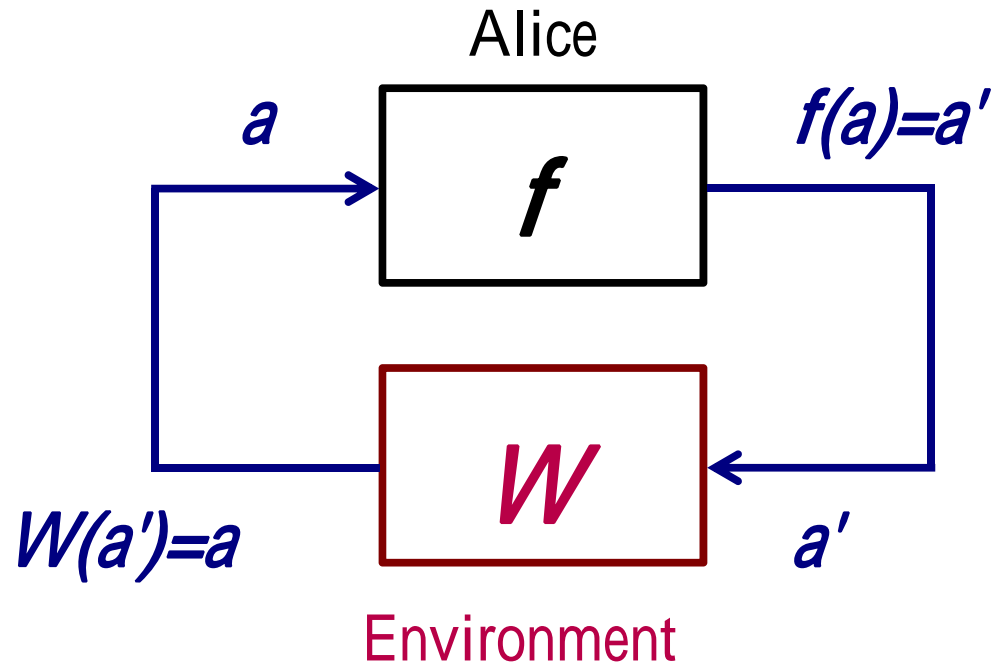
Dropping the Causal Structure



Example 3: $W(a') = 0$
No contradiction.

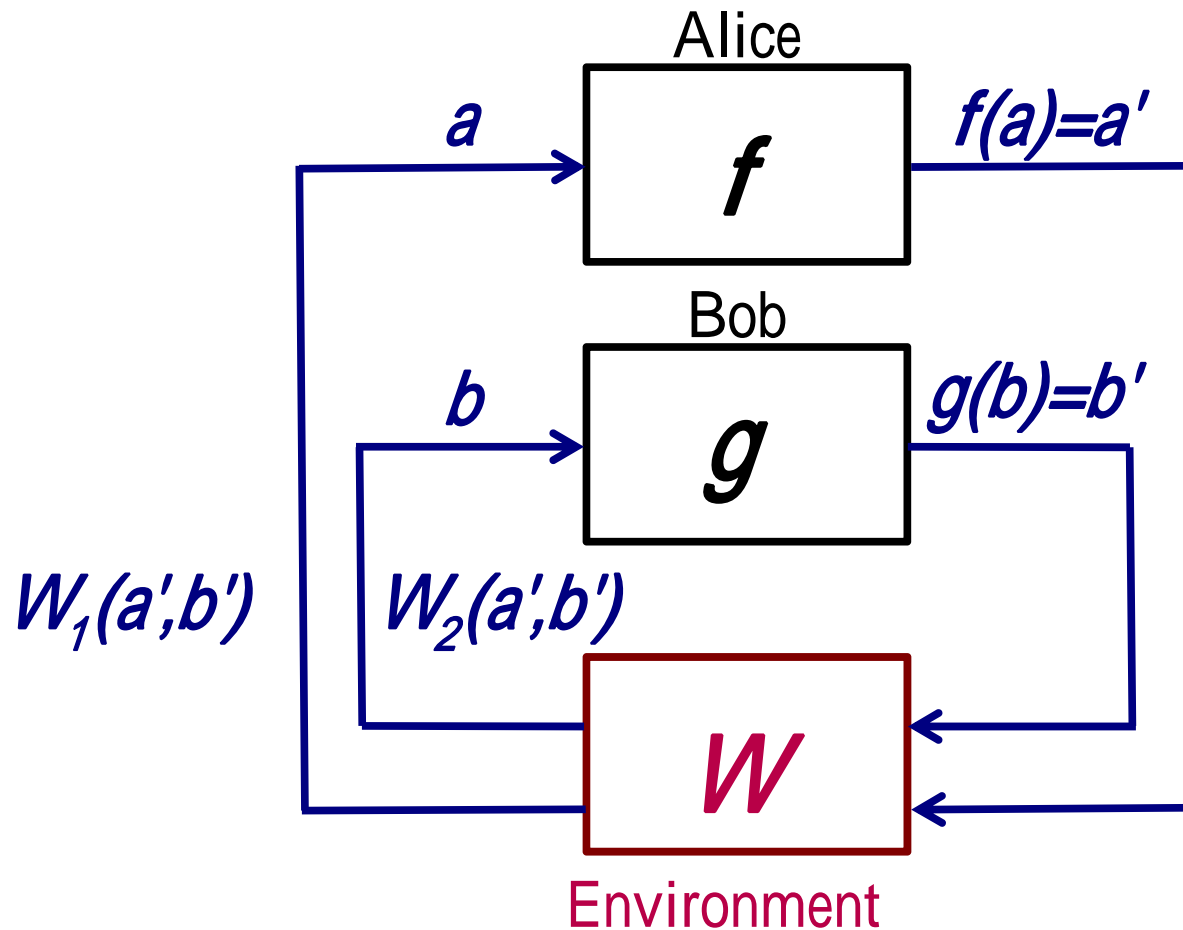
But: can be modeled *causally*: $W \xrightarrow{0} A \xrightarrow{f} 0$

Dropping the Causal Structure



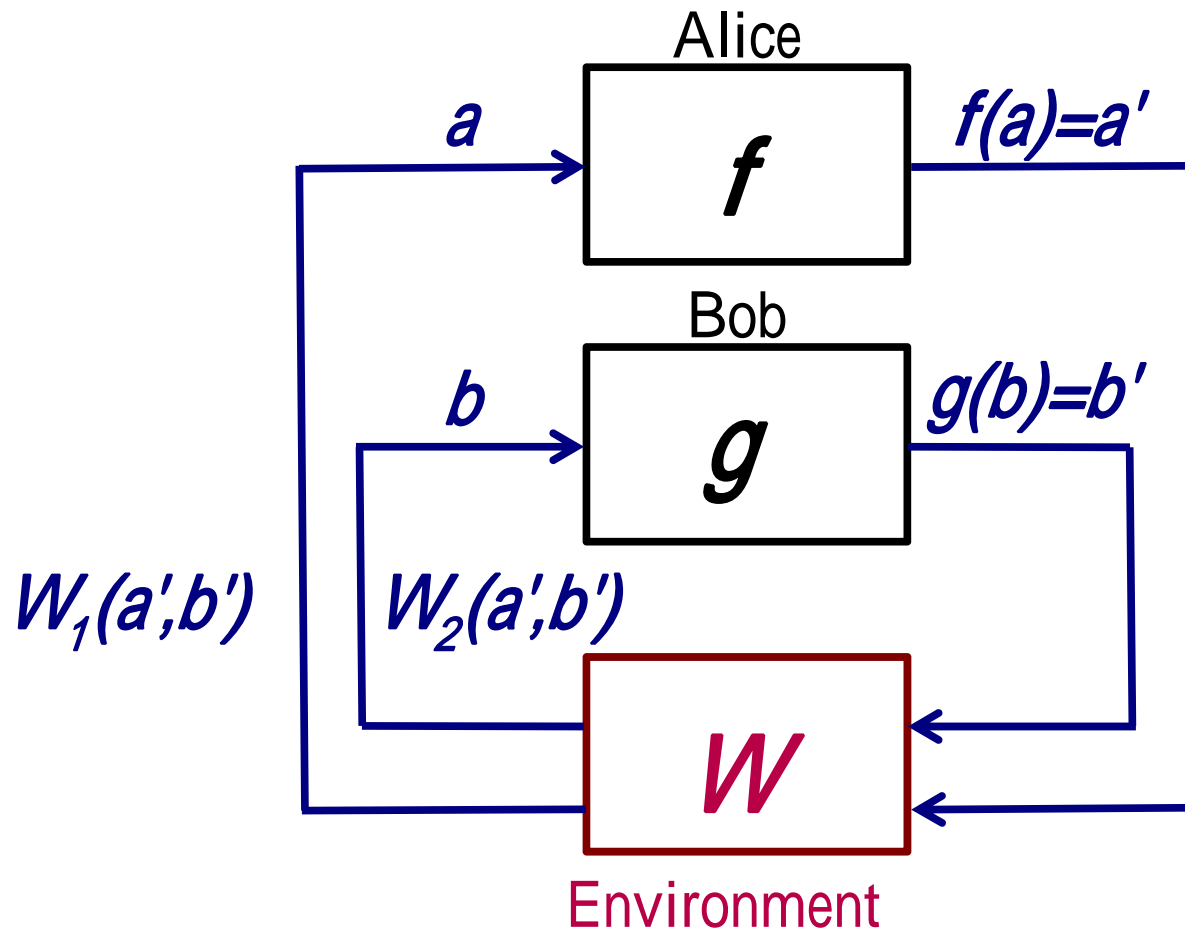
Do there exist *logically consistent* scenarios that are *not causal*?

Dropping the Causal Structure



Do there exist *logically consistent* scenarios that are *not causal*?

Dropping the Causal Structure



Oreshkov



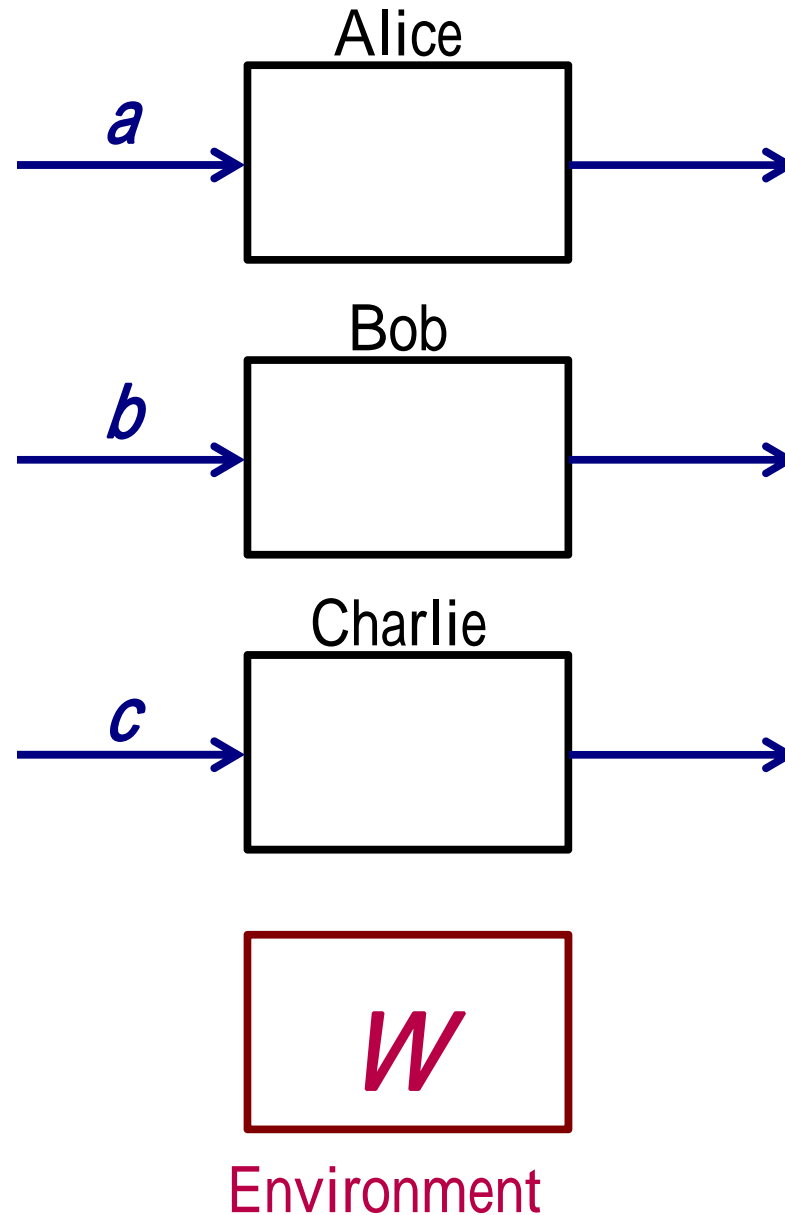
Costa



Brukner

Two parties: **No.** *Consistency implies causality*

Consistency *without* Causality is Possible



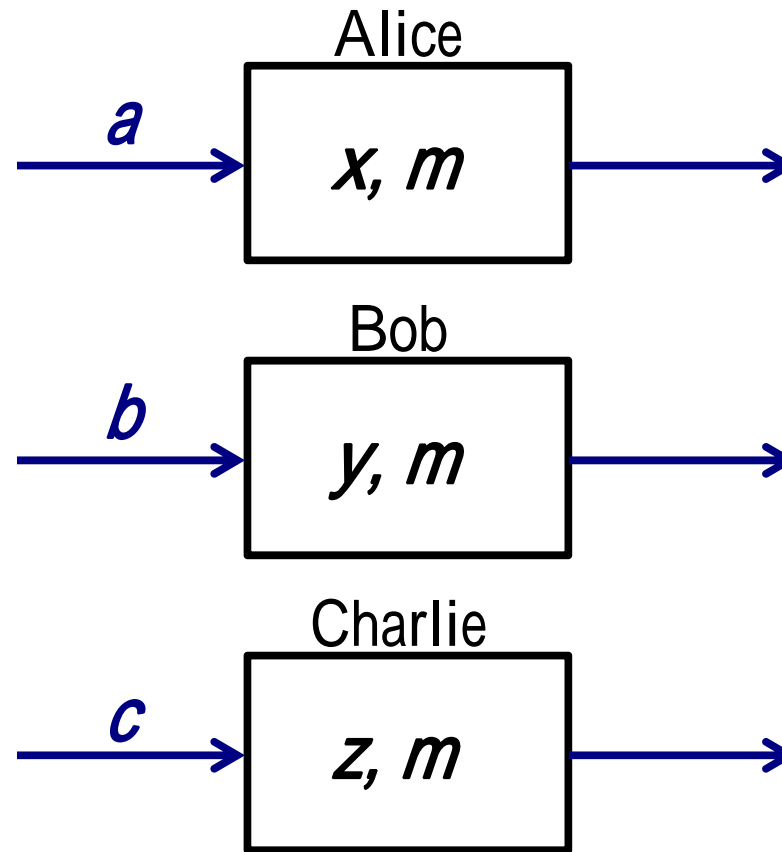
Baumeler



Feix



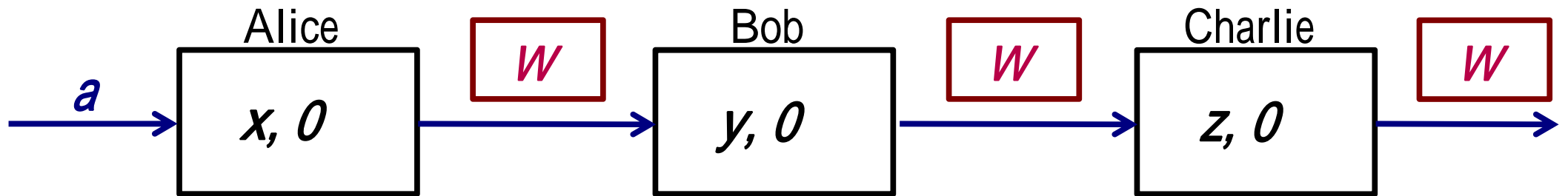
The Three-Party Game



Game won if

$m = 0$	\implies	$a = y \oplus z$
$m = 1$	\implies	$b = x \oplus z$
$m = 2$	\implies	$c = x \oplus y$

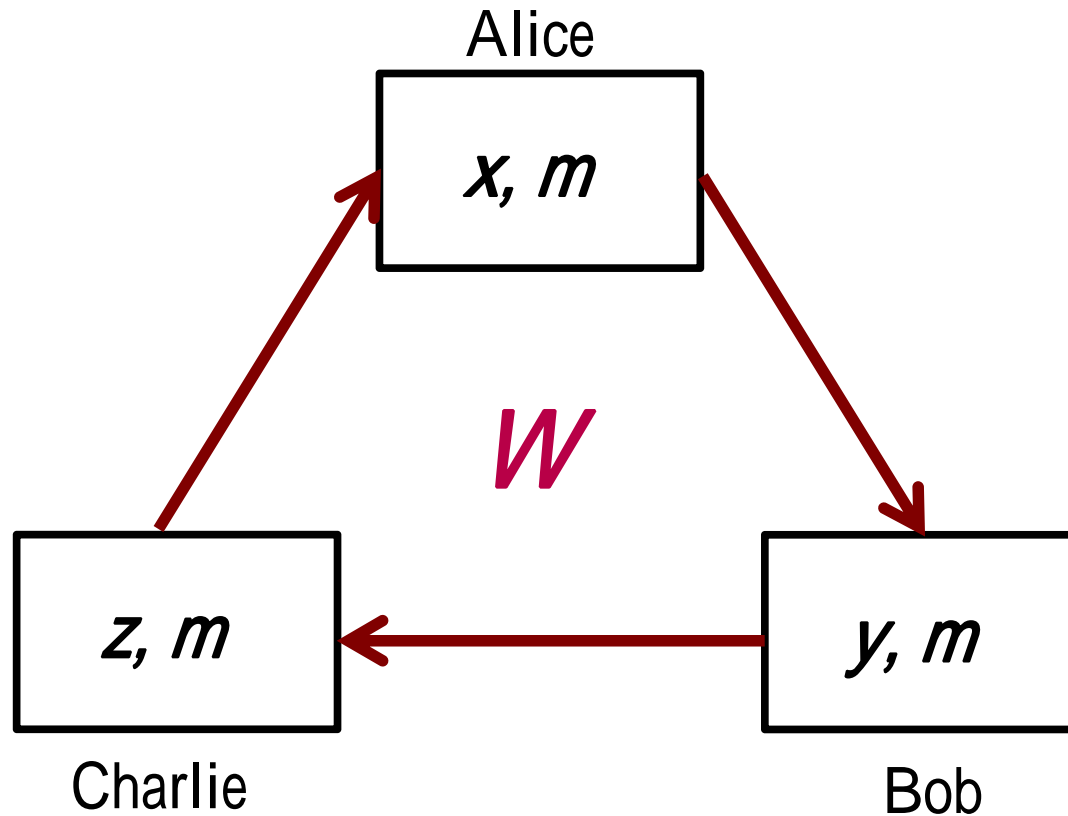
Game Cannot Always Be Won with Fixed Order



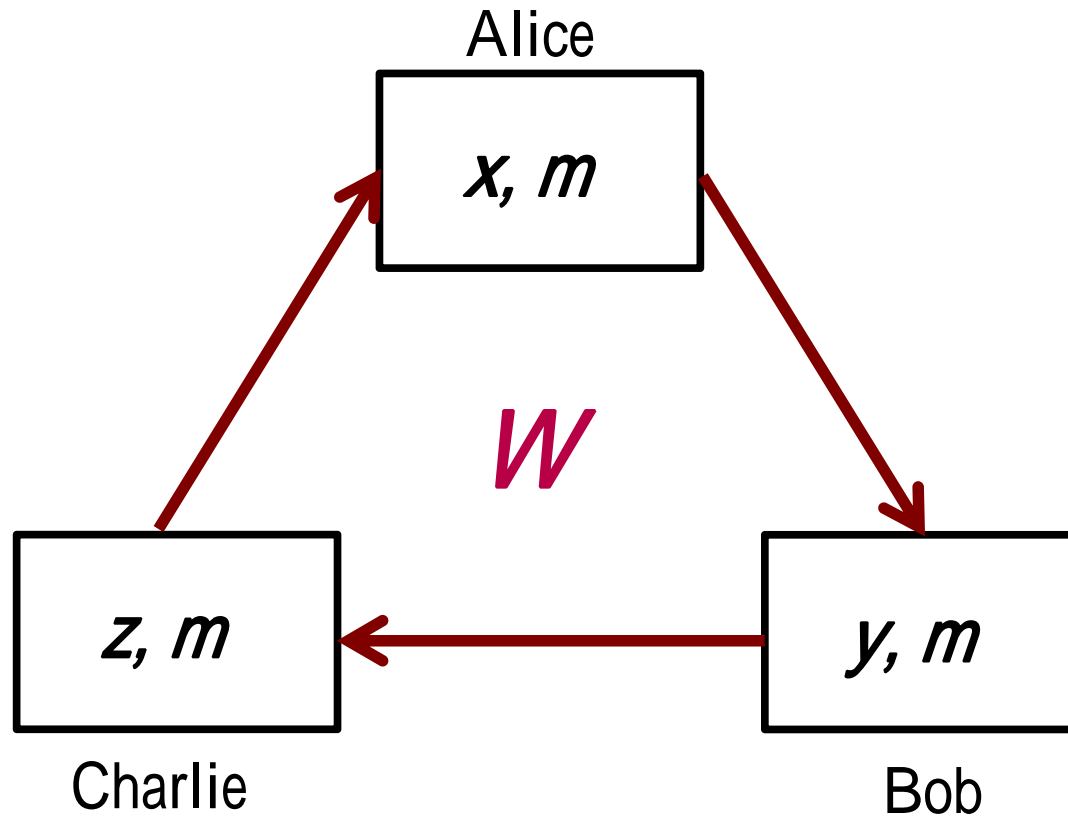
$$\text{Prob}[a = y \oplus z] = \frac{1}{2}$$

Otherwise: Game lost

Always Winning the Game with a Consistent W

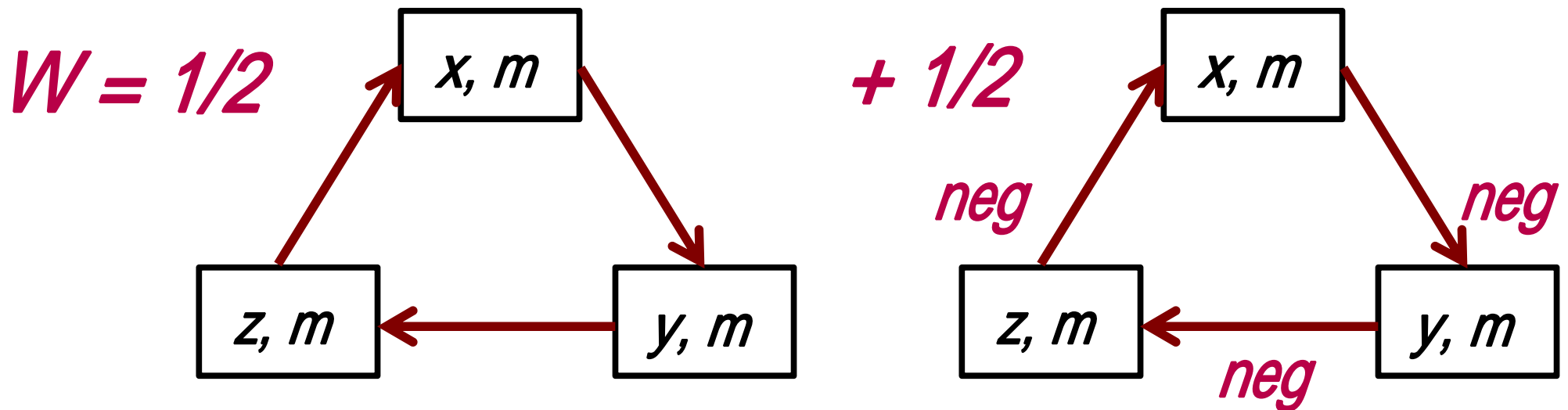


Always Winning the Game with a Consistent W



W inconsistent

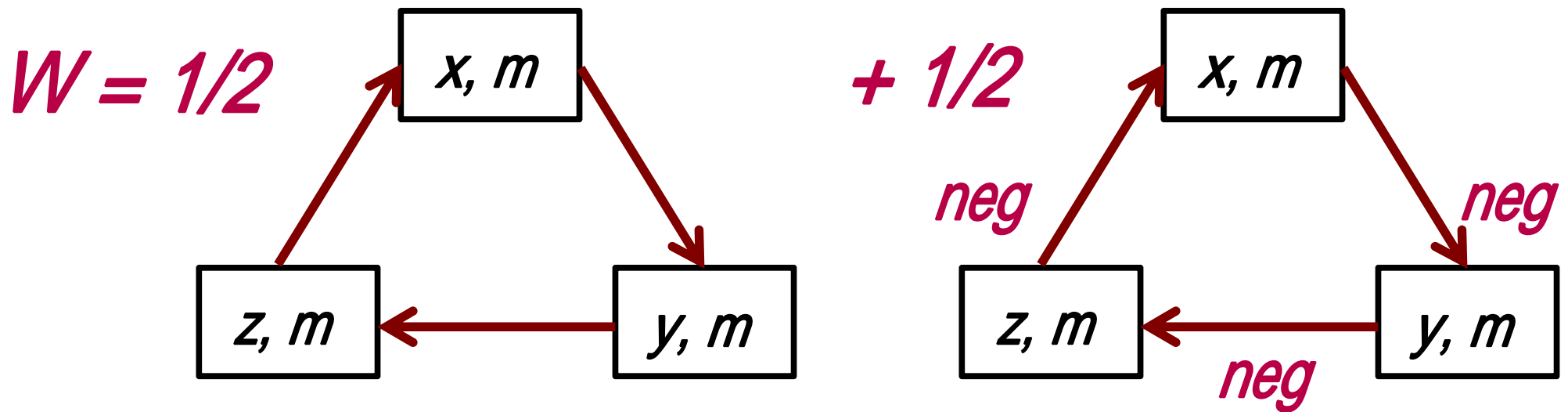
Always Winning the Game with a Consistent W



W is *consistent*:

The value a party receives is *independent* of what she

Always Winning the Game with a Consistent W

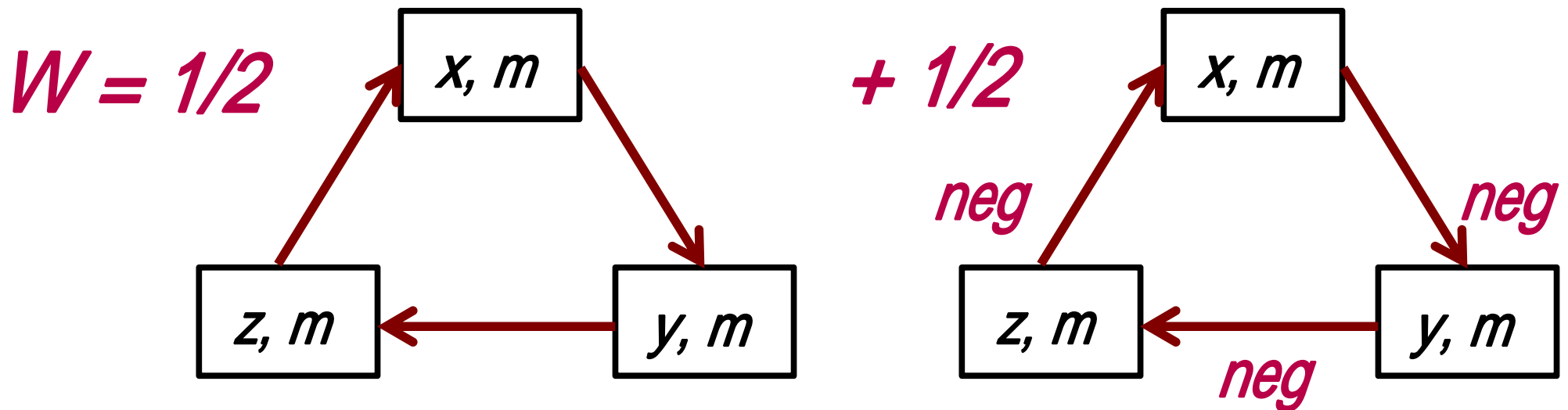


W allows for winning the game:

Each party sends the value she

Always Winning the Game with a Consistent W

Example: $m=0$

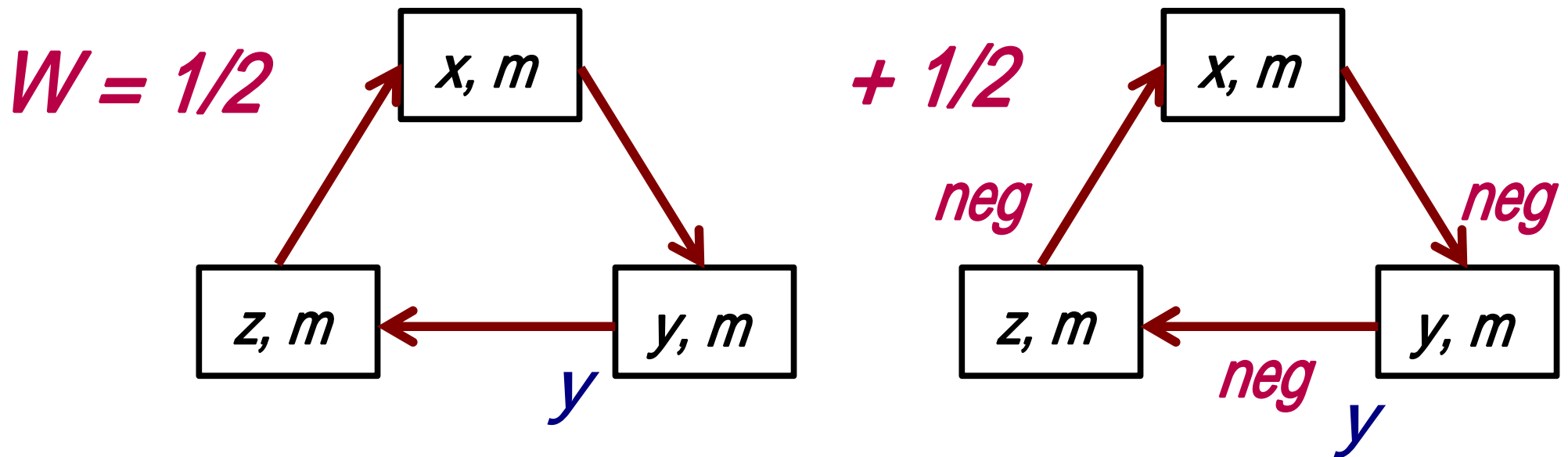


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Always Winning the Game with a Consistent W

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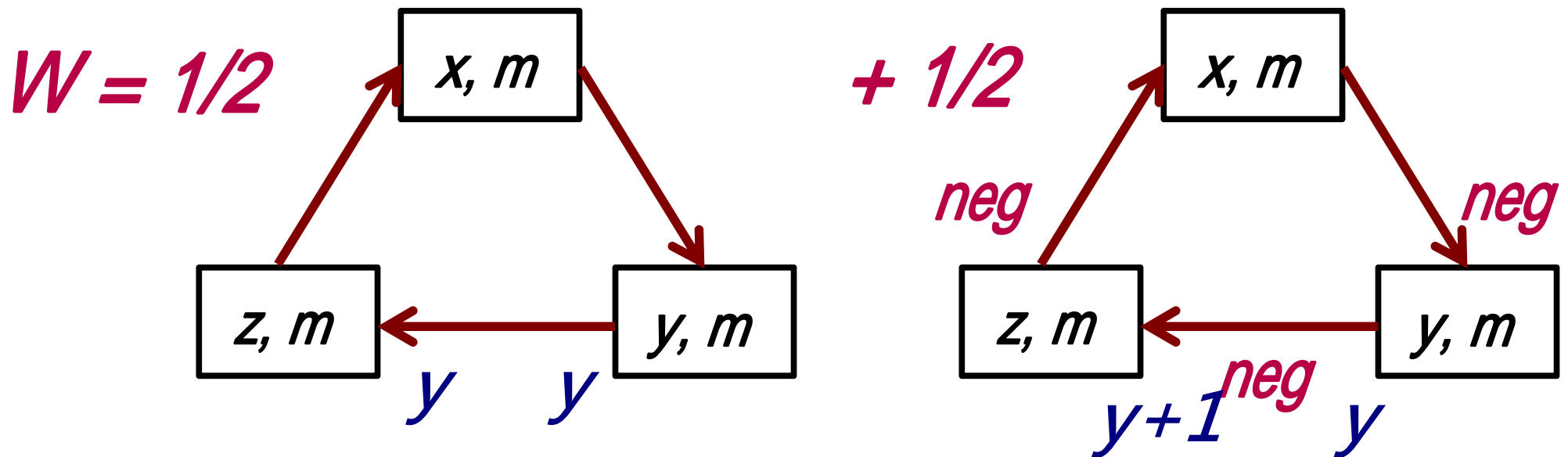


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Always Winning the Game with a Consistent W

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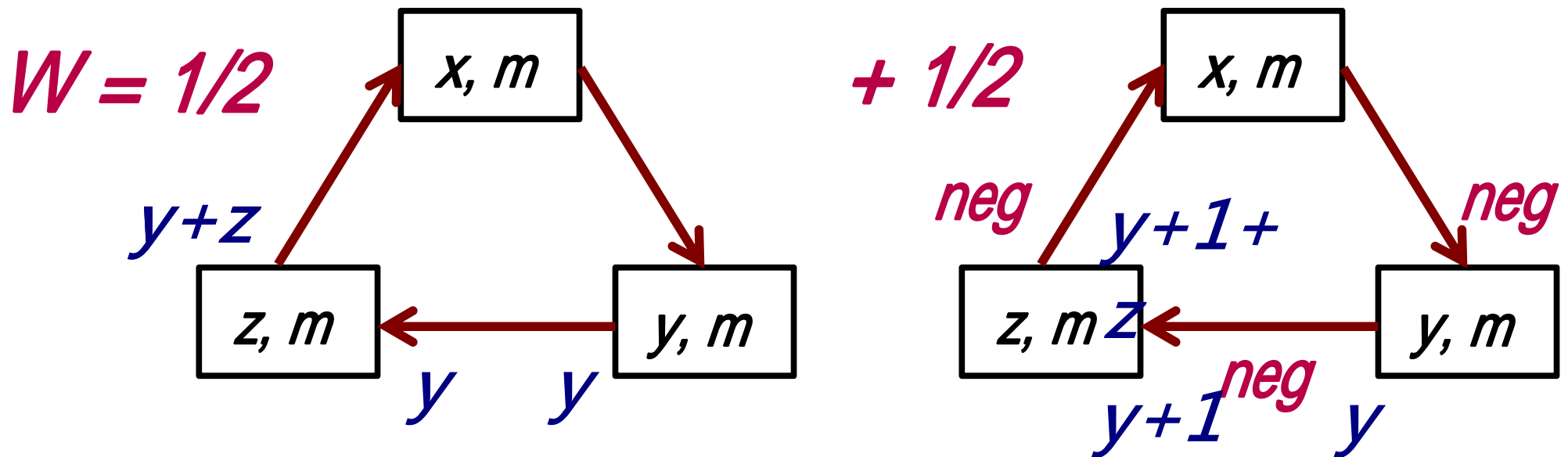


W allows for winning the
game:

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Always Winning the Game with a Consistent W

Example: $m=0$

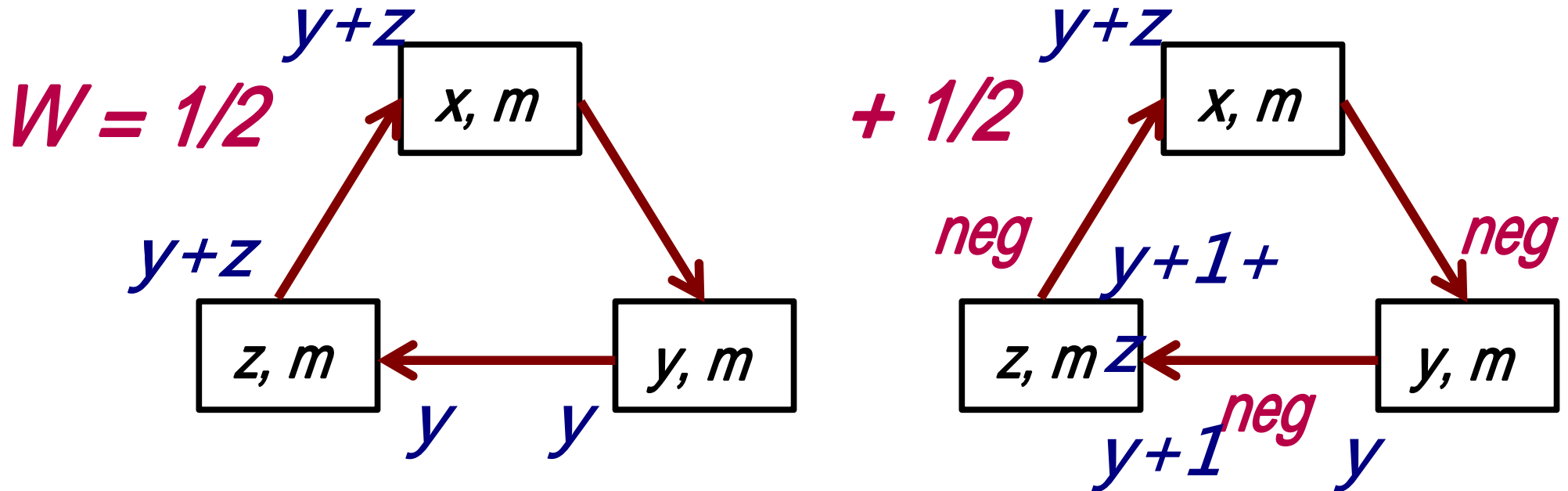


W allows for winning the game:

Each party sends the value she

Always Winning the Game with a Consistent W

Example: $m=0$. The game is won.



W allows for winning the game:

Each party sends the value she

Back to Randomness

How to define it intrinsically?

Back to Randomness

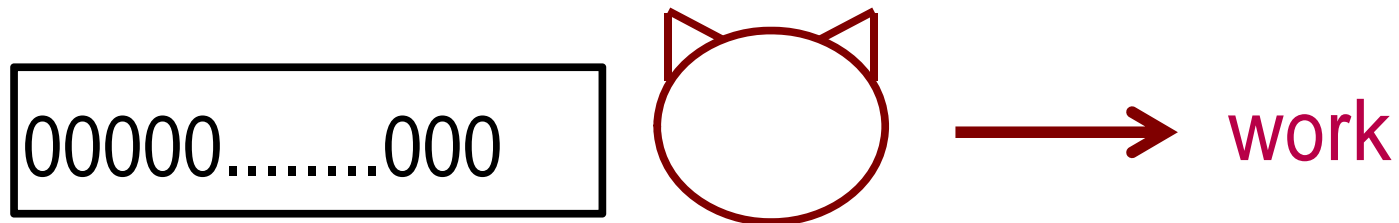
How to define it intrinsically?

00000.....000

is *not* random

Back to Randomness

How to define it intrinsically?



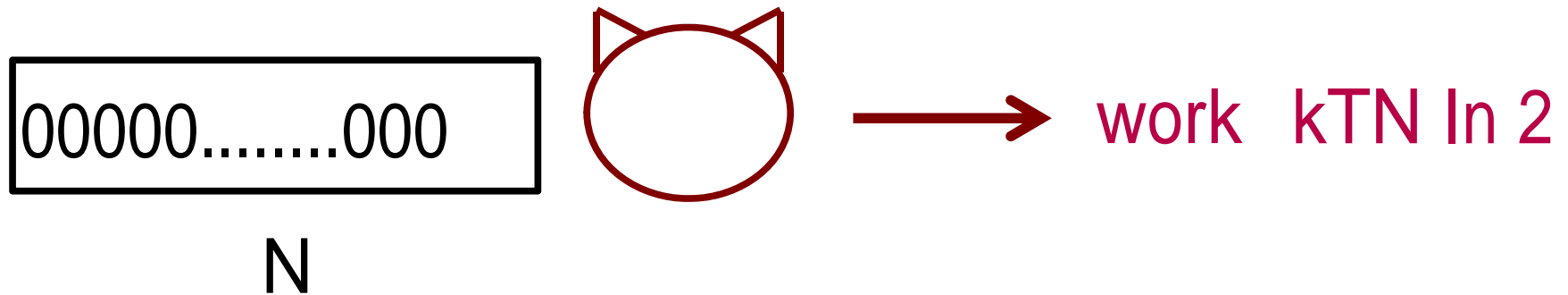
Charles H. Bennett



Rolf Landauer

Back to Randomness

How to define it intrinsically?



Charles H. Bennett



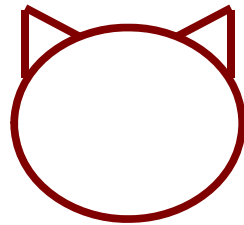
Rolf Landauer

Back to Randomness

How to define it intrinsically?

00000.....000

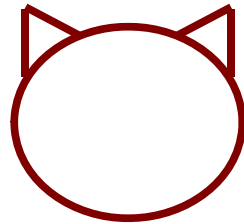
N



work $kTN \ln 2$

31415926535..

N



work $kTN \ln 10$



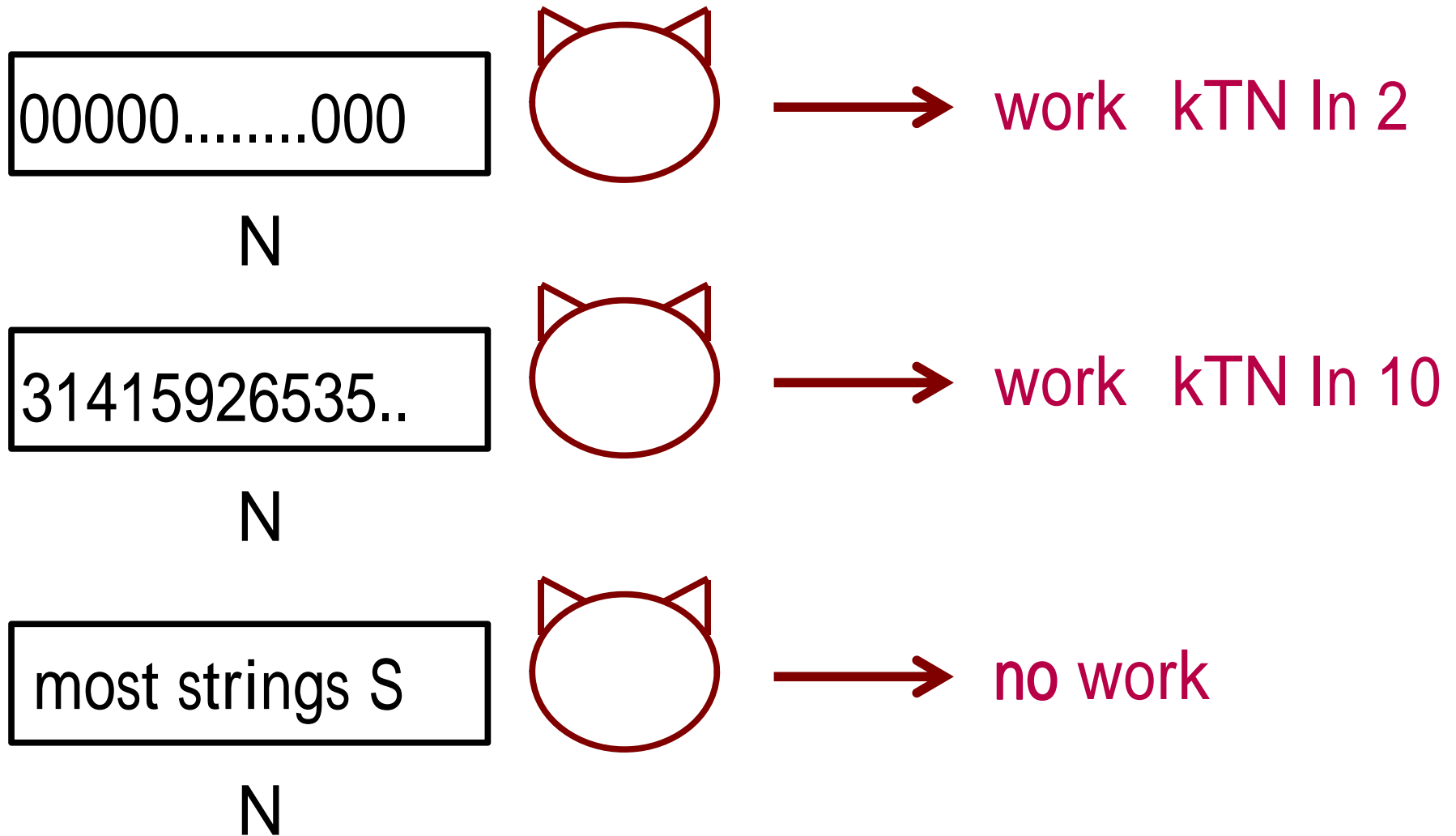
Charles H. Bennett



Rolf Landauer

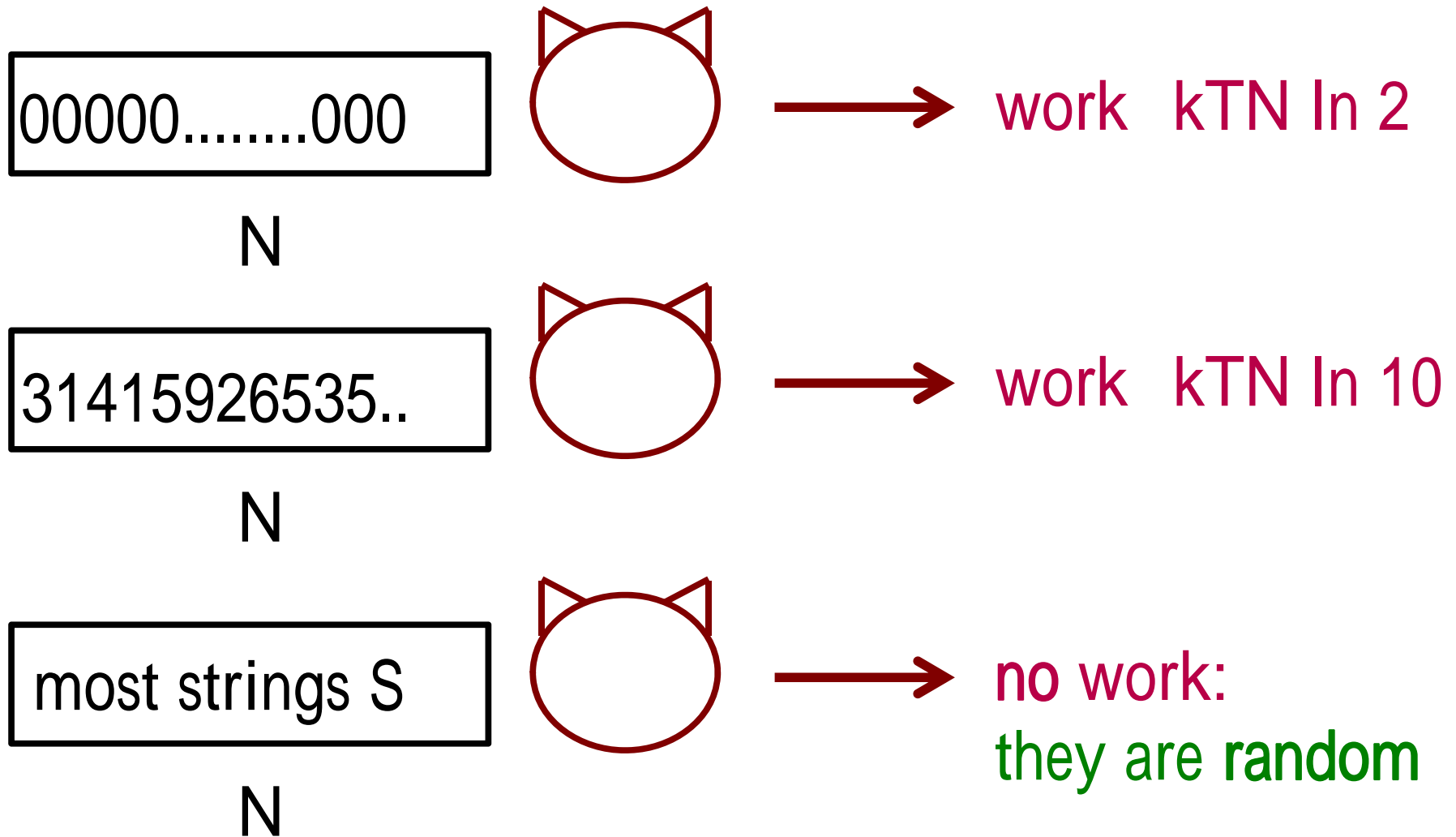
Back to Randomness

How to define it intrinsically?

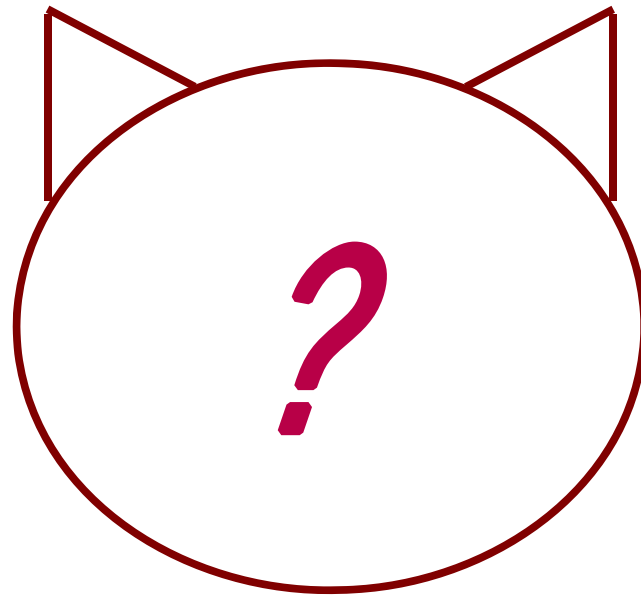


Back to Randomness

How to define it intrinsically?



Back to Randomness



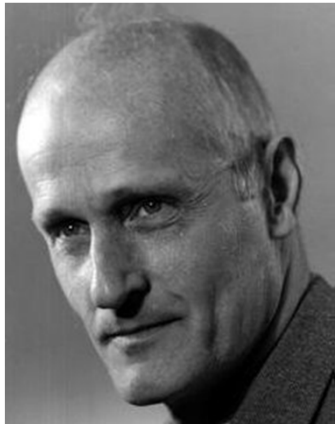
Back to Randomness



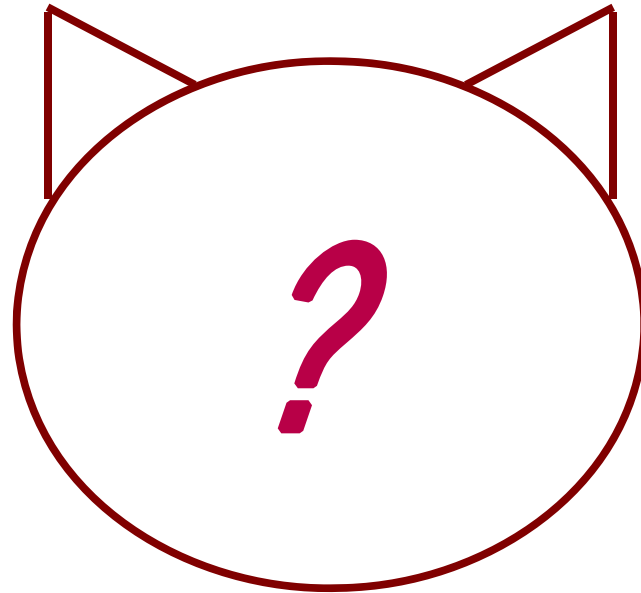
Alan Turing
1936



Alonzo Church



1943
Stephen Kleene



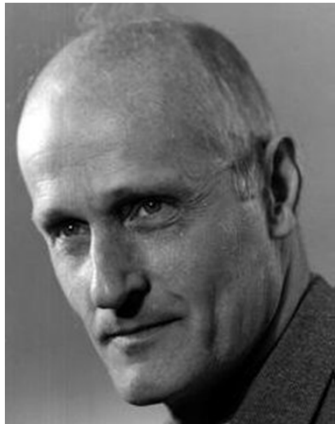
Back to Randomness



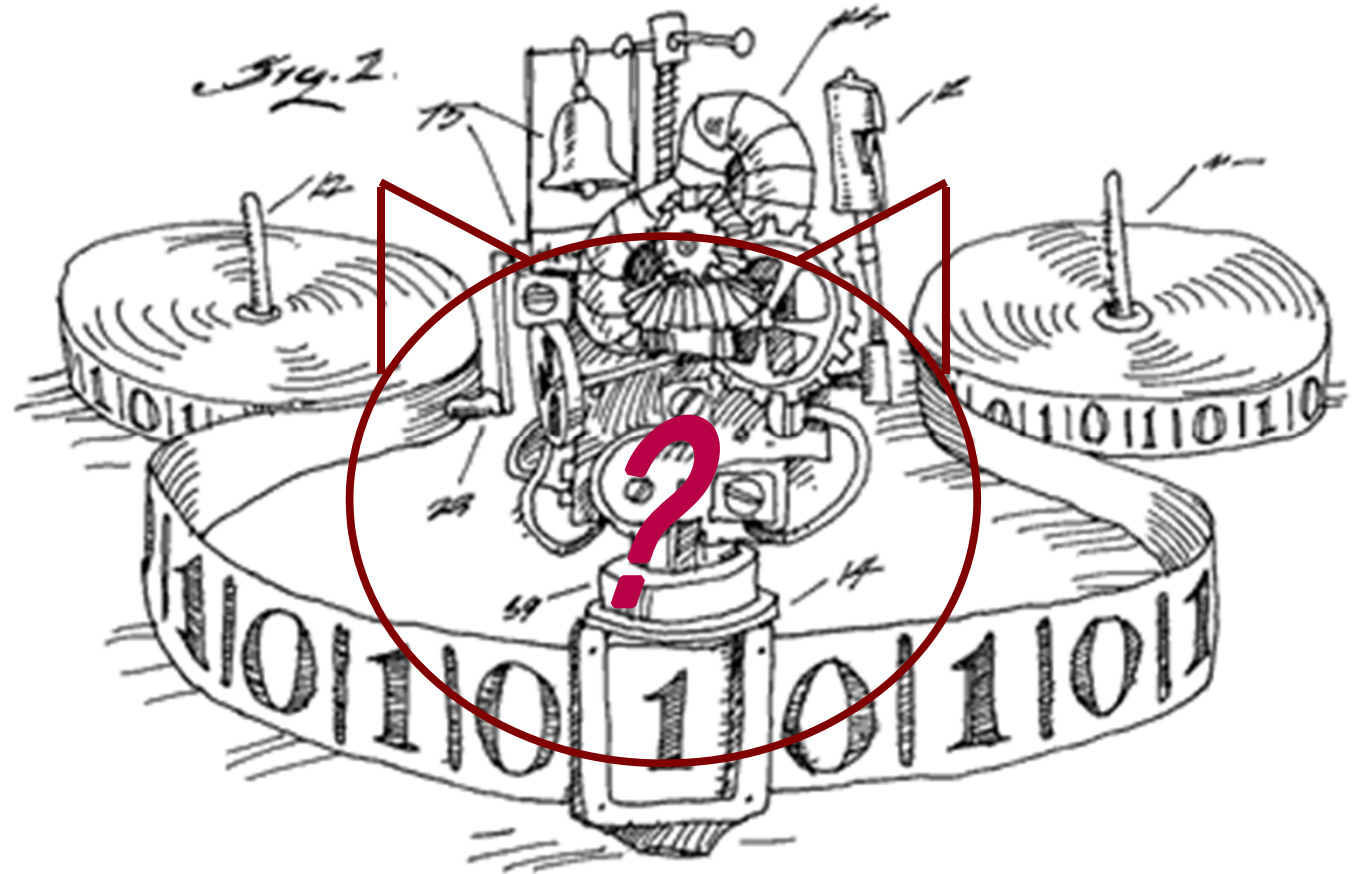
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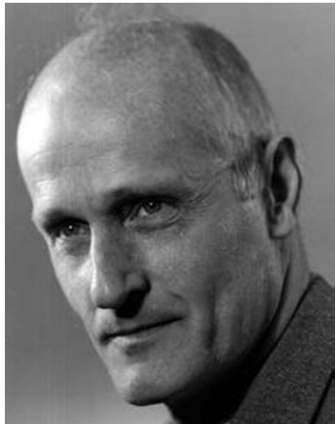
Back to Randomness



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1936



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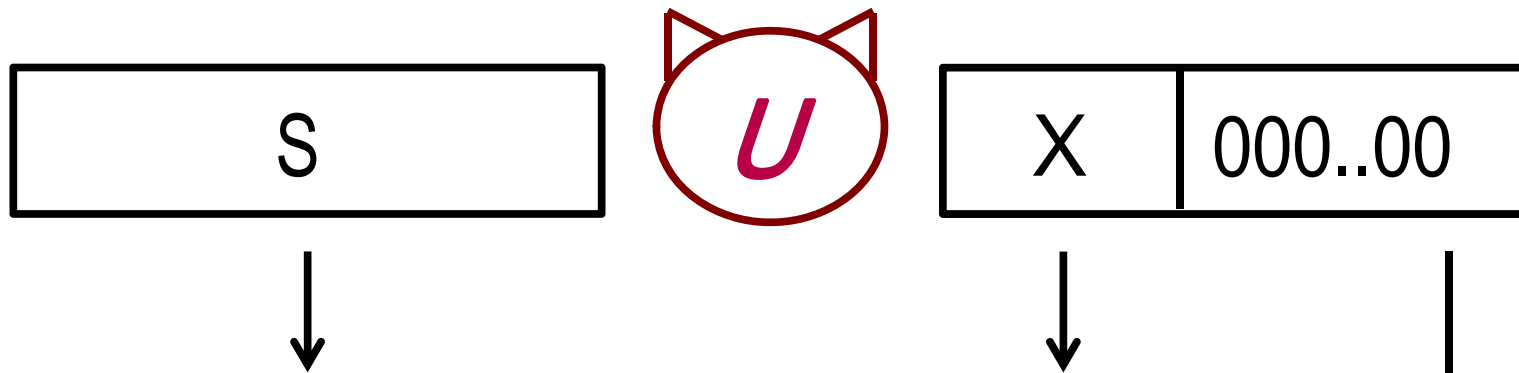


1943
Stephen Kleene



Work Extraction

The model

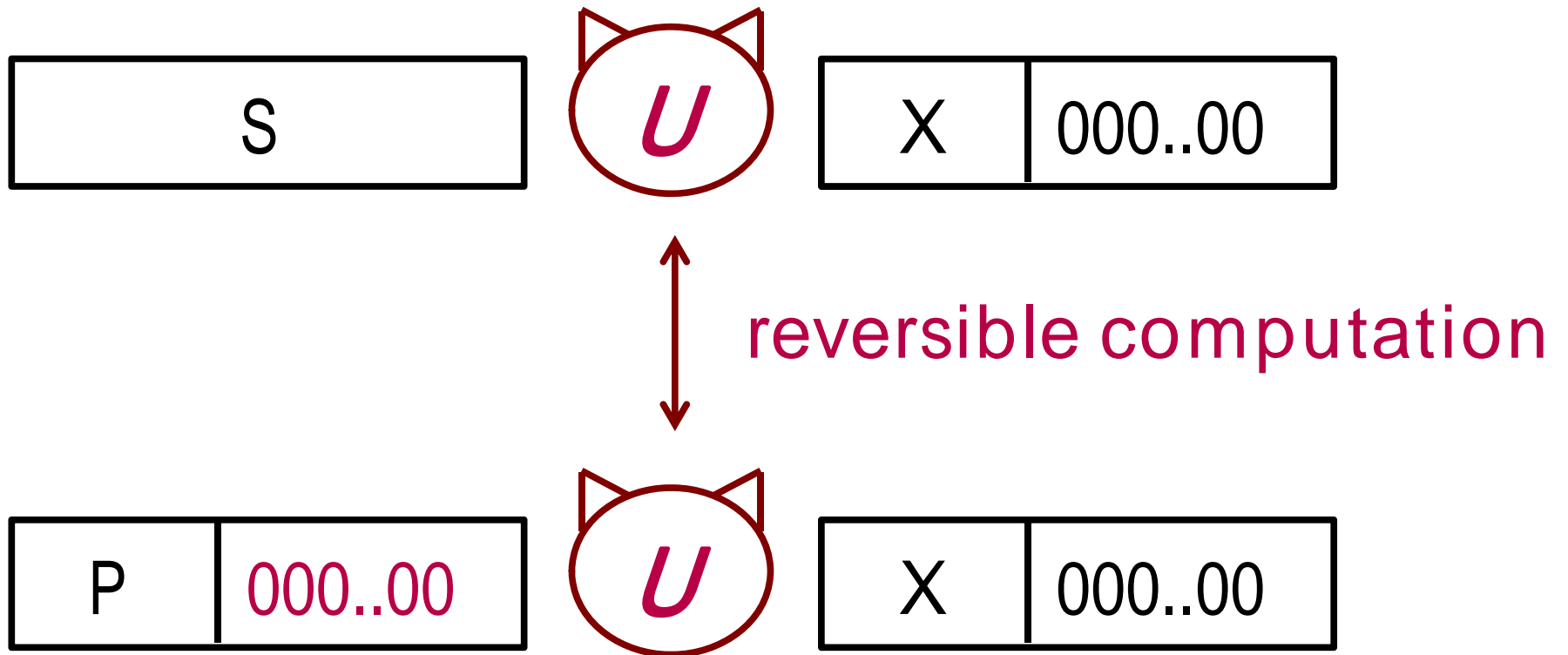


finite string from which work is to be extracted knowledge about S (initial state)

↓
finite, sufficiently long tape

Work Extraction

The model



$W(S|X)$

→ work: $W(S|X) kT \ln 2$

Bounds on the Work Value

Upper bound:

$$W(S|X) \leq \text{len}(S) - K_U(S|X)$$

|

*length of the shortest
program
for U to compute S given X*



Charles H. Bennett

Andrei Kolmogorov



Bounds on the Work Value

Upper bound:

$$W(S|X) \leq \text{len}(S) - K_{\mathcal{U}}(S|X)$$



Charles H. Bennett

Lower bound:

C compression algorithm with helper, *i.e.*,

$C : S||X \mapsto C(S, X)||X$ reversible:

$$W(S|X) \geq \text{len}(S) - \text{len}(C(S, X))$$



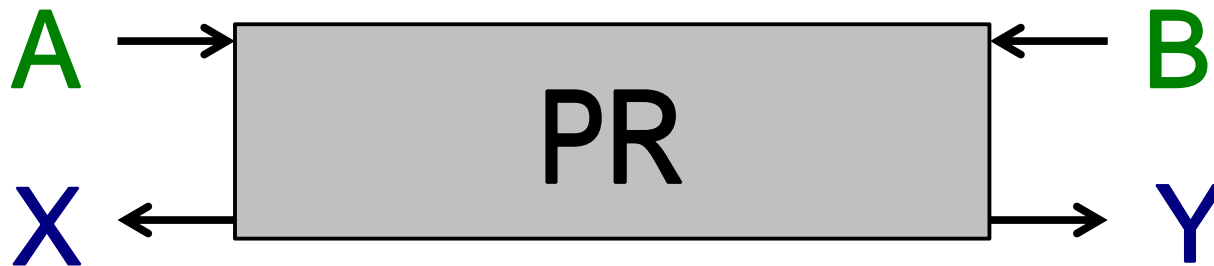
Wojciech Zurek

Back to Non-Local Correlations

Back to Non-Local Correlations *Counterfactual reasoning*

Consequence of non-

signaling:
If all (A,B) combinations are possible...



... then X and Y must be perfectly
random

Back to Non-Local Correlations

Factual-only reasoning



$$x_i \oplus y_i = a_i \odot b_i$$

Back to Non-Local

Correlations

Factual-only reasoning

(a,b) incompressible:

$$\frac{K(a^n, b^n)}{2n} \longrightarrow 1$$

for $n \longrightarrow \infty$ and $a^n := (a_1, \dots, a_n)$



$$x_i \oplus y_i = a_i \odot b_i$$

Back to Non-Local

Correlations

Factual-only reasoning

non-signaling

$$\frac{K(x^n | a^n)}{K(x^n | a^n, b^n)} \longrightarrow 1$$

for $n \longrightarrow \infty$, and symmetric

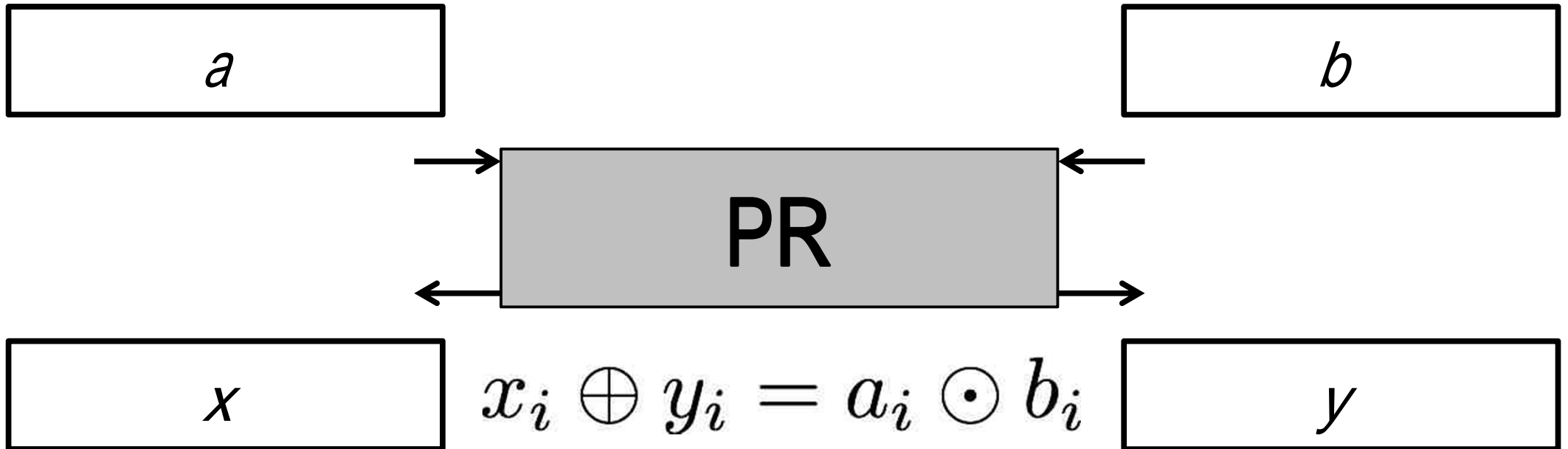


$$x_i \oplus y_i = a_i \odot b_i$$

Back to Non-Local

Correlations

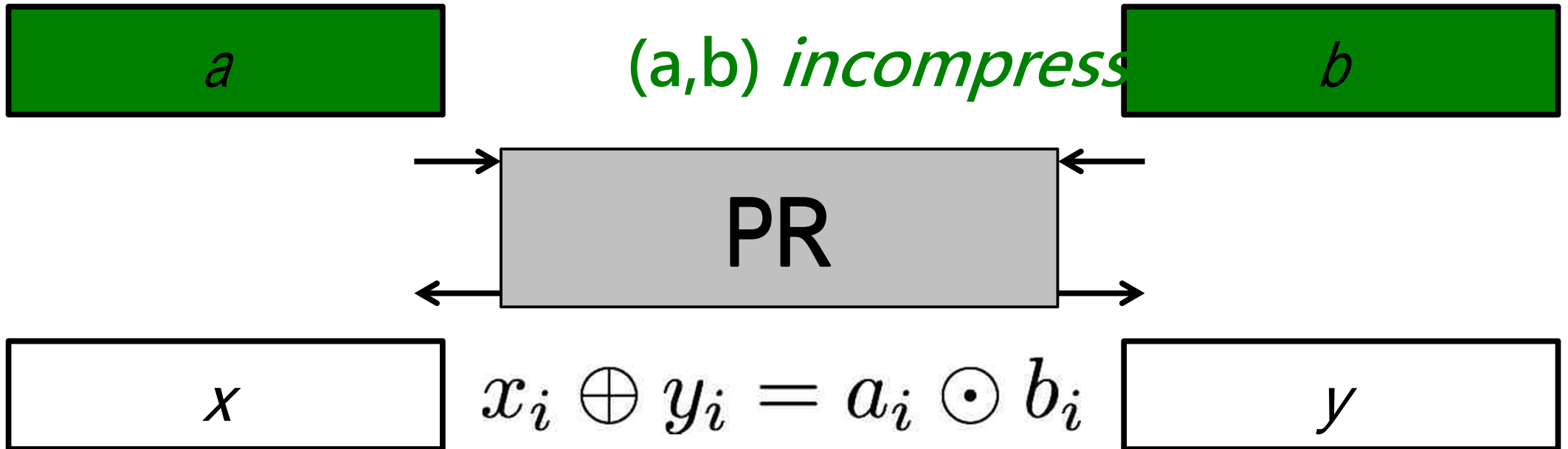
Factual-only reasoning



Back to Non-Local

Correlations

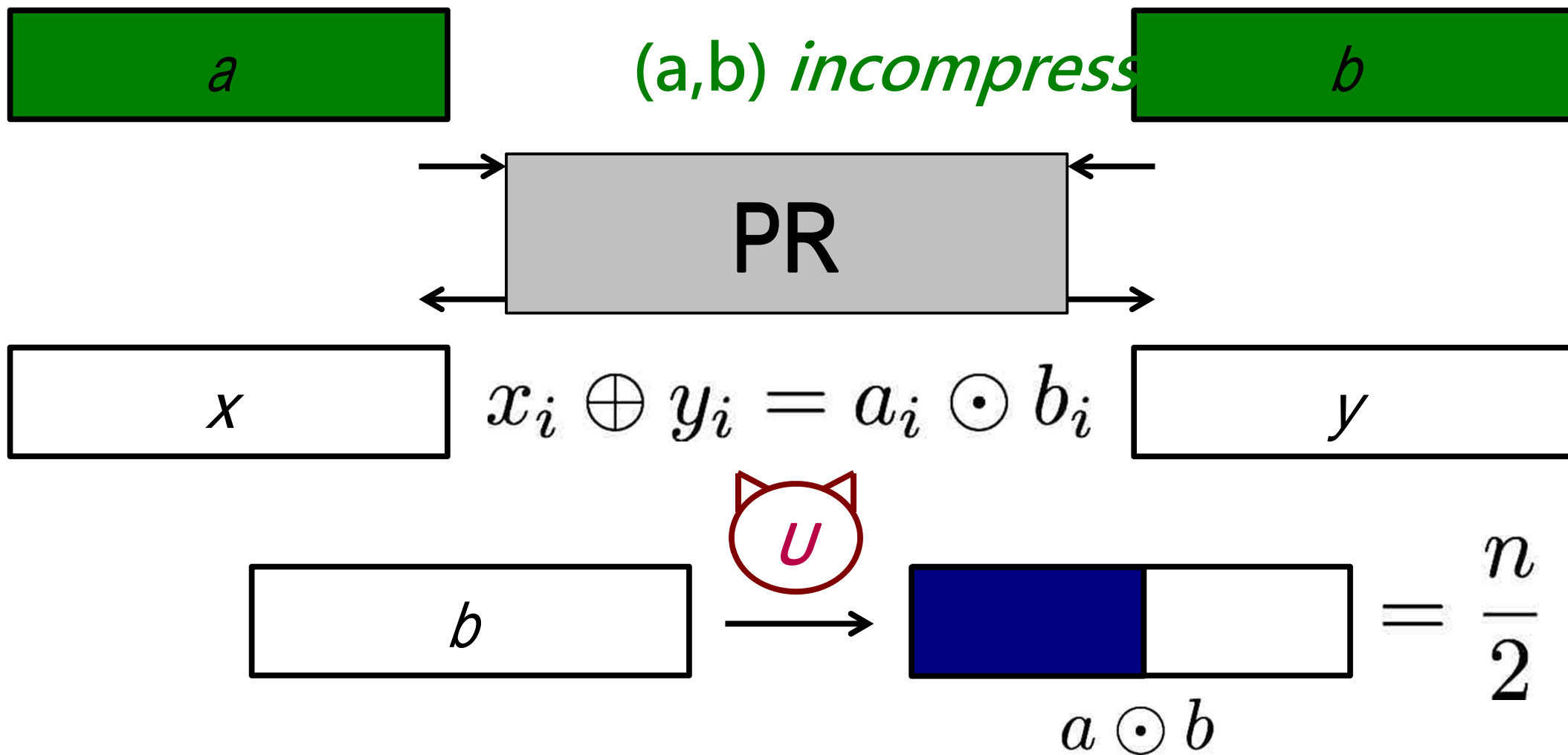
Factual-only reasoning



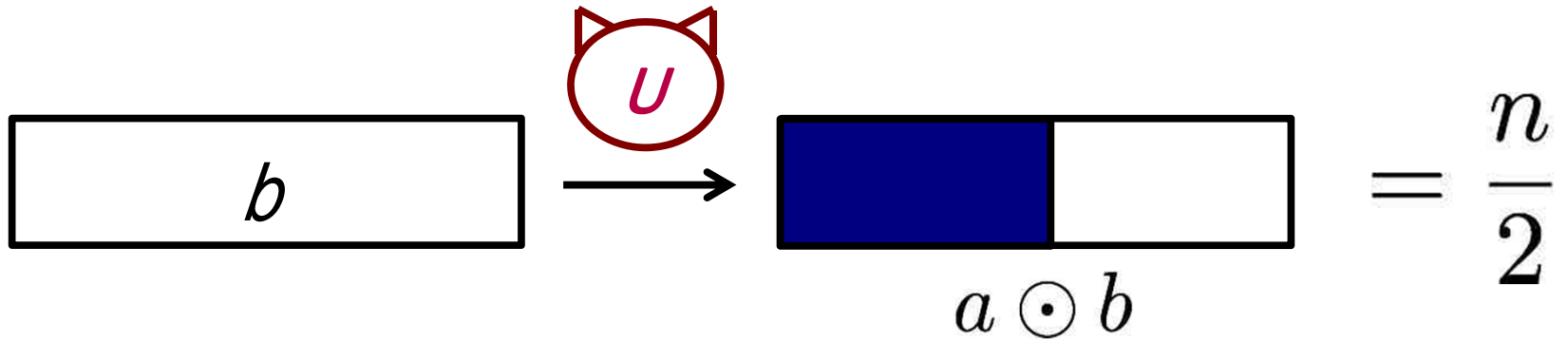
Back to Non-Local

Correlations

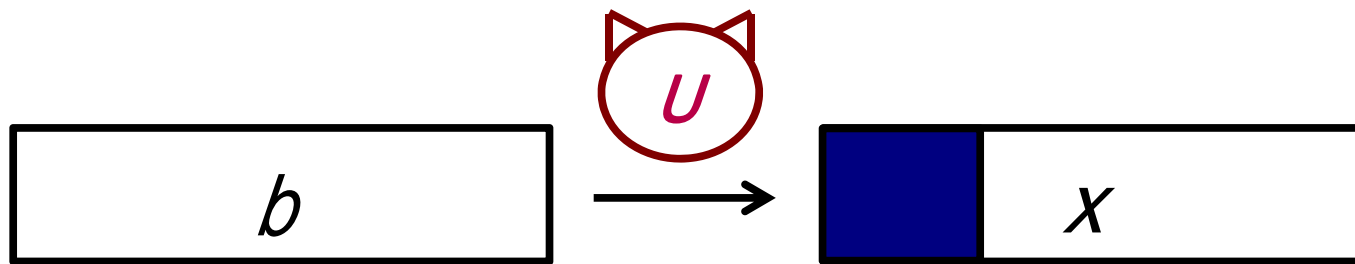
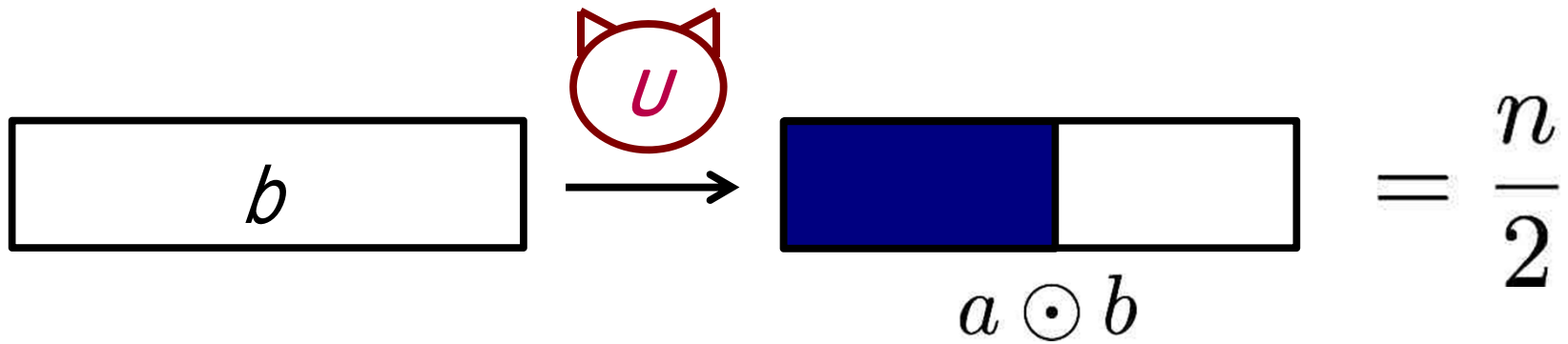
Factual-only reasoning



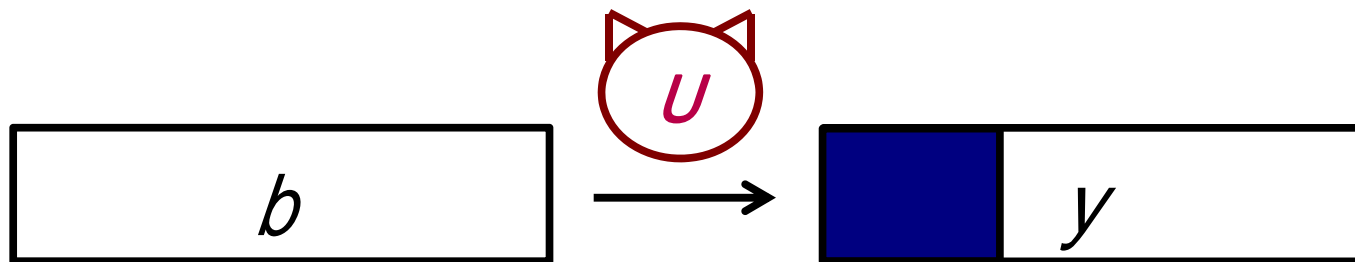
$$x_i \oplus y_i = a_i \odot b_i$$



$$x_i \oplus y_i = a_i \odot b_i$$

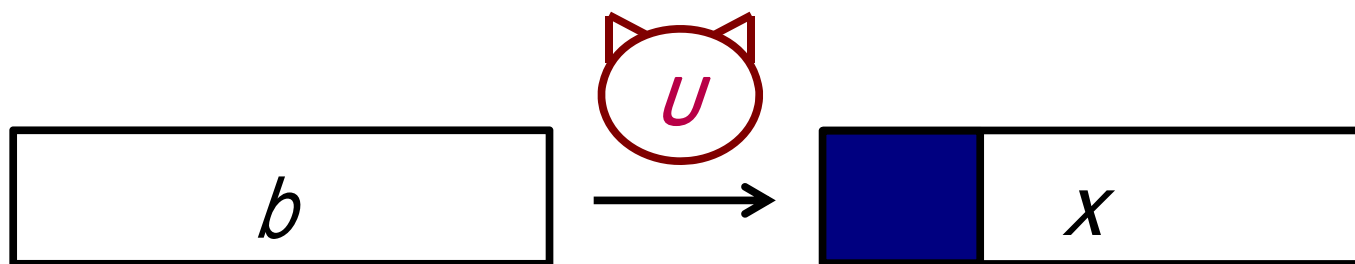


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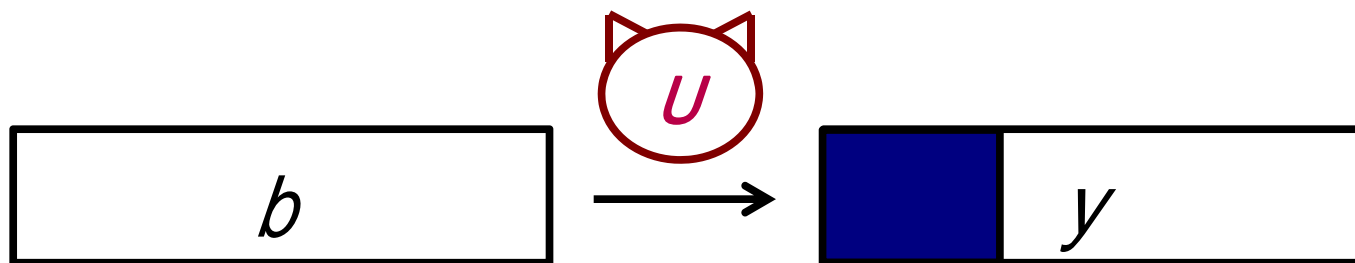


$\geq \frac{n}{2}$

$$x_i \oplus y_i = a_i \odot b_i$$

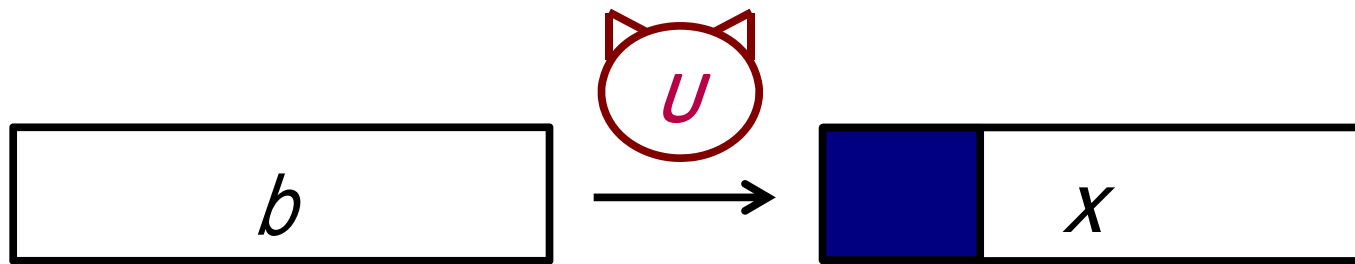


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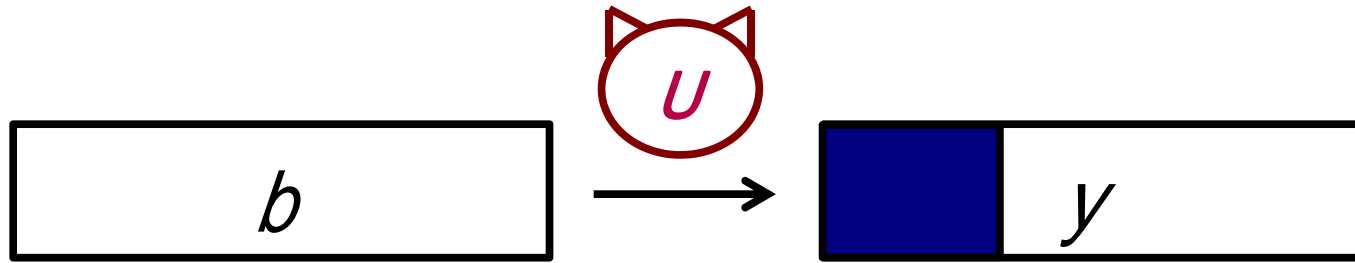


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$$x_i \oplus y_i = a_i \odot b_i$$

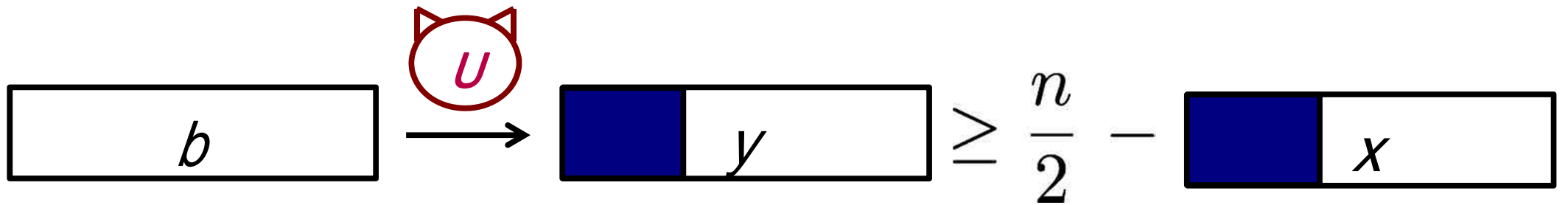


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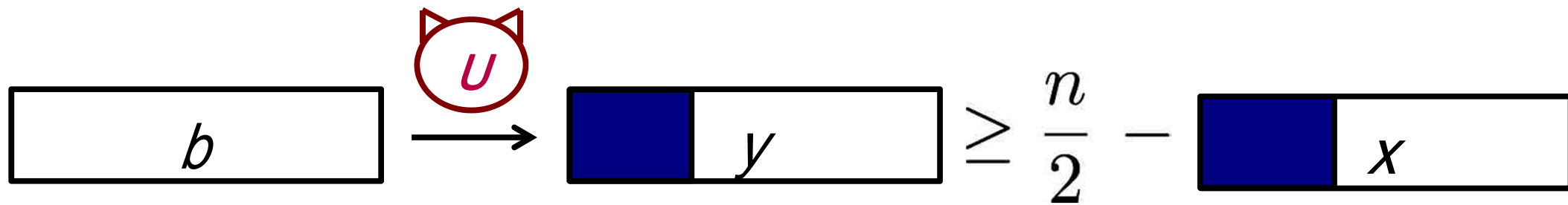


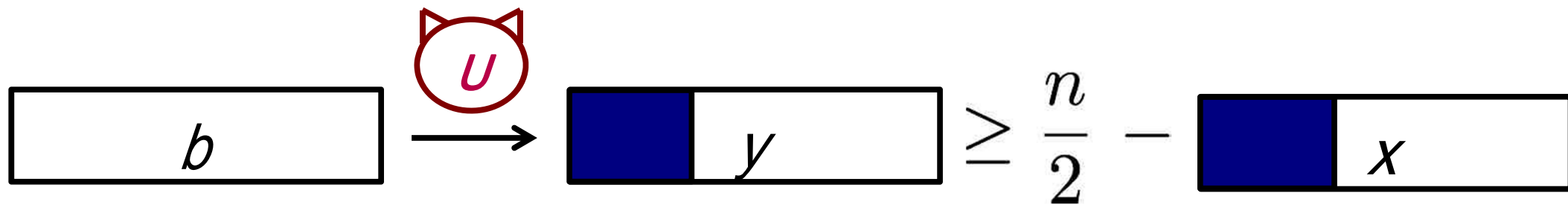
$\geq \frac{n}{2}$

\Rightarrow

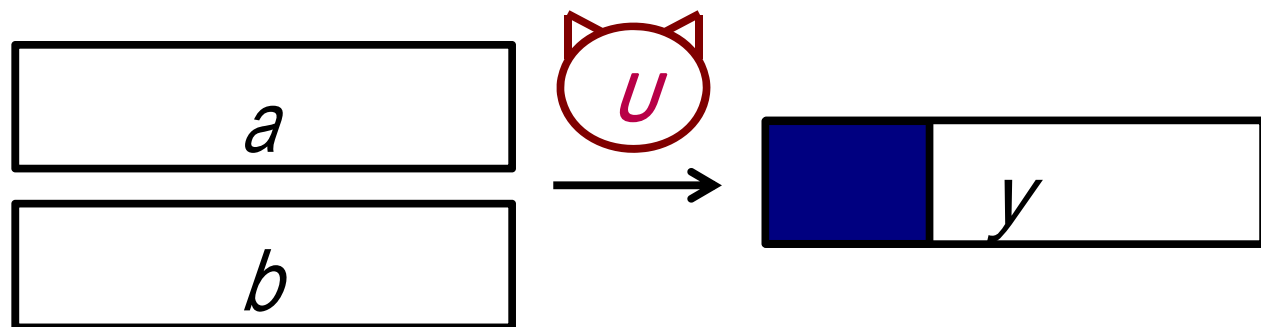


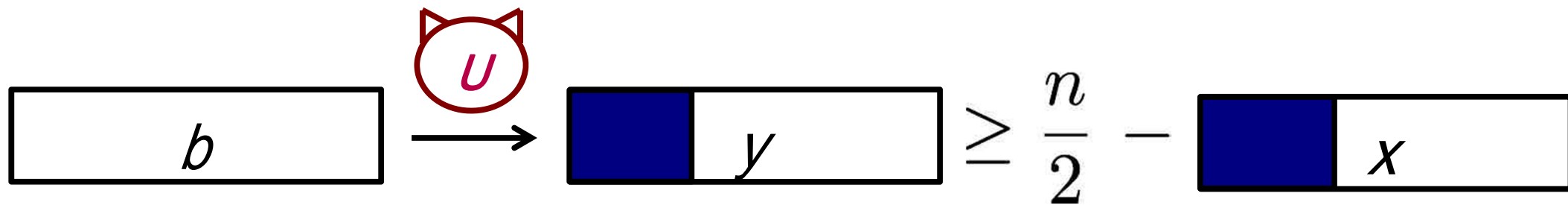
$\geq \frac{n}{2}$



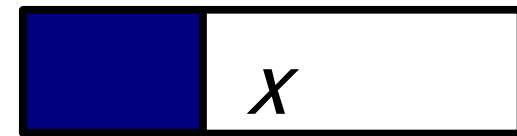


On the other hand:

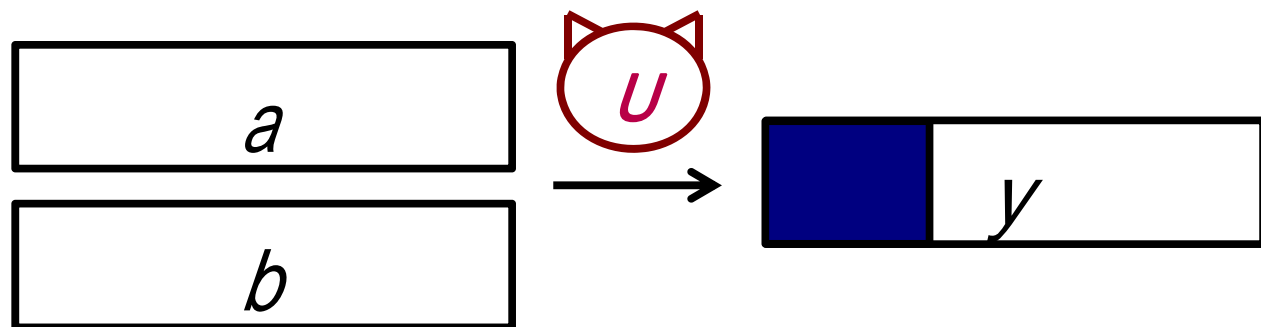




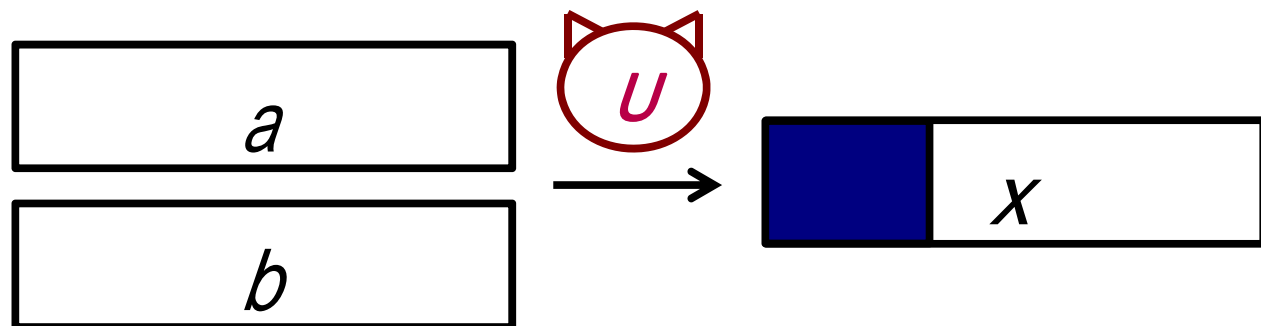
$$\geq \frac{n}{2} -$$



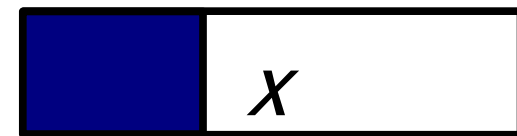
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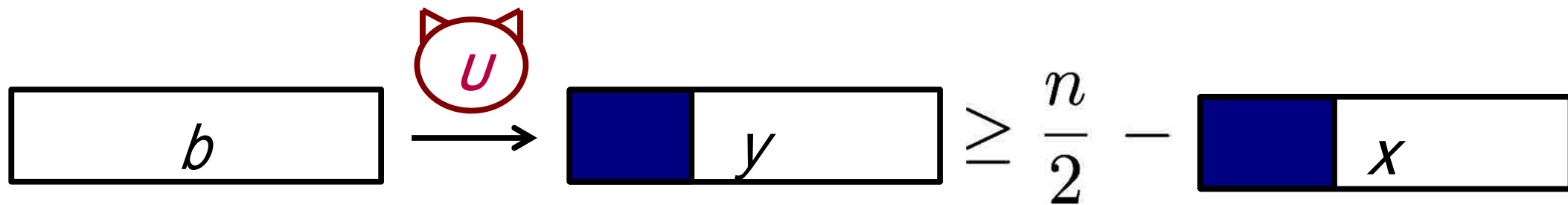


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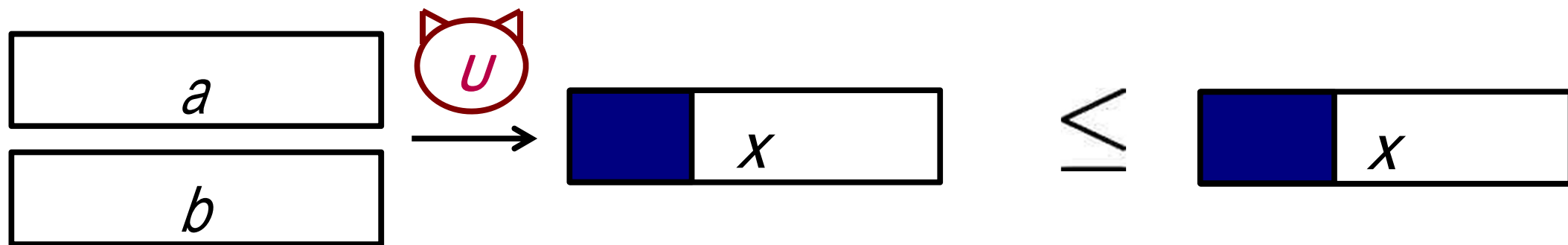
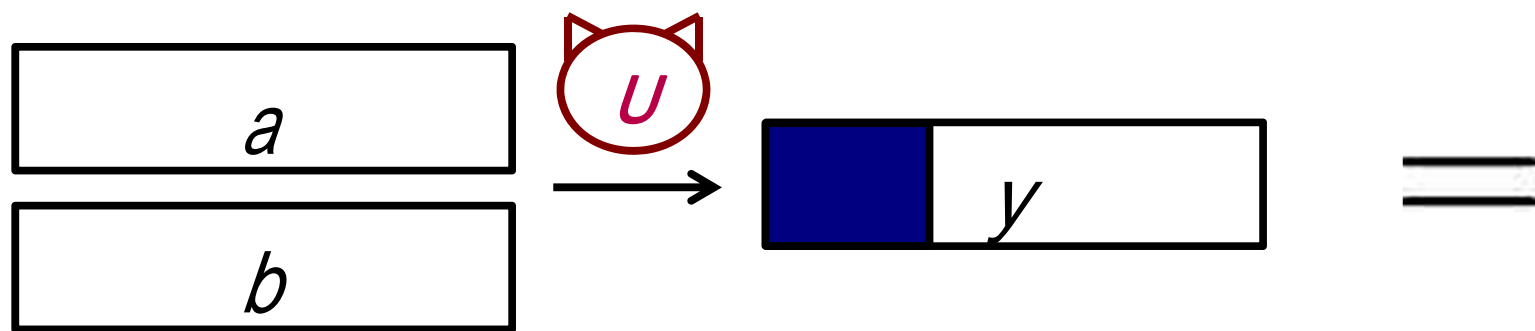


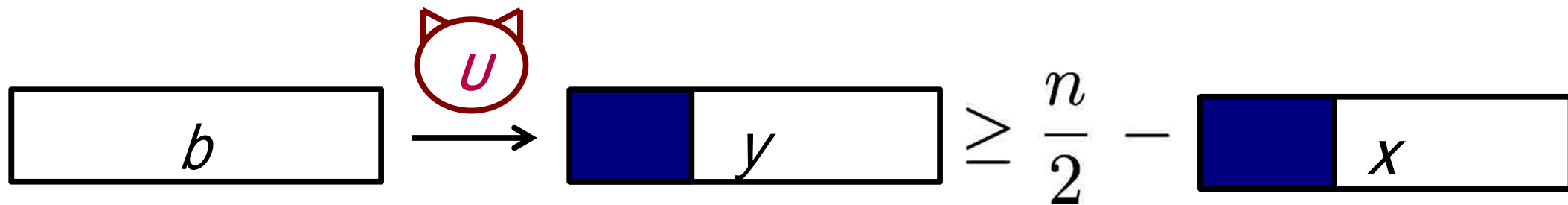
$$\geq$$



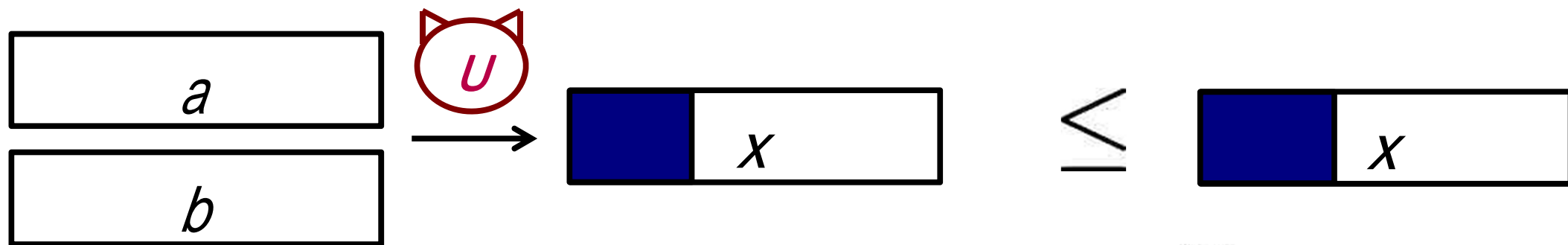
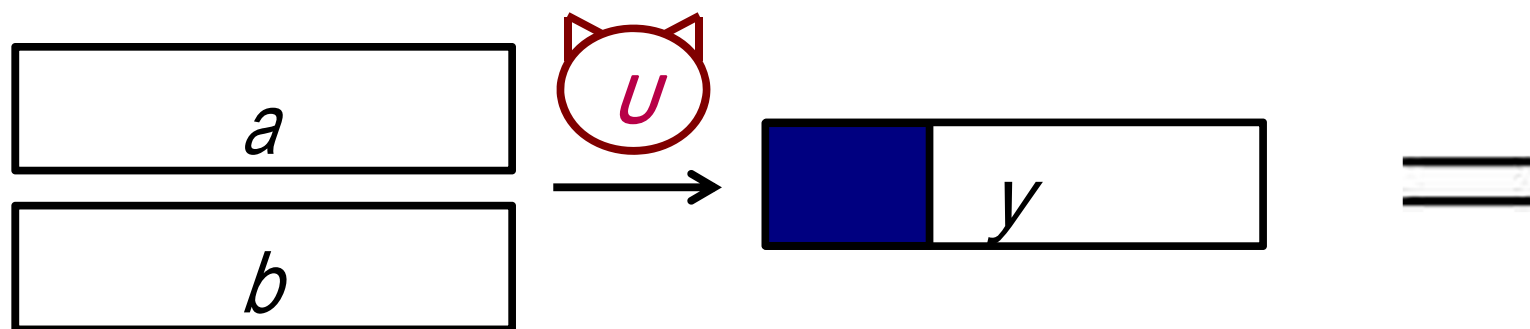


Non-signaling:

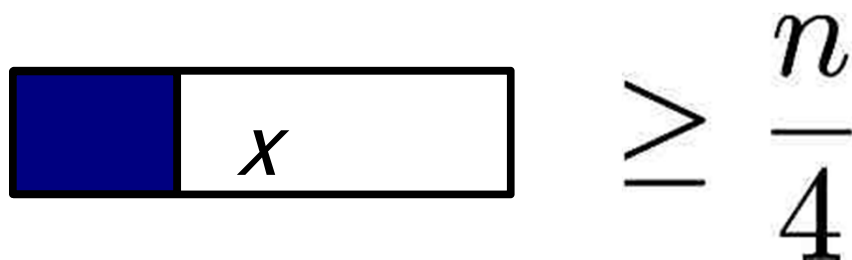




Non-signaling:



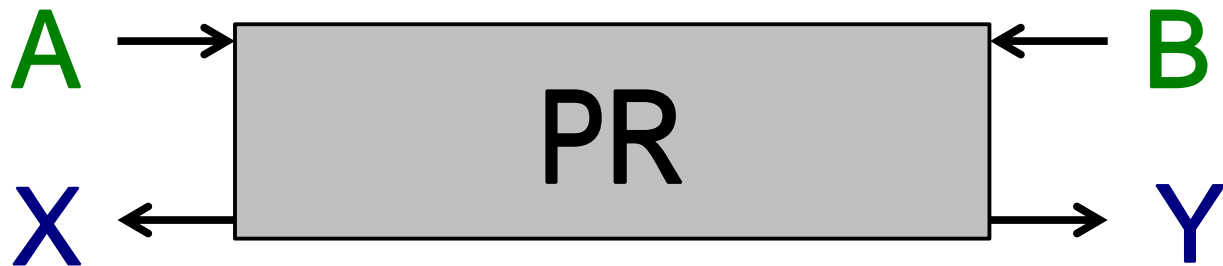
Therefore:



Non-Local Correlations

Factual-only reasoning

If (A,B) is incompressible...



Non-Local Correlations

Factual-only reasoning

If (A,B) is incompressible...



... then X and Y cannot be computable

Non-Local Correlations

Factual-only reasoning

If (A,B) is incompressible...



... then X and Y cannot be computable

Non-Local Correlations

Factual-only reasoning

If (A,B) is incompressible...



... then X and Y cannot be computable
even given the respective inputs
from parallel-repetition theorem

Ran Raz 1998



Non-Local Correlations

All-or-Nothing Feature of Church-Turing Hypothesis

Beyond TM



Beyond TM

Non-Local Correlations

All-or-Nothing Feature of Church-Turing Hypothesis

Beyond TM

Beyond TM



Beyond TM

Beyond TM

Non-Local Correlations

All-or-Nothing Feature of Church-Turing Hypothesis

Beyond TM

Beyond TM

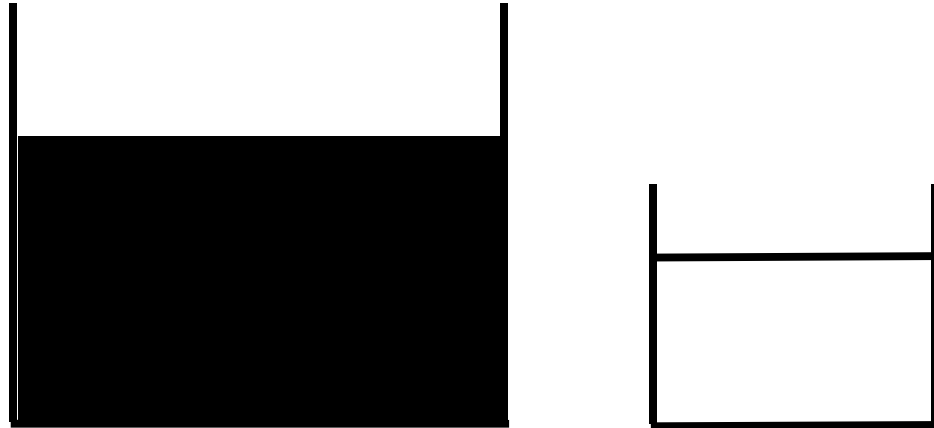


Beyond TM

Beyond TM

If the experimenter can generate (maximally) uncomputable data, then so can the measured photons

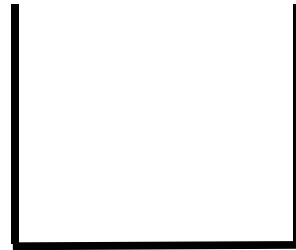
Epilogue: Complexity and Time Asymmetry



Carroll
Aaronson



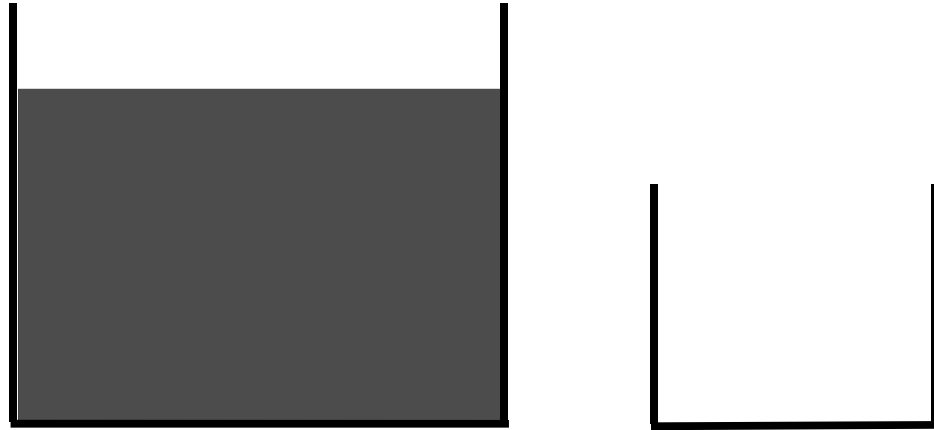
Epilogue: Complexity and Time Asymmetry



Carroll
Aaronson



Epilogue: Complexity and Time Asymmetry



Carroll
Aaronson



Epilogue: Complexity and Time Asymmetry

Which quantity is monotonic in time? Macrostate size?

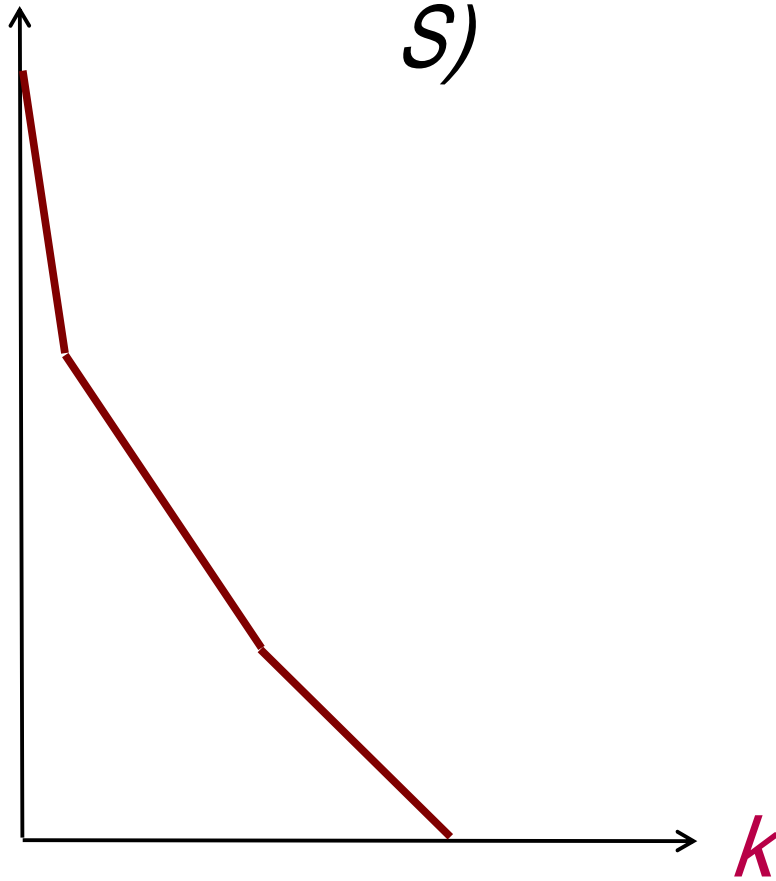


Which is not? Structure?

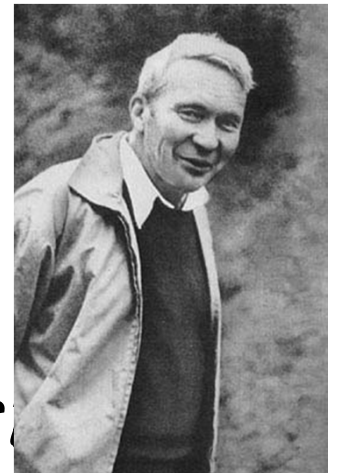
Epilogue: Complexity and Time Asymmetry

Context-free macrostate (of a finite string S)

*log-size of
smallest set
 $M(k)$ with
complexity
 k
containing S*



Kolmogorov sufficient statis

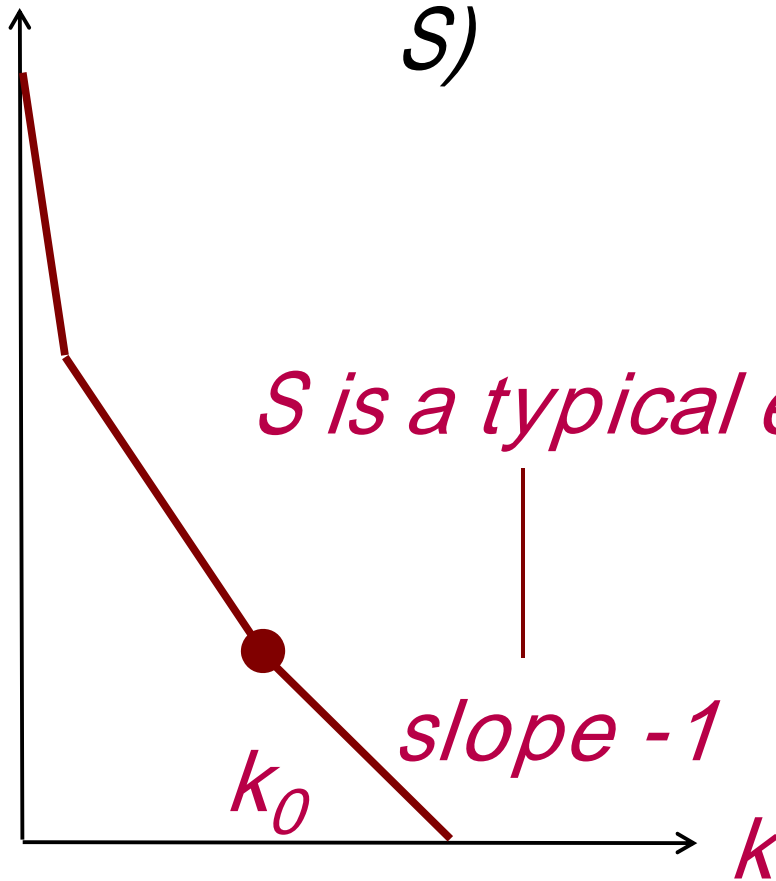


Epilogue: Complexity and Time Asymmetry

Context-free macrostate (of a finite string

S)

*log-size of
smallest set
 $M(k)$ with
complexity
 k
containing S*



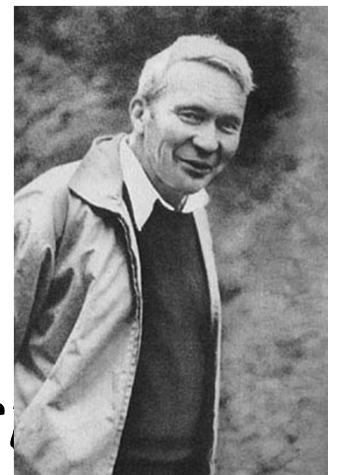
S is a typical element of $M(k_0)$

slope -1

k_0

k

Kolmogorov sufficient statis



Epilogue: Complexity and Time Asymmetry

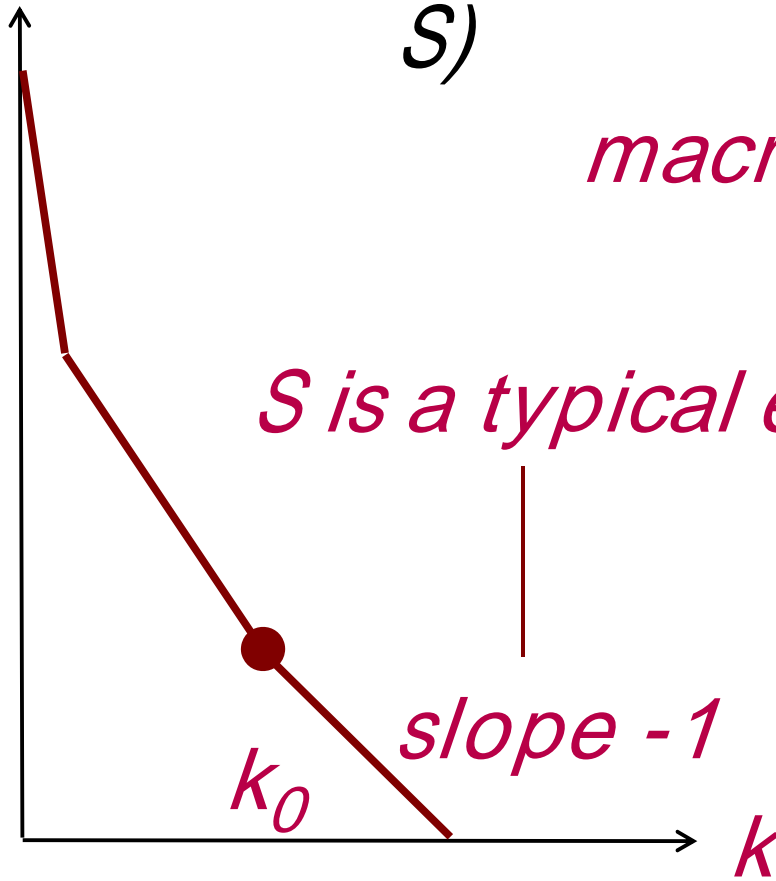
Context-free macrostate (of a finite string

S)

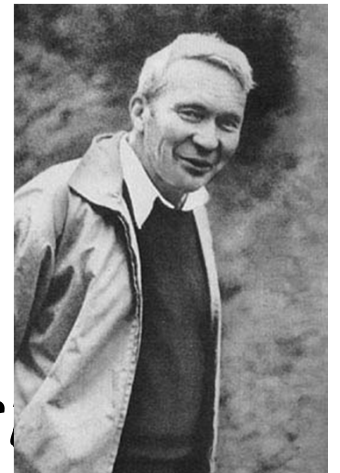
macrostate: $M(k)$

*log-size of
smallest set
 $M(k)$ with
complexity
 k
containing S*

S is a typical element of $M(k_0)$

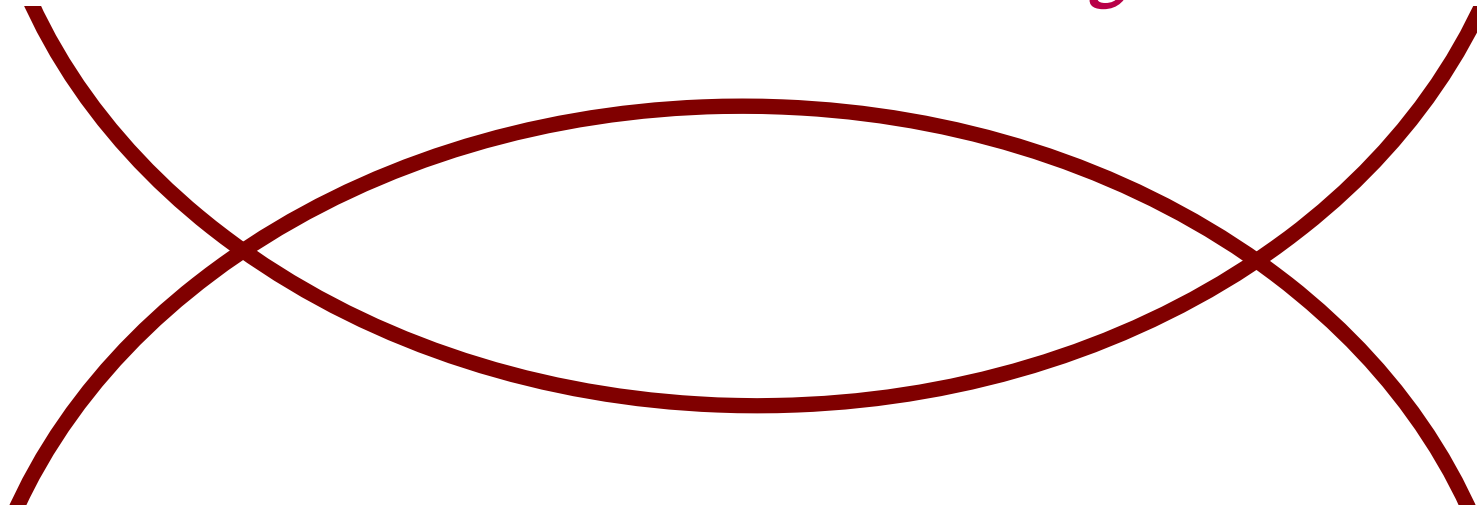


Kolmogorov sufficient statis



Epilogue: Complexity and Time Asymmetry

log-size of macrostate



complexity of macrostate

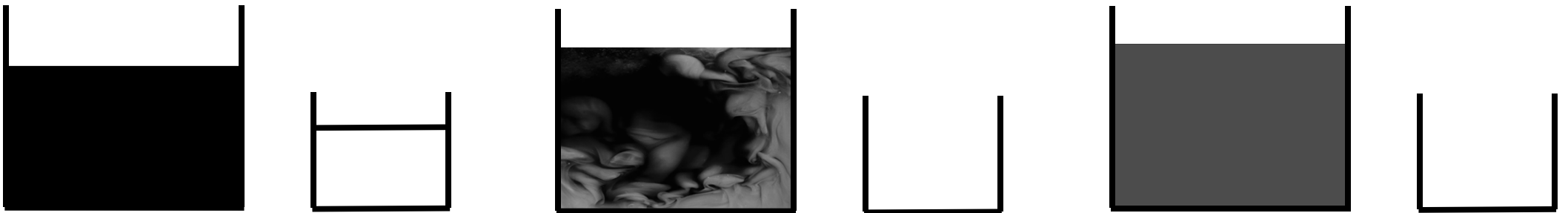


Epilogue: Complexity and Time Asymmetry

complexity of microstate

log-size of macrostate

complexity of macrostate



Epilogue: Complexity and Time Asymmetry

Observations

logical reversibility \implies $K(\text{tape/microstate})$
non-decreasing

Epilogue: Complexity and Time Asymmetry

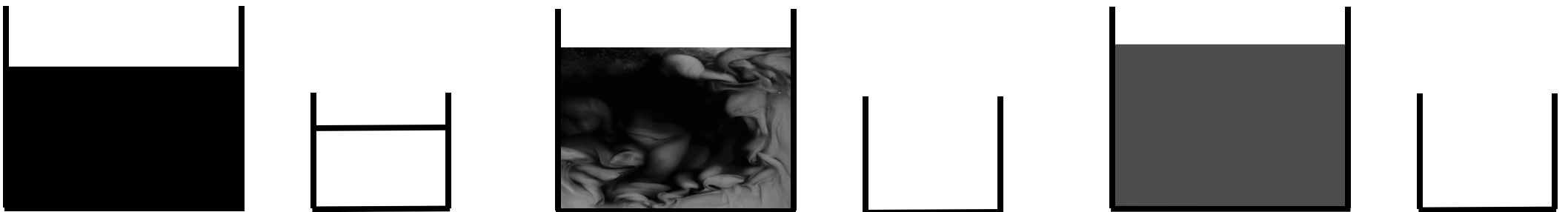
Observations

logical reversibility \implies $K(\text{tape/microstate})$
non-decreasing

complexity of microstate

probabilistic TM

deterministic TM



Epilogue: Complexity and Time Asymmetry

Observations

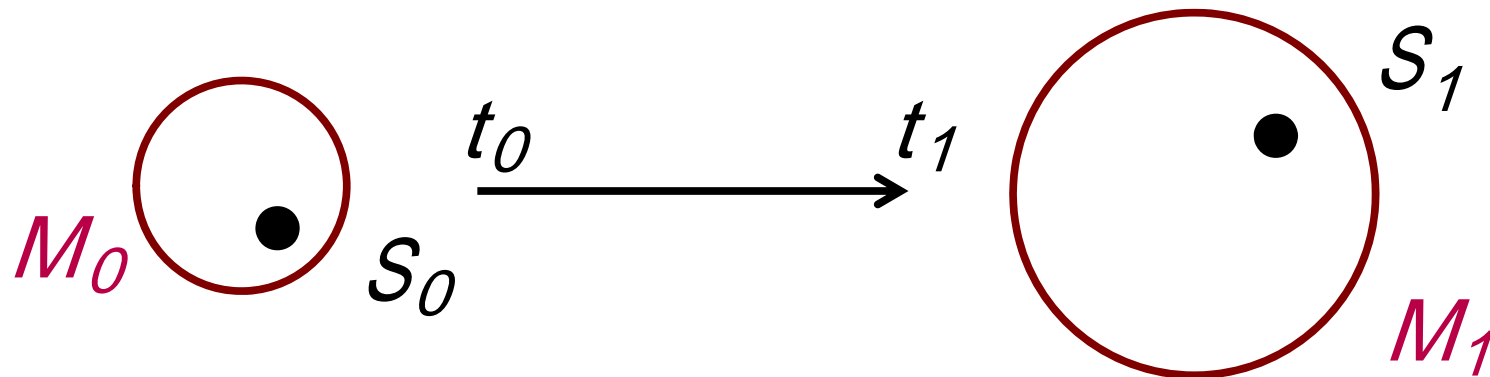
logical reversibility \implies $K(\text{tape}/\text{microstate})$
non-decreasing

Epilogue: Complexity and Time Asymmetry

Observations

logical **reversibility** \implies $K(\text{tape}/\text{microstate})$
non-decreasing

\implies simple **macrostate** not shrinking

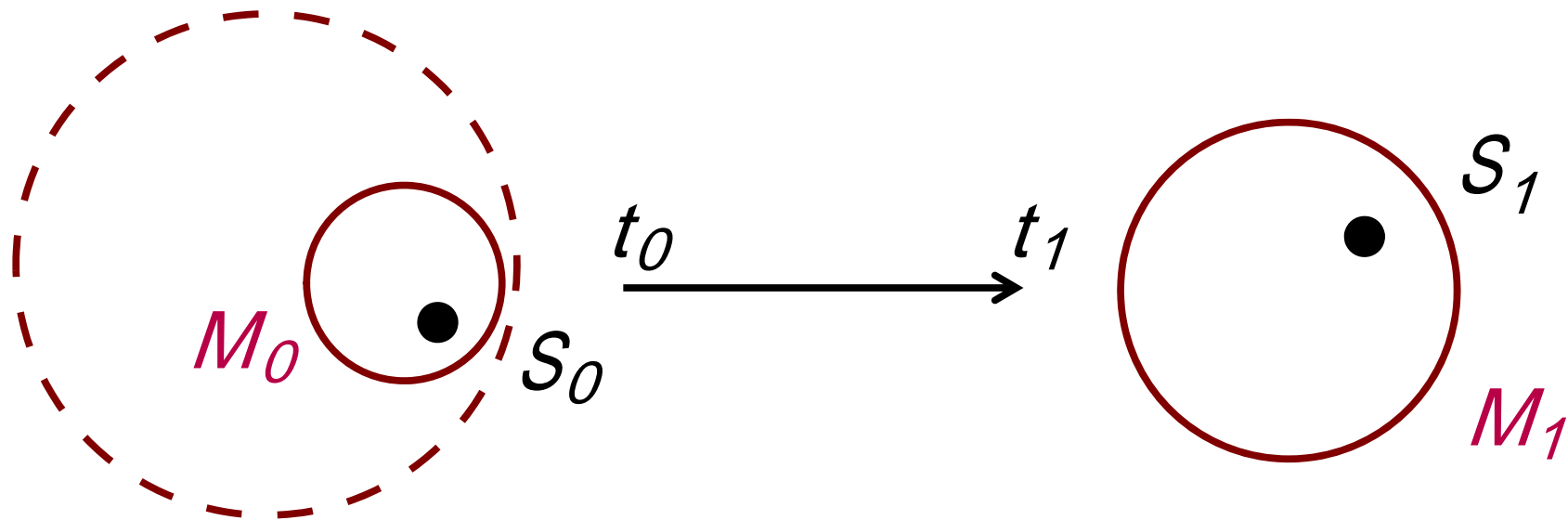


Epilogue: Complexity and Time Asymmetry

Observations

logical **reversibility** \implies $K(\text{tape}/\text{microstate})$
non-decreasing

\implies simple **macrostate** not shrinking





*Based on collaboration and discussion with:
Mateus Araújo, Veronika Baumann, Ämin Baumeler,
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Paul Erker, Adrien Feix, Jürg Fröhlich, Nicolas Gisin,
Esther Hänggi, Arne Hansen, Marcus Huber,
Alberto Montina, Benno Salwey, and Andreas Winter*